CONTINGENT PRICING POLICIES

by

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DEPARTMENT OF STATISTICS

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CHAPTER I
MULTIDIMENSIONAL INCENTIVE CONTRACTS IN
DEVELOPMENT PROJECTS

Introduction

An Air Force spokesman succinctly characterized the complexities of government procurement when he stated:

"...we are supposed to buy at the lowest sound price, on a competitive basis, but still make sure we support the operational force as required, and on time, with quality parts, at the same time giving consideration to small business and labor distressed areas and without introducing too many nonstandard parts into the inventory." [1]

One of the tools employed in the effort to achieve efficient development of advanced systems is the multiple incentive contract. Although the Department of Defense (DOD) and the National Aeronautics and Space Administration (NASA) are not the sole users of contractual incentives, the bulk of government procurement by dollar volume is performed by DOD and NASA, and these agencies have pioneered in the use of new contractual forms designed to increase the effectiveness of the procurement dollar. Consequently the examples and analysis in this paper will be most pertinent to the problems and practices of the defense and aerospace industries.

The contract types that we will be concerned with are Cost-Plus-Fixed-Fee (CPFF), Cost-Plus-Incentive-Fee (CPIF), Fixed-Price-Incentive (FPI), Firm-Fixed-Price (FFP), and Cost-Plus-Award-Fee (CPAF). We
briefly define the distinguishing features of each contract type. In a CPFF contract the government and contractor negotiate a fixed contractor fee and a target cost. The government is obligated to reimburse the contractor for all allowable costs incurred in executing the contract, whatever the cost outcome. An FFP contract defines the task to be accomplished by the contractor and the total price to be paid by the government for this effort. Any cost variation from this negotiated price is the sole responsibility of the contractor. In both CPIF and FPI contracts a target cost, target fee, and sharing arrangement are negotiated. At the completion of the project, contractor and government together share cost variation according to the negotiated formula. For example, if the sharing formula is 80/20 and the actual cost outcome is $1 million less that target cost, the contractor would receive an additional $200,000, and the government would retain the remaining $800,000. In the CPIF contract the government benefits from all cost underruns below some floor, and is fully responsible for all overruns beyond some ceiling. The FPI contract contains a ceiling beyond which the sharing formula becomes 0/100, and therefore the contractor assumes more risk with an FPI contract than under a CPIF contract. This distinction is more apparent than real since cost outcomes rarely penetrate FPI ceilings. In recent FPI contracts the cost ceilings have been as low as 112-115% of target cost.\footnote{Almost all contracts with multiple (performance and development schedule) incentives have been CPIF contracts, and DOD policy} [4] Almost all contracts with multiple (performance and development schedule) incentives have been CPIF contracts, and DOD policy
has not permitted performance or schedule incentives to be written without cost incentives operating simultaneously.\(^2\)

CPF contracts are cost reimbursement contracts negotiated with a basic minimum fee considerably lower than CPF fees. Based on the customer's evaluation of contractor performance an additional award fee of up to 10% can be earned.

Table 1 shows the trend of DOD procurement by contract type from fiscal year 1952 through fiscal year 1964. [5] The increased use of CPF contracts in the mid-1950's coincided with the accelerated

<table>
<thead>
<tr>
<th>Fiscal Year</th>
<th>CPF</th>
<th>CPIF/FPI</th>
</tr>
</thead>
<tbody>
<tr>
<td>52</td>
<td>17.9%</td>
<td>12.0%</td>
</tr>
<tr>
<td>53</td>
<td>18.0%</td>
<td>26.2%</td>
</tr>
<tr>
<td>54</td>
<td>27.0%</td>
<td>27.7%</td>
</tr>
<tr>
<td>55</td>
<td>22.7%</td>
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<td>19.0%</td>
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<td>58</td>
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</tr>
<tr>
<td>64</td>
<td>12.0%</td>
<td>32.6%</td>
</tr>
</tbody>
</table>

\(^2\) We know of one instance in which a completed contract had included the schedule incentive in an initial letter contract. (A letter contract is issued to provide temporary contractual coverage so that urgent work can proceed while a definitive contract is being negotiated.) The urgency of the requirement had justified the use of a letter contract and the buying activity provided the contractor with early motivation necessary to meet emergency programmed requirements. The letter contract provided for both penalty and bonus points applied to target periods. The contractor earned 90% of the incentive reward and the buying activity secured a delivery schedule that had been refused by two other bidders and eight other potential sources. [14]
development of advanced weapons systems and the initiation of major space projects. Between 1955 and 1961 expenditures for development and procurement of strategic aircraft, missiles, and air defense systems averaged approximately $15 billion dollars annually, and during this period the proportion of CPFF awards increased from 22% to 38%. Peck has reported that 12 major CPFF projects completed during this period had an average cost overrun of 320% and an average schedule slippage of 36%. [7, p. 22] Procurement personnel in DOD state that CPFF contracts at best tended to run funds out to the targeted amount, and the term "horror cases" entered the lexicon of defense procurement in reference to some of the CPFF outcomes. [6]

By the end of 1961 the rate of expenditure for the development of new strategic systems stabilized at about $4 billion annually. DOD attention shifted to improving the capabilities of conventional forces, which required less technological development. This allowed the institution of a "Cost Reduction Program" emphasizing both management techniques and contractual innovations. [8] The cost reduction program aimed at conducting development and procurement in an incentive environment, one major element of which was the use of contractual incentives rather than cost reimbursement contracts. Unfortunately, in the effort to explain and justify the evolution in contracting, CPFF contracts were blamed for wasting resources, sapping industrial efficiency, dragging out project schedules, and providing no incentive for outstanding performance by the contractor. In fact, it is doubtful whether any contractor could have been induced to accept management responsibility for uncertain, expensive, and technologically advanced systems such as the Atlas missile
without government assuming the financial risk. Large cost overruns were in many cases due to customer uncertainty regarding requirements, rather than poor initial cost estimates. [9] The much criticized time lag caused by the lack of prior planning and system definition has now been remedied by explicitly defining the project objectives in a separate contractual effort that may take up to six months. [10] The flexibility inherent in the CPTF contract was temporarily ignored in the rush to the incentive contract.

A revision to the Armed Services Procurement Regulations in 1962 stated the preferential order of contract types to be FFP, FPI, CPIF, and CPTF. [11] Table 1 indicates the resulting shift to incentives. This has been explained as an attempt to provide the motivation normally engendered by a commercial environment, in which competition rather than a cost outcome determines prices, and competition rather than cost analyses insures that profit remains fair. As the complexity and expense of weapons and space systems have increased, their useful life has decreased. There is now little opportunity to retrofit or modify an operational system. Consequently, early design attention in the areas of maintainability, reliability, and standardization can achieve large "downstream" savings. Systems are also now procured in smaller quantities, and development funds constitute a larger fraction of total lifetime costs. Thus while superior operational performance is desirable, it must be balanced against development cost.

\[3\] Under conditions of price competition the government may award FFP contracts to the lowest bidder without regard for profit. [12]
The use of cost incentives was relatively straightforward, but multiple incentives raised a host of questions regarding the definition and measurement of performance elements, the determination of the relative value of alternative outcomes, and the formulation of fee patterns to motivate contractor decisions consistent with government objectives. In the past year techniques have been proposed that lead to fee structures more compatible with estimated government values. [15, 16] However, an inherent weakness of contractual incentives is that the ranges and weights assigned to performance variables, costs, and profits are negotiated early in the development program, while the contractor capability and potential operational value vary throughout the development effort.

Cost-Plus-Award-Fee contracts have obvious merit in advanced development and have been used most extensively by NASA and the Navy. They retain the flexibility of CPFF contracts while providing even more

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One of the first multiple incentive contracts went to TRW Space Technology Laboratories. Ten nuclear detection satellites were to be constructed at a basic cost of $14 million, and a basic fee of $1 million. A successful launch required that two satellites be orbited simultaneously. Incentives were paid depending on the number of attempts required to achieve a successful launch.

<table>
<thead>
<tr>
<th>First Success at Launch Number</th>
<th>Award Fee</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>$125,000</td>
</tr>
<tr>
<td>2</td>
<td>$98,000</td>
</tr>
<tr>
<td>3</td>
<td>$0</td>
</tr>
<tr>
<td>4</td>
<td>-$98,000</td>
</tr>
<tr>
<td>5</td>
<td>-$125,000</td>
</tr>
</tbody>
</table>

The expected lifetime in orbit was 2-4 months. For every day short of two months in orbit STL was penalized $1600, while for every day in orbit from four to six months, STL received an additional $1600. It would be difficult to objectively explain why these fees were chosen, except to note that the $250,000 launch fee swing is 25% of the basic fee, and that the maximum fee based on operational lifetime is approximately $100,000 per launch, or 10% of the basic fee. [13]
incentive than the possibility of increased profit. The very fact that profit is based on evaluation of contractor performance causes the contractor to perceive a relation between the current and potential future contracts. This appeal to the incentive for organizational survival may far outweigh any incremental profit.

An additional consideration in attempting to increase the effectiveness of procurement policy is the increasing scope of mathematical programming models. Zschau [17] has demonstrated that under reasonable conditions a large development project can be decomposed into sub-projects, and this decomposition used to obtain minimum cost surfaces for every feasible schedule and performance outcome of the total project. This suggests that a similar approach be used to periodically evaluate the worth of alternative system performance as development progresses. Before developing this idea further we will discuss the behavioral assumptions and information required to use contractual incentives.

**Contractor Motivation**

Contractual incentives are based on the assumptions that the contractor can exert a known degree of control over the product or system, and that he will act to improve his perceived corporate position. Economists completely characterize contractor behavior by stating that program decisions are made to maximize utility, but this is not an operationally useful statement in structuring an incentive fee.

Classical models of the firm assume that production and pricing decisions are made to maximize some form of monetary profit, which may be current profit, discounted future profit, or expected profit per period. In a cost incentive development situation the firm "produces"
final cost. The fee is usually a constant plus some fraction of the
difference between initial target cost and final cost. If final cost is
controllable by the contractor, and he maximizes current fee, he will
choose the minimum possible final cost, regardless of the sharing pro-
portion. The fact that cost targets are rarely underrun by large factors
suggests that cost reduction actions cause disutility. Scherer has
explained this with a user cost model which we describe briefly.
[3, p. 407]

User cost is defined to be the future profit loss resulting from
current cost reduction actions. A contractor may initially effect cost
reductions by increasing efficiency and reducing unnecessary waste, but
beyond some point further cost reductions can be achieved only by such
actions as laying off personnel, breaking up design teams, or curtailing
effort in development. The contractor perceives this as producing a less
attractive system in the long run, thereby adversely affecting his likeli-
hood of receiving future contracts. Assume that the user cost can be
represented as a quadratic function of the cost underrun.

Let \( C_t \) = Target Cost

\( C_a \) = Actual Cost Outcome

\( X = C_t - C_a \), Amount of Cost Underrun.

Then \( U(X) = a + bX + cX^2 \), and we see that depending on the coefficients,
a, b, c, cost underruns may lead to high user costs. Assume that the
contractor maximizes the difference between accounting fee and user cost
through his choice of \( X \).
Let $\pi_a = $ Accounting Fee

$\pi_t = $ Target Fee

$\alpha = $ Sharing Proportion

Then $\pi_a = \pi_t + \alpha X$, and the condition that $X^*$ maximize $\pi_a - U(X)$ is

$$\frac{d}{d\alpha} [\pi_a - U(X)] = 0, \text{ or } \frac{\alpha - b}{2c} = \frac{\alpha}{2c}.$$ Given $\alpha$, and the user cost function, an underrun will not always be optimal. The contractor's problem in negotiation is to choose the $\alpha$ that maximizes the final net profit, $\pi_t + \alpha X^* - U(X^*)$. Scherer shows that except in special cases, the optimal $\alpha$ polarizes to 0 or 1, that is either CFFF or FFP coverage.

Since the average value of $\alpha$ observed in Scherer's case studies was 0.2, he tested the hypothesis that contractors maximize profit subject to a constraint on maximum profit. This behavior would be explained by the attempt to avoid unfavorable publicity and interest on the part of the Renegotiation Board and the General Accounting Office. The following case described by the president of Giannini Controls Corporation is an instance of the operation of these factors.

Chart No. 2 (not shown) shows the cumulative margin earned in one of our divisions on a series of contracts for heavy electronics equipment. We started at a heavy loss in 1960 because all startup costs and some developmental engineering applied against that first contract, but as production increased we showed a satisfactory trend. Obviously we made profit on every contract after the first one, since the cumulative margin rises, but that startup hole was pretty deep and we didn't break even until early 1962. During this time prices were reduced substantially, our quality record was excellent, and we were shipping equipment ahead of schedule. In every way we were rated an exemplary supplier. Looks pretty typical, doesn't it? We invested heavily in an attractive program, covered our first costs, brought the cumulative profits up to industry average, and satisfied our customer completely. And then we ran into the buzz saw. The prime agency audited a late 1962 contract that had somehow earned an unusual profit: over 20%. We had never made close to that on any prior or subsequent contract, and the order at issue was only $200,000 out of $4,000,000 total business. Our margin before that windfall was just over 4%, and with the
"excessive profits" included we barely reached 6.4%, but the contract stood alone and none of this background carried any weight. The result was exactly 11 months of transcontinental debate with the agency demanding a profit rebate. We wouldn't agree to refunding a prior year's earnings, but we did accept a downward redetermination of the current production price. That's why you see the margin slipping off a quarter point in 1964--we're paying a penalty now for conservative estimating on a single contract in 1962! [19]

If such occurrences were frequent the weight assigned to profit as a motivator of contractor behavior would be low. Peck concluded, however, that risk aversion rather than concern over renegotiation motivates the low sharing proportions observed, and Scherer also stated:

Much more frequently, contractor representatives mentioned risk aversion as the principal reason for their efforts to negotiate low values of \( \alpha \). They stated explicitly that in many instances, given the uncertainties which pervaded advanced weapons systems and sub-system production programs, they were willing and eager to sacrifice the higher average profit expectations associated with firm fixed price contractual coverage for the greater security against an occasional short run loss afforded by CPFF and incentive contracts. [18, p. 275]

Baumol suggests that many large commercial firms make pricing, advertising, and marketing decisions with the objective of maximizing total sales revenue, subject to a constraint on the minimum profit that will satisfy stockholders. [20] This hypothesis is only partially applicable to the large defense contractors, whose sales are measured by the dollar volume of research, development, and production contracts held. Too rapid an expansion of this type effort results in a dilution of product quality, leading to a poor corporate reputation. Consequently, defense contractors prefer to expand gradually, keeping key research and design teams intact and productively employed.

There are less quantifiable motivators pertinent in predicting corporate behavior. These include organizational survival, security of employment, sales, and profit, freedom from harassment, the desire for
public approbation, the desire to enhance national security, and the
desire to advance science and technology. [3, p. 7] The research oriented
contractor is especially interested in retaining his key scientists and
maintaining the continuity of his research effort. Thus while contractor
negotiators may act as if they intend to maximize current accounting
profit, there is evidence that program decisions are based on the longer
range factors. [3, p. 159] In some cases the advantages accruing from
pioneering in a new technical area may be sufficiently strong to motivate
a cost-plus-no-fee contract, and in all cases the profit motivation must
be related to both the project uncertainties and the intangible contractor
rewards. In the current atmosphere of limited and selective development,
sharing proportions are more often 0.4 or 0.5 than 0.2, and the revised
profit policy of DOD emphasizes profit as a stimulant to contractor
effort. Fewer projects are being carried into engineering development,
and projects in research, exploratory, and advanced development are
subject to immediate cancellation if more promising concepts appear.
[23, 24] Under these conditions concern for organizational survival is
very real, and there is less of a requirement for extensive contractual
safeguards and stimulants.

Scherer's user cost model is descriptive in explaining contractor
behavior in cost reimbursement and cost incentive contracts. We would
proceed one level deeper in an effort to understand the cost-time-
performance choices made by the contractor. We hypothesize that the
defense contractor acts to maximize the probability of program continu-
ation subject to a minimum profit constraint. A single very large loss,
or a succession of unprofitable projects will drive down the market
price of the firm's securities, adversely affect the firm's image among the agencies awarding contracts, and raise doubt as to the capability of the firm's management. Hence the program manager is constrained in the short run against outright loss, and in the long run to some minimum profit. During any given development contract, we feel that the program manager and other corporate personnel have a perception of the customer valuation of the different parameters of the system.\(^5\) Adjustments to these perceived values are obtained from the in-plant agency representatives, visiting service personnel, liaison personnel in Washington and other headquarters, and from knowledge of the progress of potentially competitive projects. The program manager can therefore judge the relative desirability of alternative schedule, performance, and cost outcomes from the buyer's point of view. His initial system choices are determined by the acceptable time-performance outcomes defined by the development contract, modified by the current view of the buying agency, and constrained by the cost-time-performance surface. If the cost resulting from a desirable outcome is less than target cost, the manager decides whether to take the cost underrun as a saving for the buyer, or

---

\(^5\)Our survey of aerospace contractors revealed one development project incorporating an incentive for the successful firing of a missile by a certain date. Failure to successfully fire by that target date would result in successively reduced incentive payments, until at the end of six months a substantial penalty would be levied against the contractor. The buying agency was aware that pressure existed to halt program funds and realized that an early successful firing was necessary. The contractor also recognized that failure to demonstrate the system successfully would increase the likelihood of program cancellation, and therefore put forth a "crash" effort. It is our opinion that the pressure for program continuation was more important than the contractual reward at stake.
to utilize it to increase performance in profitable areas. The latter decision enhances system attractiveness relative to potential competitors. If, however, the characteristics considered essential by the customer are only attainable at a cost above target, the program manager must decide whether performance or monetary considerations are more important in influencing program continuation.

At any program decision point the contractor behavior is then conceptually described by a nonlinear programming problem in which the firm chooses performance variables, schedule outcomes, and development cost to maximize the probability of program continuation subject to a constraint on minimum profit and contractual constraints on the values of performance, schedule, and development cost.

Let \( c \) = cost outcome aimed at by the contractor,
\[
t = \text{schedule outcome aimed at by the contractor,}
\]
\[
p = \text{vector of performance outcomes aimed at by the contractor,}
\]
\[
p = (p_1, \ldots, p_n),
\]
\[
c = \text{current cost ceiling desired by the buyer, as perceived by the contractor,}
\]
\[
t = \text{current development schedule desired by the buyer, as perceived by the contractor,}
\]
\[
p = \text{vector of performance outcomes desired by the buyer, as perceived by the contractor,}
\]
\[
p = (p_1, \ldots, p_n).
\]

Let \( \psi(c,p,t) = 0 \) represent the minimum cost attainable at every time and performance outcome, and let \( \varphi(c,p,t) \) denote the contractor fee at any outcome. Assume that the probability of program continuation approaches one as the projected project outcome approaches the outcome.
desired by the customer. Assume further that the program manager estimates the weights \( \beta_i \), and that minimizing

\[
\sum_{i=1}^{n} \beta_i (p_i - p_i) + \beta_{n+1} (c-c) + \beta_{n+2} (t-t)
\]

approximates maximizing the probability of program continuation. The contractor problem is then

\[
\text{Minimize} \quad \sum_{c,p,t} \sum_{i=1}^{n} \beta_i (p_i - p_i) + \beta_{n+1} (c-c) + \beta_{n+2} (t-t)
\]

such that

\[
\psi(c,p,t) \leq 0
\]

and

\[
\phi(c,p,t) - \psi(c,p,t) \geq 0
\]

In view of the information and decision variables entering the contractor decisions under this hypothesis, we conclude that expected fee alone is a weak and uncertain motivator. Even if the cost-time-performance tradeoff surface were stationary and available to the buying agency, and if a set of multiple incentives could be structured consistent with true government values at every outcome and assuring maximum contractor fee at that outcome desired by the government, there is no assurance that
this outcome would be chosen by the contractor. Contractual incentives can be analyzed for consistency, but not for motivation.\textsuperscript{6}

We conclude that in research and exploratory and advanced development, the primary contractor motivation is to maintain the firm's capability to compete for and retain projects. There is increasing realization that only a few systems will proceed into engineering development, and fewer still into production. A greater amount of effort will be directed into research and exploratory development, and profit will be based on the cost of the research effort as well as results. These conditions create incentives for efficiency, and the marginal value yielded by the superposition of contractual incentives may not be worth the effort required of the contractor and procurement personnel.

**Government Objectives in Development**

Regardless of the source of the requirement for a new weapons or space system, a preliminary study of projected system utilization is necessary to specify initial performance characteristics and operating conditions. In structuring an incentive contract the government must have a ranking of the possible development outcomes, and even in a CPTF contract where decisions on specific details may be made later in the program, recommendations should be based on some consistent evaluation of

\textsuperscript{6} This model of contractor behavior is consistent with Peck's observation that firms are more prone to take risks if their projects depend on it. [7, p. 540] For example, if the purchasing agency is debating program continuation and has suspended funds, the contractor may supply corporation funds to continue critical development of production until a final decision is reached. Not to do so might affect the final decision unfavorably, while in the event of program termination some of the expended funds may be recovered.
system worth. Yet this task of determining operational value is extremely
difficult and controversial. The Holifield Report criticizing Defense
management of satellite communications is quoted as stating:

"The unfortunate experience (of DOD analysis of alternative satellite
communications systems) suggests to us," the report continues, "that
considerations of economy from a Defense Department budget stand-
point are not sufficient criteria for making decisions in fields
which involve government-wide policy, politics, and international
diplomacy." [25]

Sidestepping the question of whether a particular system should have
been brought to the development stage, the buyer may select various
objectives in the development effort. The more obvious are

(i) maximizing some measure of effectiveness subject to a cost
constraint, and

(ii) minimizing the lifetime cost of the system subject to the
attainment of some fixed level of effectiveness.

There may be other valid objectives such as minimizing the number of
military personnel required to achieve a given level of effectiveness,
or obtaining a system flexible enough to operate in a number of possible
environments. In any event development objectives and government negoti-
ation objectives are not identical. The former are determined by evalua-
tion boards, cost-effectiveness analysts, and military personnel, while
the contract negotiations are performed by procurement specialists beset
by a welter of regulations, guidelines, and hindsight. Hence the Director
for Procurement Policy, OSD, testified:

A requirement (for a weapons system) having been established, and
feasibility determined, what is our objective? It is to procure
the development, production, and delivery of the weapons system
into inventory in the shortest possible time and at the minimum
cost. [27, p. 106]
But Scherer's case studies indicated that actual government procurement practice has been a compromise between maximizing the incentive for contractor efficiency, minimizing contract outlay, and minimizing the risk that unnecessary and excessive profits will be paid.\textsuperscript{7} [3, p. 147]

There is an effort in current procurement policy to consider total system costs and not be hesitant about spending funds in development to achieve operational savings, but at the initiation of development the potential value of a new system is extremely uncertain. The cost-effectiveness analyst assumes that development will be successful and tries to foresee the implications of the successful system in the relevant time frame. The essentiality of specific characteristics will vary with developments in associated systems, in opposing systems, and in the operational environment.\textsuperscript{8} Preliminary models of the system attempt to determine the key parameters, and the variation in performance as a function of these parameters. After preliminary studies have shown a concept to have merit, the Project Definition Phase (PDP) may begin. PDP is a formal step in which the cost-effectiveness model is refined.

\textsuperscript{7}Minimizing contract outlay is especially shortsighted in view of the operating and maintenance expenses of modern weapon systems. In an expository example Hitch and McKeen point out that a hypothetical air transport fleet obtained under the minimum contract outlay would cost over $1 billion more than the least cost fleet on a lifetime basis. [26, p. 140]

\textsuperscript{8}An appropriate approach in this situation would seem to be a sequential decision process, with specification details supplied to the contractor as late as possible to take advantage of available information. According to one theory Defense programs may not be managed in this manner because the members of development staffs recognize the transitory nature of their assignments and hesitate to leave vital decisions to their unknown successors. Hence the system is "cast in concrete" at an early point. [28]
further, and the total system costs and development schedule estimates are determined. It would be desirable to have government personnel determine the relative value of the possible outcomes, but normally two or more contractors will conduct PDP, competing for the follow-on work. During PDP these contractors will determine the "optimum balance between total cost, schedule and operational effectiveness for the system." [29] Following PDP a development contract may be awarded on an FFP or FPI basis, using the tradeoffs generated in PDP as a basis for the incentive structure.

Consistent incentives require the capability to compare performance increments to cost increments. Therefore the analyst must find some method of ultimately comparing different aspects of performance in dollar terms. This may be straightforward in some cases. For example the value of early delivery of a military satellite communications system might be indexed by the current dollar usage rate of commercial facilities. The problem is more usually akin to determining the worth of additional "rate of climb" in an anti-submarine aircraft. Quantitative statements about the worth of accelerated development undoubtedly include consideration of possible system obsolescence due to oncoming competitive systems, or the need of the agency to demonstrate a successful prototype to obtain further development funds. Development programs are basically either urgent, and rapid development receives priority, or else the purchasing agency is interested in showing cost efficiencies, and only enough is spent to keep the schedule from making uneconomic use of resources. Scherer concludes that the fee allotted to schedule incentive would be better saved for cost reduction incentive, and in fact OSD procurement
personnel state that the early emphasis on schedule incentives appears to be diminishing. [14]

Several approaches have been used in quantifying the worth of additional reliability and maintainability. Van Tijn developed a model which measures variation in support costs as a function of subsystem reliability and maintainability. [31] The model accumulates support costs based on the maintenance effort predicted by reliability models of the system. Douglas Missiles and Space Systems Division performed a reliability sensitivity analysis prior to suggesting a reliability incentive for a missile system. [32] Data derived from their study is shown in Table 2.

<table>
<thead>
<tr>
<th>Missile Reliability</th>
<th>.84</th>
<th>.97</th>
<th>.99</th>
<th>.999</th>
</tr>
</thead>
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<tr>
<td>Operational Reliability</td>
<td>.67</td>
<td>.78</td>
<td>.80</td>
<td>.80</td>
</tr>
<tr>
<td>Cost/Missile ($10^3$)</td>
<td>230</td>
<td>275</td>
<td>311</td>
<td>355</td>
</tr>
<tr>
<td>Number Required</td>
<td>148</td>
<td>128</td>
<td>125</td>
<td>125</td>
</tr>
<tr>
<td>Total Missile Cost ($10^6$)</td>
<td>35.2</td>
<td>35.2</td>
<td>39.0</td>
<td>44.3</td>
</tr>
<tr>
<td>Total Operational Cost ($10^6$)</td>
<td>37.0</td>
<td>35.2</td>
<td>38.8</td>
<td>44.1</td>
</tr>
</tbody>
</table>

This allowed Douglas to determine the optimal reliability around which the incentive would apply. A report by the Management Systems Corporation contains an example in which incentives were desired for both accuracy and reliability in a tactical ground-to-ground missile. [16, p. 22] In this case the overall effectiveness criterion was the ability to inflict a specified amount of damage on a given target with a specified level of confidence. Thus the number of missiles procured would depend on the number required to achieve the fixed level of damage,
and accuracy and reliability could be directly related to the damage expected per missile launched. Variations in accuracy and reliability could then be evaluated by the cost savings in missile procurement, and the incentive fee related to the cost saving incurred. When it is not possible to relate the performance variation directly to a cost saving, the Management Systems Corporation study advises:

... it is necessary to estimate the relative dollar worth of effectiveness over the range of project outcomes. This judgement might be expressed as follows:

\[
\Delta \text{Value} (\$) \quad \Delta \text{Effectiveness}
\]

The dollar value of effectiveness defines a tradeoff between total cost and effectiveness which the government wishes the contractor to use when making decisions which could affect both quantities. [16, p. 77]

Even comprehensive analysis prior to development contract negotiation leaves many unanswerable questions. In the next section we will consider how the ability to decompose a large project into smaller activities can aid in obtaining an accurate cost-time-performance tradeoff surface, but this capability does not aid appreciably in resolving the basic uncertainties that exist in estimating system effectiveness. To decide between all potentially important alternatives at the very initiation of a program is at best a guided guess, and at worst the cause of an inadequate system or costly development redirection.

Cost, Time, Performance Tradeoffs

The Incentive Contract Guide confidently states:

... one method for making this determination (of tradeoff functions) is to request the contractor to provide forecast performances for several combinations of cost and schedule, (with analysis by the government as to reasonableness). This of course will be standard practice under Project Definition Phase. [4, Ch. VIII, p. 7]
The difficulty in obtaining such a performance forecast is that there are many equal cost methods of achieving very different final sets of specifications. A specified missile range may be attainable with alternative thrust capability, fuel capacity, and vehicle weight. Given only the cost and schedule the contractor usually will have wide latitude in the systems he can produce. However, the derivation of a realistic cost-time-performance surface is desirable in any development project. It is especially necessary to formulate a consistent incentive contract. And it may make possible submission of several FFP development bids with different specifications by the same contractor.

There are quantitative tools available and under development which yield estimates of the cost-time and cost-time-performance surfaces. A PERT network, for example, is established by defining a set of activities, an order or precedence relation, and a spread of activity completion times for some fixed funding. Methods have been available for some time to derive the expected completion time for the total project. The distribution of project completion time can be computed, analytically if the activity completion times are exponentially distributed, and by simulation in the case of more complex distributions. [30] If cost is a convex function of completion time for every activity, the minimum cost required to complete the project in a specified time may be derived.9

9Roderick W. Clark points out that time-cost data is not readily available because management accounting systems are not structured to provide ex ante estimates of costs as functions of completion times for various activities. This is due to the fact that accounting systems are designed to facilitate rapid reimbursement for incurred costs with only secondary emphasis on internal control. [34]
Thus, theoretically, the PERT network could be reformulated for several representative sets of performance specifications, and the general nature of the tradeoff surface approximated. However, even a system with a relatively small number of performance characteristics may have thousands of separate activities. Each time the key variables are assigned new values, the contractor must find the subsystems and sub-sub-systems that achieve the required performance levels with minimum cost for some fixed schedule. The requirement to reformulate in detail the entire PERT network would require an inordinate amount of administrative time, and would interfere with the normal work of the development personnel.

There are thus two separate problems. First the minimum cost subsystems that achieve desired performance in fixed time must be determined. Then the minimum cost schedule to develop specified subsystems must be found. Edward Zschau has formulated a procedure called "Project Modelling" which considers technical interrelationships among design specifications, precedence relationships among development activities, and determines the minimum cost of developing within an allotted time a system with specified performance characteristics. [17] The model assumes that cost in each activity can be expressed as a convex function of improving performance or decreasing development time, and that the total project cost can be expressed as the sum of the activity costs. Rather than attempting to solve the entire problem simultaneously, the procedure uses the tradeoff surfaces generated at activity levels as inputs to higher level optimization models and iteratively determines the minimum cost activity specifications for fixed time. For this set of activity specifications the minimum cost schedule is derived, and the specification
model is then re-entered with the new schedule. Zschau has shown that this iterative procedure converges to a solution optimal in the following sense.

When the project demands (e.g., the desired project duration and system performance characteristics) together with the activity cost functions are fed into the project model, it outputs the minimum project cost, the optimal engineering specifications, and the optimal project schedule for these demands and cost functions. By comparing the optimal project costs with the alternative sets of project demands that give rise to them, the optimal tradeoff functions relating time, cost, and performance can be generated. [17, p. 23]

Decentralized project management offers further advantages. Less information flow is required, decisions are made closer to sources of data, and it may be possible for the buying agency to decompose a large development project into subprojects for several contractors. This would allow each contractor to compete for each subsystem, and yet allow the buyer to insure the proper meshing between subsystems. Furthermore the contractor could furnish tradeoff surfaces without divulging either activity costs or proprietary process information.

The feasibility of project modelling depends on the accuracy with which development personnel will predict costs. Some incentive for accurate estimation rather than protective estimation may be required. The project modelling approach should be tested in a tractable development program to determine its further applicability.

One additional consideration in basing decisions on cost-time-performance surfaces is the uncertain nature of contractor control. For example, top level program management may intend that a limited number of development hours be devoted to reliability considerations, but engineers at the working level may be motivated by professional or personal interest to spend more than the allotted number of hours on this area.
The time may be charged to an appropriate task, but the individual may still be thinking about ways to improve reliability. Overall project tradeoffs should be viewed as estimates to be systematically improved.

**Profit**

If system value and development cost were known as functions of availability date and performance characteristics, a consistent fee function could be defined by applying a profit rate to the difference between value and cost at any outcome. We have indicated that this would in general not motivate the contractor to aim at the outcome which maximizes monetary profit. However, the actual profit attainable at any outcome can be a strong determinant of contractor behavior if profit rapidly decreases as characteristics drop below those desired by the buyer. Our hypotheses about contractor behavior imply that contractors will be motivated away from areas of potential monetary loss, but are relatively indifferent between small positive profit increments. In August, 1963 the Weighted Guidelines revision to ASPR 3-808 was issued to aid contracting officers in determining appropriate profit rates. [11] Prior to this, contracting personnel were provided with a wealth of suggested factors and considerations, but were required ultimately to base their profit proposals on experience and judgement. In a report on actual contracting practices, Summer Marcus states:

Many contracting officers choose the expedient solution to their quandry (of conflicting and diverse decision elements). Through experience they arrive at a profit or fee rate that is well below the maximum permitted, but high enough so that the contractor will accept it, and they use these few rates over a long period for all contracts they negotiate, regardless of contractor situation. As time goes on they tend to lower the rate slightly to establish themselves as good bargainers. [35]
The Weighted Guidelines procedure establishes categories such as Contractor's Input to Total Performance, Contract Cost Risk, Past Performance, and Special Profit Considerations. Contracting officers use available information to assign allowable weights to appropriate categories, and then rely on judgement in assigning weights to risk and special factors.

Based on interviews conducted prior to September, 1964, the most frequent criticism of Weighted Guidelines was that most government activities used the technique to establish a rationale for the final negotiated profit rate, rather than establishing an initial profit objective. Table 3 shows eleven FFP procurement actions negotiated at one buying activity using weighted guidelines.

<table>
<thead>
<tr>
<th>Weighted Guideline Objective</th>
<th>Profit Proposed</th>
<th>Historical Rate</th>
<th>Profit Negotiated</th>
</tr>
</thead>
<tbody>
<tr>
<td>12.37</td>
<td>11.7</td>
<td>10.5</td>
<td>11.7</td>
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<td>11.45</td>
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<td>11.11</td>
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<td>11.87</td>
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<td>6.0</td>
<td>6.0</td>
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<td>11.98</td>
<td>11.11</td>
<td>11.1</td>
<td>11.11</td>
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<td>14.0</td>
<td>11.1</td>
<td>11.1</td>
<td>11.1</td>
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<td>12.02</td>
<td>6.0</td>
<td>10.0</td>
<td>6.0</td>
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<tr>
<td>12.1</td>
<td>10.0</td>
<td>10.0</td>
<td>10.0</td>
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<td>14.0</td>
<td>11.1</td>
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<tr>
<td>15.0</td>
<td>21.0</td>
<td>---</td>
<td>15.0</td>
</tr>
<tr>
<td>13.2</td>
<td>11.0</td>
<td>11.0</td>
<td>11.0</td>
</tr>
</tbody>
</table>

10 Collection and retention of contractor performance data is formally required by the Contractor Performance Evaluation program. Contractor success in meeting cost, performance, and schedule requirements will be recorded, and source selection boards will be required to review this information prior to making contract awards. [36]
This particular set of data confirms that contracting officers use high initial profit objectives to avoid later explanation of failure to negotiate at a point below or near a valid initial objective.

The contractor cannot press for too high a profit rate or target cost for fear of later unfavorable attention or the label of "profiteering." One government procurement officer discussed the environment in which the contractor's negotiator must operate.

It is quite evident that the lead negotiator's performance is staged in a "fishbowl" visible alike to corporate and division, but the end product of his negotiation is also subject to much scrutiny. In a government negotiation he can't afford to extract a profit too high even if the opportunity arises. Two of the interviewed companies recently had instances where they felt the profit was too high and the contract was not executed. In the only corporation interviewed that hired negotiators from outside the Government/Aerospace complex, it was stated that the most difficult point to instill in their thinking was not to always seek the maximum profit as is customary in normal commercial bartering. [37]

When target costs are negotiated a good deal higher than projected cost outcomes, the contractor rarely risks taking a windfall profit, but instead voluntarily rebates profit, or commits additional funds for performance improvement, schedule acceleration, or related research. This behavior is explained not only by the existence of the GAO and Renegotiation Board, but by the low incremental cost of supporting additional research in this way. 11

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11 The Renegotiation Board annually reviews profits of contractors with more than $1 million in sales in order to recover excessive profits. The Board attempts to take the individual contract types into consideration, but it is difficult to determine those contracts in which profits are due to contractor efficiency as opposed to skill during negotiation, and the board's judgement has been questioned. [3, p. 261]
Recent DOD profit policy has aimed at negotiating higher profit rates in return for the contractor assuming more of the cost risk. The number of cost incentives negotiated in the 50/50 to 70/30 range has increased substantially, and there are currently very few low risk sharing arrangements concluded. [4, 6] This indicates the awareness at procurement policy levels that a number of factors motivate the contractor, and that by using higher profit at the desired outcome, and lower profit at less desirable outcomes, the government can motivate contractor behavior that is more consistent with government values.

**Formulation of Contractual Incentives**

A well structured incentive arrangement has two objectives. It will create a financial motive for the contractor to achieve superior performance in all variables. If this is impossible, the fee attainable at every possible outcome will reflect the government's ranking of that outcome, thereby guiding the contractor's choices.

The initial use of incentives was restricted to cost incentives, and the fees were almost always linear in the development cost. Justification for linear fees in such contracts is their simplicity of negotiation and ease of interpretation. The previously referred to Management Systems Corporation study justifies linear fees in terms of their flexibility. A contractor can always be motivated to aim at a certain outcome (if he is a fee maximizer) by means of a linear fee. But the example they present is misleading.

Let $C_D = \text{Development Cost}$,

$C_o = \text{Operating cost during lifetime},$

$C_T = \text{Total Cost} = C_D + C_o$.
Assume that $C_0$ is dependent on $C_D$ in the following way:

$$\mu_t[C_0 = x | C_D] = C_D^{-\delta}$$

Then $E[C_0] = 1/C_D$, where $E$ denotes the expectation operator. Hence $E[C_T] = C_D + 1/C_D$. Now assume that the buying agency desires to find that $C_D$ which minimizes $C_T$. Then $\frac{\partial C_T}{\partial C_D} = 0$, which implies that $C_D = 1$.

Let $f$ denote the contractor fee as a function of development cost, and assume that the contractor maximizes expected fee. In order to motivate the contractor to set the development cost, $C_D$, at 1, we must have $\frac{\partial f}{\partial C_D} = 0$ for $C_D = 1$. If the fee is linear in the total cost, $C_T$, we have $f = \alpha - \beta C_T$, for some $\alpha$, $\beta$, and $E[f] = \alpha - \beta E[C_T]$ or $E[f] = \alpha - \beta [C_D + 1/C_D]$. This is maximized at $C_D = 1$ for any choice of $\alpha$, $\beta$. If however, the fee is linear in the development cost, $C_D'$ (e.g., $f = \alpha - \beta C_D'$), then the development cost chosen by the contractor is indeterminate without further assumptions.\(^\text{12}\) Using linear fees, we must be careful about the variable the fee is based on, for the fee in the above example is certainly nonlinear in development cost.

When CPFF contracts were extended to include multiple incentives, the fees were generally independent, linear, and additive. For example, development cost might be weighted 40%, schedule 20%, and performance 40%. Individual performance elements such as thrust, or range, were then assigned portions of the 40% performance weight. This approach had

\(^{12}\)Recall that we were required to postulate the existence of a user cost function or disutility resulting from cost reduction to explain the contractor's behavior in this case.

28
several defects. It required the subjective ranking of the separate performance, schedule, and cost elements. Any set of weights valid in expressing government value at one point were not in general representative of value over the possible performance spectrum of the system. And in such a fee arrangement it was possible for the contractor to earn target fee while submitting poor performance in several variables. Consider the following example which illustrates this fault of compartmentalized incentives.
If the contractor delivered an 1150-mph aircraft three months after target at a cost of $130 million he would receive target fee, and if he delivered the same high performance aircraft six months late at a cost of $160 million he would still earn target fee. [2, p. 43]

The linearity of the fee structure oversimplifies the guidance to the contractor. When the performance level of some element is close to minimum acceptable, the value of an increment in that variable's performance is worth much more than an increment in performance in some variable already close to the maximum desired. Linear incentives do not reflect this. One partial solution to compartmentalization is to extend the range of the cost incentive to retain some control over cost overrun. This is termed "overlap." And the partial solution to the independence of the fees is to define the acceptable outcomes in terms of several variables. For example, acceptable minimum range might be 4000 miles, and acceptable minimum speed 1600-mph, but no more than minimum fee would be paid unless range plus speed exceeded 6000. This is simply taking care to bound the region within which a linear fee will be acceptable to the government.
Two methods have been proposed to define a fee surface which more accurately reflects government values than do linear fees. The Logistic Management Institute proposed a "Tabular Model" which uses a set of multipliers to alter the fee structure in regions of contractor performance defined as "outstanding" or "poor." [15] After the government and contractor negotiate the regions deserving special attention, fees are structured independently for each variable, and total fee is the sum of the individual fees. The fee for the outcomes in the specially designated regions are multiplied by suitable constants, greater than one for superior performance, and less than one for below standard performance. This corresponds to modifying a planar fee surface by steeper planes in these designated regions as shown in Figure 1.

![Figure 1](image)

The name "Tabular" arises since the buying agency must define grades of performance (tables) for the designated variables, and then must determine appropriate multipliers for each grade. For example, assume the basic fee for control of development cost is linear, and the buyer desires
to reward exceptionally low development cost and penalize a combination of high development cost and low performance. The buyer therefore might divide the range of expected development cost outcomes into ten grades. The lowest two of these grades could be assigned the multipliers 1.4, and 1.2, while the combination of high development cost and low performance could be assigned the multiplier 0.5. The use of multipliers is an explicit method of requiring the procuring agency to determine its own utility function, and then introduce any significant nonlinearities into the incentive fee. Using multipliers for combinations of variables requires a method of ranking the relative desirability of outcomes. Without a model of system effectiveness, determining the relative value of the outcome of three parameters, each having ten grades, could be tedious.

We have seen that a basic problem in incentive fee formulation is that the effectiveness of a system is not additive, but is a complex function of the variables characterizing the system. The Management Systems Corporation study recognizes this and suggests that effectiveness be measured by a cost-effectiveness model, and that the incentive fee be a monotonically increasing function of government value. The fee structure they propose is:

\[ f = f_t + \alpha(\Delta V - \Delta C) \]

where \( f \) = final fee,
\( f_t \) = target or initial fee,
\( \alpha \) = sharing proportion,
\[ \Delta W = \text{incremental increase in effectiveness value over target effectiveness value,} \]

\[ \Delta C = \text{incremental increase in cost over target cost.} \]

We discussed the problem of translating system effectiveness into dollar terms. Even when it is difficult to make such comparisons they will be implicitly used in any case, and the use of cost-effectiveness techniques requires that they be made formally. One of the criticisms heard against government contract management is that the buying agency is vague about its requirements and choices among performance capabilities. This objection could be met by furnishing the contractor with the cost-effectiveness model, or data derived from it. The contractor would then have a more objective basis for tradeoff decisions, and the government would have a more objective method of deciding whether the contractor did, in fact, act in the government's best interest.

A combination of Zschau's approach to decentralized project management, the periodic updating and use of system effectiveness models, and use of an award fee appear to offer substantial efficiencies in development contracting. An initial cost-effectiveness model may allow an incentive structure that reflects government desires. But as development progresses and both the cost-time-performance surface and the external environment change, the most desirable system will also change. By using project models and system models capable of being updated, the contractor and buying agency could respond to the operating environment adaptively. If tradeoffs are to be evaluated in detail, decentralization would be as desirable in the user's cost-effectiveness model as in the development project model. Decentralization allows more detailed and accurate
estimates of effectiveness value but requires a higher level model to avoid suboptimization. Use of an award fee is a step closer to the "After-The-Fact-Evaluation" contracting method proposed by Scherer.\textsuperscript{13} The award fee allows the buying agency to take account of outstanding innovation or management, and to compensate for unforeseen requirement changes or technical problems that may have precluded the contractor from earning a fair profit in an otherwise satisfactory program.

Implementation of such changes would have to be gradual. The defense and aerospace contractors are just becoming familiar with the multiple incentive contracts initiated in 1962. After an initially lukewarm reception, industry reaction now appears favorable. [13] A Douglas Missiles and Space System Division report states:

An incentive program usually stimulates the development of efficient techniques resulting in higher reliability and more contract dollars. [32]

Contracting officials at two aerospace corporations commented that the magnitude of the incentive fees did not seem to be very important. The knowledge that the project was "on incentive" was sufficient to induce higher performance and better planning. Some corporate executives felt that performance incentives were superfluous. Other writers argue that reliability and maintainability incentives are required with cost incentives to prevent corner cutting in quality. W. C. Frederick states:

\textsuperscript{13} Under Scherer's proposal the government would maintain data describing contractor efficiency in cost control, reliability, operational performance, and adherence to schedule, which would be used in determining profit rates, and awarding contracts. In addition, Scherer proposes that the growth, long term profits, and direction of activity of the contractor be planned by a Performance Evaluation Board. This would place emphasis on the long range incentives of survival and growth. [3, p. 327]
The risk to the purchaser (in a firm-fixed-price contract) is a bit more subtle, but no less real. Under the FFP type of contract the contractor's profit varies inversely with the cost of the system; therefore he has a tremendous incentive to cut costs, perhaps to the point where the quality of the system is degraded. Thus unless the specification or work statement is extraordinarily precise and complete, the purchaser may be forced to accept a product of a quality lower than he had expected. [39, p. 7]

And again in arguing against single dimensional (cost) incentive contracts,

The provisions of one incentive contract of which the author is aware were such that the contractor's fee would be greatest if his cost were exactly $1,000,000. It was the stated intent of the program manager to accumulate costs of $1,000,000, and since costs were in danger of exceeding that amount, the planned reliability program was curtailed accordingly--despite the fact that the reliability objective was conceded to be extremely important to the customer and difficult to meet. [39, p. 8]

Scherer, however, repeatedly claims that weapons system contractors have an extremely high aversion to cutting quality or reliability to gain short run profit. The incentives toward high quality--corporate reputation, the reluctance of the purchasing agency to authorize design changes that impair quality, and the possible effect of low quality on future awards--do not appear applicable to commercial firms, while Frederick's comments, directed at defense contractors, do appear to be more applicable to non-defense firms with several markets. The utility of current profit is higher for these firms, and the renegotiation constraint is not present.

Government procurement personnel are pleased with the results of contractual incentives. The Incentive Contract Guide states:

The success of the program (expanded use of performance incentives) seems clearly revealed in performance levels achieved at target or high than target in the majority of cases and cost being controlled in almost every instance within 10% of negotiated target costs. [4, Ch. VIII, p. 2]
It is not clear whether these results have been obtained because of the existence of incentive fees, or because of the detailed system analysis and capability estimation that precedes contract negotiation. A set of experiments in which contractors submit proposals for FFP, CPlF, and CPAF awards would illuminate the preferences for risk and contract type within individual firms, and between firms.

Summary

We have indicated the important trends in contracting for the development and procurement of complex weapons and space systems. The first is the greater selectivity and limited production instituted by Secretary of Defense McNamara, which creates a competitive atmosphere among contractors, and places the government in a more powerful bargaining position. Thus the government is able to institute management controls and practices which it considers conducive to the more efficient use of procurement resources and development personnel. The second operative factor is the growing bank of data recording contractor performance. Over a period of time chance performance variations will be averaged out, and the long term management competence which a given contractor brings to his projects will be quantitatively demonstrable in terms of development costs versus estimates, and field performance versus projected performance. The final factor is the growing arsenal of management techniques and tools for control and analysis of alternatives in project definition and development. Adaptive project modelling by the contractor and a similar government effort could operate together to allow sequential decision making during the development process. This
would be more efficient than fixing the ranges of system performance far in advance of any firm information regarding desirability, feasibility, and cost.

These three factors interact to allow government initiation of the project modelling technique, by creating general system models and project models of various space and weapons systems. Contractor performance data would provide initial model data which could be continually updated by completed projects. Similarly the system models could be refined by operational testing. As government and industry gain familiarity with the formulation, use, and properties of system models, the government could ask for development program bids which would have tradeoff decisions made by the government, sequentially, with the profit award being negotiated after project completion with final cost, performance, risks, and special factors available for consideration. We feel that this approach to program management promises a more responsive contractor effort in research and development than do the more rigid contractual incentives.
CHAPTER II
UNCONSTRAINED CONTINGENT PRICING MODELS

Introduction

In reviewing the use of incentives in development we observed that performance incentives were intended to motivate superior technological progress in the design and development of new products. Once high performance features are designed into equipment, production inefficiencies can degrade operation, but even outstanding production practices cannot improve uninspired design. We also noted that the difficult problems in designing contractual incentives were encountered in the advanced system area. It is useful to point out that relatively few contractors are involved in large advanced system development. The majority of producers supply standard or slightly modified equipment, and generally do not sell on an incentive basis. Government practice is to procure noncomplex products (shoes, medical supplies, hand tools, ammunition) in large lots and to specify the Acceptable Quality Level (AQL), or percent defective allowable in the lot.¹ The acceptability of complex items such as communication equipment or vehicles is determined by defects-per-unit,

¹Sampling of incoming lots is conducted in accordance with approved sampling procedures. A fixed price is paid for accepted lots. Items which cannot be designated defective or nondefective (such as capacitors with a range of capacitance, or ammunition propellants with variable muzzle energy) are classified by conformance to specification, and lot disposition is determined by the percent of the lot estimated to fall outside tolerance limits. [41]
demerits-per-unit, or reliability tests such as the AGREE procedures.

[22]

In production, as opposed to development, there is less cost uncertainty and little cause for changes in product specifications. Conformance to these specifications can be more accurately determined than in development, and it appears that fixed-price-quality-incentives, or "contingent pricing contracts" are applicable to these production situations. A contingent pricing contract would provide for variable payment based on the degree of product conformance to design specification as measured by an agreed on testing or sampling method. Such adherence to specifications is commonly termed "quality" and contingent pricing is therefore more of a quality related concept than are performance incentives in general. Note that in discussing performance incentives we implicitly assumed that at the completion of product development the performance parameters were precisely known. In fact these parameters (MTBF, range) can only be estimated or predicted from the observation of other variables. Hence there is sampling uncertainty in the payment received by the contractor or producer. Contingent pricing policies are designed to handle this sampling uncertainty directly. The concepts involved in analyzing the quality situations are also applicable to structuring performance incentives based on measured performance variables.

When the buying agency samples an incoming lot or tests an incoming equipment and finds that it does not conform to specification, it may either waive the nonconformance or reject the lot or item. When items cannot be rejected due to delivery or urgency considerations a single
unit price is not an effective tool to motivate superior production performance. The item marginally acceptable may be improved later by supplier effort, but the extra time and possible government expense involved should detract from the unit price in some way. Contingent pricing provides a method of rewarding contractor excellence in production and penalizing poor contractor quality control.

There have only been three models that explain contractor response to performance incentives, and of these only two attempt to design incentive policies optimal in some sense. In this chapter we briefly discuss these models, state the assumptions in them which we feel prevent the derivation of realistic policies, and then in the next chapter develop a method for generating more realistic incentive pricing policies.

Hill Model

The forerunner of the contingent pricing models is the effort by Hill [42] in 1960 to explain the incentive effect of both payment and acceptance sampling plans. Hill suggested that the primary purpose of acceptance sampling was to motivate higher production quality rather than simply detect and reject deficient batches. He assumed as axioms:

(i) the distribution of the batch percent defective submitted to the purchaser is entirely dependent upon the actions of the manufacturer; and

(ii) the manufacturer will take those actions which maximize expected profit.

From these axioms we see that if the lot quality (percent defective) submitted is to improve, the manufacturer must take action, and that he
will do so only if it is economically advantageous to him. A further implication is that it should be to the manufacturer's economic advantage to offer good quality and to his disadvantage to offer poor quality. The sampling and payment schedules to which Hill restricts his discussion provide for a fixed payment if the lot is accepted and no payment otherwise. In situations of this type, higher quality should be rewarded by greater frequency of lot acceptance. While all operating characteristics produce this result, it is standard practice to shift sampling plans as the sampling results provide information on the average percent defective submitted in incoming lots. Consider the following case which illustrates how certain features of the Dodge-Romig sampling tables operate against good quality. [43]

Let \[ LTPD = 4.0\% \]

Consumer's Risk = 0.10

Batch Size = 500.

OC curves for four different plans are shown in Figure 1, and each curve is to be used for a particular range of process average. In Hill's example producer A submits lots with an average 0.4% defective, B submits 0.5% defective, and C submits 1.3% defective. Yet A will have 80.2% of his lots accepted, B will have 92.5% accepted, and C 86.7% accepted. The lack of incentive to good production is obvious.
Hill outlined a basic economic theory of sampling inspection to account for the incentive features. He assumed

1. the distribution of defectives in submitted lots is binomial with parameter $p$, and
2. the unit cost of manufacture as a function of outgoing quality, $h(p)$, can be represented by $K/\sqrt{P}$.

Let $x =$ number of defectives observed in the sample,

$N =$ batch or lot size,

$LN =$ price paid per accepted batch,

$T =$ average profit to the producer per batch,

$n =$ sample size,

$a =$ acceptance number,

$P =$ proportion of accepted batches.
Since \( x \) is binomial, \( P = \sum_{x=0}^{a} \binom{n}{x} p^x (1-p)^{n-x} \). Under the assumption that a rejected batch is a total loss to the manufacturer, \( T = NLP - Nh(p) \), or
\[
T = KN \left[ \frac{L}{K} P - \frac{1}{\sqrt{p}} \right].
\]
Given \( K \), \( L \), and \( a \), the manufacturer aims to adjust \( p \) to maximize \( T \).

We can derive the relationship the optimal \( p \) must satisfy. First
\[
\frac{dP}{dp} = \frac{d}{dp} \sum_{x=0}^{a} \binom{n}{x} p^x (1-p)^{n-x} = \frac{-n!}{a!(n-a-1)!} p^a (1-p)^{n-a-1}.
\]
Then
\[
\frac{dT}{dp} = KN \left[ \frac{L}{K} \frac{dP}{dp} + \frac{1}{2\sqrt{p}} \right] = 0.2 \text{ Or } \frac{L}{K} \left[ \frac{n!}{a!(n-a-1)!} p^a (1-p)^{n-a-1} \right] = \frac{1}{2\sqrt{p}}.
\]
Figure 2 indicates the variations in producer profit, \( T \), with submitted quality under a sampling plan in which \( n = 150 \), and \( a = 11 \). The plan is a MIL-STD-105A plan for an AQL of 4%.

Hill notes that the variation of the optimal \( p \) is relatively insensitive to changes in the price-cost ratio. Furthermore the sampling plan used causes the producer's optimal choice of \( p \) to be in the neighborhood of 4%, the quality the plan is designed for. Hill did not extend his remarks beyond this example, but his essential points are clear for the purchasing agency. If the producer attempts to maximize expected profit through choice of quality level, both the sampling plan used and the price paid per lot will determine this choice.

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To insure that \( T \) is actually maximized the producer would in fact have to test the end points and all interior points at which \( dT/dp = 0 \).
Figure 2

Johns-Lieberman Model

In 1961 Vernon M. Johns and Gerald J. Lieberman of Stanford University formulated the contingent pricing problem in a more general setting. [44] They described a situation in which the buyer agrees to purchase batches of items, for a price to be determined in a prescribed way according to the results of a sample of items from the lot. The basic assumptions in the Johns-Lieberman model were that

1. the producer is capable of controlling the quality of his product to a known extent and at a cost known to both producer and purchaser, and

2. given any sampling plan and pricing policy the producer will choose the quality to maximize expected profit.
They further assumed that the cost of sampling is charged to the buyer and that the buyer chooses the sample size in advance. Johns and Lieberman note,

The situation is then formally that of a two person non-zero-sum game in which the strategy of one player (i.e., the pricing policy of the consumer) is revealed to the other player (the producer) in advance of his choice of a strategy. This type of game has a well defined notion of a solution. The consumer knows the producer will adopt a strategy (quality level) which will net him (the producer) the maximum expected return under the pricing policy chosen by the consumer. The consumer must therefore select the pricing policy which will net him the greatest return under the producer's corresponding optimal strategy.

The following example illustrates the model and certain problems that arise in this formulation.

Let

\[ N = \text{number of items in a batch}, \]
\[ p = \text{probability each item produced is defective}, \]
\[ h(p) = \text{unit cost of producing items at an average quality level } p, \text{ where } h(p) \text{ includes a normal loading for unit profit}, \]
\[ n = \text{sample size}, \]
\[ x = \text{number of defectives observed in the sample}, \]
\[ c = \text{sampling cost per item}, \]
\[ N^p(x) = \text{price paid for the lot when } x \text{ defectives are observed}, \]
\[ T_c = \text{expected consumer profit}, \]
\[ T_p = \text{expected producer profit}, \]
\[ g(p) = \text{expected value of } \varphi(x), \text{ given } n \text{ and } p. \text{ (If } x \text{ is a binomial random variable} \]
\[ g(p) = \sum_{x=0}^{n} \varphi(x) \binom{n}{x} p^x (1-p)^{n-x}. \]
For simplicity in exposition let the value to the consumer be one per non-defective item, and zero per defective item. Denote the expected unit value by $V(p)$. In this case $V(p) = (1-p)$. The expected net returns to consumer and producer are then

\begin{align*}
(2.1) & \quad T_c = N(1-p) - Ng(p) - nc, \quad \text{and} \\
(2.2) & \quad T_p = Ng(p) - Nh(p). 
\end{align*}

Since a unit profit is included in the $h(p)$ function, we assume that the producer will agree to a contract if $T_p > 0$ for some $p$. For fixed $n$, $T_c$ will be maximized when $Ng(p)$ is minimized, and from this we see that in an optimal policy we would like

\begin{align*}
(2.3) & \quad Ng(p) = Nh(p) \quad \text{for some } p. 
\end{align*}

At that $p$,

\begin{align*}
(2.4) & \quad T_c = N(1-p) - Nh(p) - nc. 
\end{align*}

For any $n$, (2.4) will be maximized by the $p^*$ for which $(1-p) - h(p)$ is a maximum. In this case $p^*$ satisfies $h'(p^*) = -1$. The situation is shown in figure 3. Consumer profit, $T_c$, will be maximized by using the smallest sample size, $n$, at which the producer can be motivated to produce at $p^*$. 

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Consider the case \( n = 1 \). To motivate production at \( \hat{p} \), \( T_{\hat{p}} \) must attain its maximum at \( \hat{p} \), or

\[
(2.5) \quad \frac{d}{dp} [N_g(p) - N_h(p)] = 0, \quad p = \hat{p}.
\]

If \( x \) is a binomial random variable, we have from the definition of \( g(p) \),

\[
(2.6) \quad g(p) = \varphi(0) \cdot (1-p) + \varphi(1) \cdot p,
\]

and from (2.5) \( -\varphi(0) + \varphi(1) = h'(\hat{p}) \), while from (2.3)

\[ \varphi(0) (1-p^*) + \varphi(1) p^* = h(p^*). \]

Solving for \( \varphi(0) \) and \( \varphi(1) \) we obtain the pricing schedule,

\[
(2.7) \quad \varphi(0) = h(p^*) + p^*, \quad \text{and}
\]

\[
(2.8) \quad \varphi(1) = h(p^*) + p^* - 1.
\]
A sample of one produces an optimal policy! Consider the following numerical example based on the preceding development and illustrated in Figure 4. Let the value to the customer, $V(p)$, be $(1-p)$ as before.

Let $h(p) = \frac{.022}{\sqrt{p}} + .4$, unit production cost, and $c = .05$, sampling cost, $N = 100$, lot size.

We determine optimal quality, $p^* \approx .049$, and $h(p^*) = 0.495$. Therefore

$$\varphi(0) = 0.544,$$ and $$\varphi(1) = -0.456$$

![Figure 4](image.png)

This policy yields an expected consumer gain, $T_c$, of 49.55 and no producer gain beyond the unit profit included in the $h(p)$ function. That is, $g(p^*) = h(p^*)$.

Note that if the sample of size one is defective the producer incurs an actual net loss of $(.495) \times 100 + (.456) \times 100 = 95.1$, while if the sample is non-defective his actual gain is only 4.9. This is based on an investment of close to 49.5, and while the expected profit and actual
profit are acceptable at \( p^* \), we feel that it is unlikely that a producer would agree to this type of policy.

One of the elements we feel is intuitively unsatisfying is the magnitude of the penalty levied against the producer when a defective is discovered. We will discuss the implications of bounding the minimum payment in a later section, but to illustrate that lower bounds do not provide the solution, consider the producer to have no fixed costs, so that his unit production cost, \( h(p) \), is \( \frac{0.22}{\sqrt{p}} \), and restrict the payments, \( \varphi(x) \), to be non-negative. The optimal \( p^* \) is still 0.049, but \( h(0.049) = .1 \). Let a sample of nine items be drawn and the price schedule be as follows:

\[
\begin{align*}
\varphi(x) &= 0.175, \quad x = 0 \\
\varphi(x) &= 0.0, \quad x > 0.
\end{align*}
\]

The producer is still motivated to produce at \( p^* \), \( (T_p \) is maximized at \( p^* \)), and expected producer profit at \( p^* \) is 1.10. One defective observation in this case is not as disastrous for the producer as in the previous case since he loses only his investment and is not penalized additionally. However Figure 5 indicates that \( g(p) \) is very close to \( h(p) \) in the vicinity of \( p^* \), and therefore the producer is not severely penalized for quality poorer than \( p^* \). For example, if the producer chooses to produce at \( p = 0.10 \) rather than \( p = 0.049 \) his expected net loss is only .175. He will obtain an actual gain of 10.5 39% of the time, and will lose 7.0 61% of the time.
A producer prone to corner cutting or gambling might find this policy attractive, and we suspect that a buyer would want tighter discriminatory ability.

**Flehinger-Miller Model**

The last numerical example used a sample size of nine and a single payment of 0.175 if no defectives were observed. This was not arbitrary, but was an optimal policy in a way that will now be described. In 1964 Betty J. Flehinger and James Miller reported a "product improvement" model very similar to the Johns-Lieberman model. [45]

They assume that a producer is already manufacturing at some verified quality level and is presumably receiving an acceptable fee. The producer can spend additional funds to improve the product quality. If this improvement benefits the consumer more than it costs the producer, such improvement will be advantageous to both parties. An acceptance test
will be performed to verify that the product improvement has been achieved. The fact that product improvement is advantageous to both parties prompts Flehinger and Miller to seek only admissible policies—those that maximize joint profit.

The important assumptions in this formulation not included in the previous models are:

1. If the producer makes no attempt to improve his quality it will be characterized by the basic quality at which he is currently producing, $p_0$.

2. A strategy or policy, $(n, \varphi_0, \varphi_1, \ldots, \varphi_n)$, or $(n, \varphi)$ is defined to be admissible if it yields a positive expected profit to both consumer and producer, and if no other policy yields a greater expected profit to both sides.

Assume the expected value to the consumer is $V(p) = 1 - p/p_0$. The expected profits can then be written as in (2.1), (2.2) as functions of $p$, $\varphi$ and the sample size, $n$:

\[
T_C(p, \varphi, n) = NV(p) - Ng(p) - nc , \quad \text{and}
\]

Flehinger and Miller feel that this test is motivational rather than informative, since the test is the mechanism by which the producer is motivated to a particular quality. Actually the test is informative to the consumer. While the producer is assumed to have perfect control and can set the quality precisely where he chooses to, the consumer has no protection other than the assumed rationality (i.e., profit maximization behavior) of the producer. Hence the acceptance test both motivates the producer, and gives the consumer some indication of the quality actually attained.

We will denote the unit prices both by $\varphi(x)$ and $\varphi_x$. 

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(2.10) \[ T_p(p, \varphi, n) = N_g(p) - N_h(p). \]

Given any policy \((n, \varphi)\) the producer will choose \(P\) to maximize (2.10). Denote this \(P\) by \(P^*\). We may express \(P^*\) as \(P^*(\varphi, n)\), and express the consumer and producer profits at \(P^*\) as

\[
T^*_c = T_c(P^*, \varphi, n), \quad \text{and} \quad T^*_p = T_p(P^*, \varphi, n).
\]

Each \((n, \varphi)\) is thus mapped into a point \((T^*_c, T^*_p)\) and by Assumption 2 a policy is admissible if it maps into a point \((T^*_c, T^*_p)\) such that \(T^*_c \geq 0, T^*_p \geq 0\) and no other policy maps into a point \((T'_c, T'_p)\) that dominates \((T^*_c, T^*_p)\). This implies that an admissible strategy maps into a point \((T^*_c, T^*_p)\) such that \(T^*_c + T^*_p\) is maximum over all policies leading to the same value of \(T^*_p\).

Since \(T^*_c + T^*_p = N[V(p) - h(p)] - nc\), we see that by seeking admissible strategies, we arrive formally at the relation (2.4) which says simply that if a quantity is to be divided with a fixed amount going to one party, the other party maximizes gain by finding the point at which the initial quantity is greatest.

Flechinger and Miller derive the following procedure to find the set of all admissible strategies:

a. Find the \(P^*\) which maximizes \(V(p) - h(p)\),

b. Out of the class of all policies, \((n, \varphi)\), select those for which \(T^*_p\) is maximum at \(P^*\).

c. Of these policies select those for which \(T^*_c \geq 0, \text{ and } T^*_p \geq 0\).
d. Classify the policies selected by the value of \( T_p^* \), and out of all policies yielding the same value of \( T_p^* \), choose those for which the sampling cost is minimum. These remaining policies are admissible.

We present an example to illustrate the method.

Let \( p = \) the probability that each item produced is defective,
\[
(p \leq p_0)
\]
\( N = \) lot size,
\( n = \) sample size,
\( N^p(x) = \) price paid for the batch,
\( NV(p) = N[1 - p/p_0] \), expected batch value,
\( h(p) = h_1 \ln(p_0/p) \) the unit cost of improving quality to \( p \),
\( c = \) unit cost of sampling,
\( q = 1 - p \).

To obtain the \( p \) that maximizes \( V(p) - h(p) \) in this example we set
\[
\frac{d}{dp} [V(p) - h(p)] = 0 \quad \text{or} \quad -\frac{1}{p_0} = \frac{-h_1}{p^*} \quad \text{and hence},
\]

\[
(2.11) \quad p^* = p_0^{1/h_1}.
\]

To find the policies \( (n, p) \) which maximize \( T_p \) at \( p^* \), set
\[
\frac{d}{dp} [Ng(p) - Nh(p)] = 0 \quad \text{for} \quad p = p^*, \quad \text{which is}
\]
\[
\sum_{x=0}^{n} \varphi_x \frac{d}{dp} \binom{n}{x} p^*(1-p)^{n-x} + \frac{h_1}{p} = 0 \quad \text{for} \quad p = p^*, \quad \text{or}
\]

\[
(2.12) \quad \sum_{x=0}^{m} \varphi_x \binom{n}{x} p^x q^{x-1} = \frac{h_1}{p} \quad \text{for} \quad p = p^*.
\]
This is a condition that $\varphi(x)$ must satisfy for every $n$ in order that $T_p$ will be maximum at $p^*(0,1)$.\footnote{We remark again that every policy, $(n,\varphi)$, must be tested at the endpoints, $p = 0$, $p = 1$, and at every interior point for which the derivative $= 0$. If any policy yields more than one point at which the derivative $= 0$, an additional constraint can be used to find a policy which does maximize $T_p$ at $p^*$.}

Consider all policies of the form $\varphi(0) = L$, a constant to be determined, and $\varphi(x) = 0$, $x > 0$. (2.12) then reduces to $\varphi(0) = \frac{h_1}{p^* n q^* n - 1}$, or since $p^* = p_q h_1$,

\begin{equation}
L = \frac{1}{p_q n q^* n - 1},
\end{equation}

and $g(p)$ is then given by $\frac{q^*}{np_q}$. We thus obtain

\begin{equation}
T_c^* = NV(p^*) - Nq^*/np_q - nc, \quad \text{and}
\end{equation}

\begin{equation}
T_p^* = Nq^*/np_q - Nh(p^*).
\end{equation}

For this special class of policies we now choose sample sizes, $n$, that yield $T_c^* \geq 0$, and $T_p^* \geq 0$. Since $L$ is a function of $n$, we first determine the minimum $n$ such that $T_c^* \geq 0$. Denote this by $n_1$. This is given by

\begin{equation}
n_1 \geq \frac{NV(p)}{2c} - \frac{1}{2} \sqrt{\left(\frac{NV(p^*)}{c}\right)^2 - \frac{4Nq^*}{p_q c}}.
\end{equation}
Denoting by $n_2'$ the largest $n$ such that $T^*_p \geq 0$, we obtain

$$\tag{2.17} n_2' \leq q^* / p_o \ h(p^*) .$$

The requirement that for a fixed value of $T^*_p$ we choose those policies $(n, L)$ such that $n \cdot c$ is minimum implies that the forward difference of $T^*_c$ with respect to $n$ must be non-negative. Therefore $n_2$ may not be given by (2.17) but rather by $\Delta n[NV(p^*) - Nq^*/p_o - nc] \geq 0$, or

$$\tag{2.18} n_2''(n_2'' + 1) \leq Nq^*/cp_o .$$

Hence $n_2$ is given by $\min(n_2', n_2'')$. This is shown in Figure 6.

![Figure 6](image.png)

To summarize, all admissible policies of the special form $\varphi(0) = L$, $\varphi(x) = 0$ $(x > 0)$ are given by

$$\tag{2.13} L = 1 / p_o nq^* h^{-1}$$
for \( n \) in an interval \((n_1, n_2)\) where \( n_1 \) is given by (2.16), and \( n_2 \) satisfies both (2.17) and (2.18). This special class of policies will be proven admissible, and except for one special case exhausts the class of admissible policies.

The special case arises in the following way. The assumption 1 and the cost function, \( h(p) \), imply that even if no improvement effort is made by the producer, quality will remain at level \( p_0 \), and at that level there is a positive probability, \((1-p_0)^n\), that zero defectives will be observed in a sample of size \( n \). We have observed that \( T_p(p^*, L, n) \) decreases in \( n \), and eventually becomes negative. However, \( T_p(p_0, L, n) \) never becomes negative, which indicates that in some cases there may exist a sample size, say \( \tilde{n} \), at which the producer will choose to forego product improvement. This situation is illustrated in Figure 7 and Figure 8.

![Graph showing expected payment](image-url)

\[ h(p) \]
\[ n < \tilde{n} \]
\[ T_p(p^*) > T_p(p_0) \]

**Figure 7**
Flehinger and Miller constrain this situation from occurring by having the upper bound on sample sizes for the policies \((n, L)\) also satisfy the condition that \(T_p(p^*, \phi, n) \geq T_p(p_0, \phi, n)\), and if necessary, add an extra payment \(\varphi(1)\) at sample sizes above \(\bar{n}\), that "remotivate" the producer to \(p^*\).

The following numerical example illustrates the results just derived.

Let \(N = 100\)
\[p_0 = 0.1\]
\[h(p) = .1 \ln .1/p\]
\[c = 0.4\]

From (2.11) the optimal quality, \(p^*\), is 0.01, and the unit production cost, \(h(0.01) = 0.23\). From (2.13) \(L = 10/n(0.99)^{n-1}\). (2.16) yields the minimum sample size, \(n_1\), of 11. (2.17) yields \(n_2 = 43\), and (2.18) yields \(n_2 = 49\). Several of the admissible policies are listed in Table 1, and the situation is displayed in Figure 9.
Table 1

<table>
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<tr>
<th>n</th>
<th>L</th>
<th>$T_p^* \times 100$</th>
<th>$T_c^* \times 100$</th>
<th>$T_c^* + T_p^* \times 100$</th>
<th>Probability of Producer receiving payment</th>
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<td>10</td>
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<td>72.00</td>
<td>.904</td>
</tr>
<tr>
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<td>70.00</td>
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<td>41.50</td>
<td>68.00</td>
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<td>58.80</td>
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<td>.65</td>
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<tr>
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<td>.351</td>
<td>-0.50</td>
<td>58.90</td>
<td>58.40</td>
<td>.64</td>
</tr>
</tbody>
</table>

Figure 9

We now prove that the special class of policies given by $\varphi(0) = L$, $\varphi(x) = 0, x > 0$, is admissible. By the definition of admissibility any admissible policy, $(n, \varphi)$, must maximize $T_p^*$ at $p^*$. Therefore consider
any policy, \((n', \varphi')\), not of the type given by (2.13), (2.16), (2.17), (2.18) that maximizes \(T'_p\) at \(p^*\), and assume that \((n', \varphi')\) maps into \((T'_c, T'_p)\) such that \(T'_c(p^*, \varphi', n') + T'_p(p^*, \varphi', n') > T'_c(p^*, L, n) + T'_p(p^*, L, n)\). This is \(V(p^*) - h(p^*) - c'n' > V(p^*) - h(p^*) - c'n\) and hence, \(n' < n\).

Since \(n_2\) was chosen such that \(T'_c\) increases in \(n\) for \(n_1 \leq n \leq n_2\), we have \(T'_c(p^*, L, n') < T'_c(p^*, L, n)\). By definition

\[
T'_c(p^*, L, \varphi_1, \ldots, \varphi_{n'}, n') = NV(p^*) - c'n' - \varphi(0)q^{n'} - \sum_{x=1}^{n'} \varphi(x) \cdot \binom{N}{x} p^* x q^{n'-x},
\]

and since the last term, \(\sum_{x=1}^{n'} \varphi(x) \cdot \text{pr}[x|n', p^*]\), is positive,

\[
T'_c(p^*, L, \varphi_1, \varphi_2, \ldots, \varphi_{n'}, n') < T'_c(p^*, L, n')\]

Therefore

\[
T'_c(p^*, L, \varphi_1, \ldots, \varphi_{n'}, n') < T'_c(p^*, L, n)
\]

which proves that \((n, \varphi)\) satisfying (2.13), (2.16)-(2.18) is admissible.

Are there admissible strategies other than this class of policies that might yield a larger \(T'_c\)? To answer this denote the special class by \((n^*, \varphi^*)\) and assume first that \(n_2\), the maximum sample size, is given by \(T^*_p = 0\), as in the numerical example. Then if there existed an admissible policy \((n', \varphi')\) such that \(T'_p \geq 0\) and \(T'_c > T'_c\) we would have \(T'_c + T'_p > T'_c + T^*_p\) contradicting the admissibility of \((n^*, \varphi^*)\).
Alternatively assume that \( n_2 \) is given by \( \Delta n T_0 C \geq 0 \). Then if there were an admissible policy \( (n', \varphi') \) such that \( T_p' > T_p^* \) and \( T_c' > T_c^* \), we would have \( \frac{T_p'}{T_p} > \frac{T_c'}{T_c} \) implying \( \frac{T_c'}{T_c} < \frac{T_p'}{T_p} \) which again is a contradiction.

Since \( n \) is a discrete variable, \( n_2 \) may be given by \( (2.17) \) with \( \frac{T_p^*}{T_p} > 0 \). In this case an additional positive payment, \( \varphi(1) \), at the sample size \( n_2 + 1 \), might motivate the producer to \( p^* \), yet increase the consumer's profit. This would depend on the cost of sampling. The paper reporting the Flehinger-Miller model uses an example in which the "sample size" is really a test length. Hence the test length may be varied continuously to drive the \( \frac{T_p^*}{T_p} \) exactly to zero.

We are left with two somewhat unsatisfying results. The Johns-Lieberman formulation implies that the consumer can both maximize his own profit and satisfy the producer with a sample size of one, and the Flehinger-Miller model derives optimal policies which are sampling plans with an acceptance number of one, except in special cases.

To understand these results consider the magnitudes of the actual prices and the confidence with which the results of these plans would be viewed. Consider first the implications of a lower bound on the prices in the Johns-Lieberman formulation. Let \( \varphi(x) \geq m \) for all \( x \), and let \( n = 1 \). From the conditions that both \( T_c \) and \( T_p \) be maximized at \( p^* \), or \( (1-p^*) \varphi_o + p^* \varphi_1 = h(p^*) \), and \( -\varphi_o + \varphi_1 = h'(p^*) \) we derived the solutions \( \varphi_o = h(p^*) + p^* \), and \( \varphi_1 = h(p^*) + p^* - 1 \). If \( h(p^*) + p^* - 1 < m \), we must choose \( \varphi_1 = m \), and hence
\[ q^*\varphi_0 + p^*\varphi_1 = q^*\varphi_0 + p^*\varphi_1 \]
\[ > q^*[h(p^*) + p^*] + p^*[h(p^*) + p^* - 1] \]
\[ = h(p^*) . \]

Thus the existence of a lower bound on the prices shows that in general \( n = 1 \) is not optimal, and the greater the lower bound, the greater will be the sample size necessary to allow the equations

\[ g(p^*) = h(p^*) , \]
\[ g'(p^*) = h'(p^*) \]

to be satisfied simultaneously. The numerical results on page 58 are an instance in which \( m = 0 \), and \( n = 43 \) is the first sample size at which both \( (.99)^n\varphi_0 = .23 \) and \( -n(.99)^{n-1}\varphi_0 = -10 \) can be satisfied.

Now bound the prices above by \( M \), and again let \( n = 1 \). Then since we must have \( -\varphi_0 + \varphi_1 = h'(p^*) \), and \( \varphi_1 \) is bounded below by \( m \), the minimum \( \varphi_0 \) is \( \varphi_0 = m - h'(p^*) \). For large \( m \) or large negative \( h'(p^*) \) it may occur that \( m - h'(p^*) > M \). Again a larger sample size may be required so that \( \varphi_0 \leq M \).

---

6Let \( \pi_x = (\frac{n}{x}) p^x(1-p)^{n-x} \) and \( dx = d/d\pi_x \). Then since

\[ \sum_{x=0}^{n} dx = 0, d_0 = - \sum_{x=1}^{n} dx. \]

Now for \( n > 1 \), and \( \varphi(x) = m, (x > 0) \),

\[ g'(p^*) = h'(p^*) \]

becomes \(-n\varphi_0 q^{n-1} + m[nq^{n-1}] = h'(p^*)\), or

\[ \varphi_0 = m - h'(p^*)/nq^{n-1} \]. For \( p^* < 1/n+1, \Delta n[nq^{n-1}] > 0 \), and since \( h'(p^*) < 0, \varphi_0 \) decreases in \( n \).
Therefore the unbounded prices in the Johns-Lieberman model account for the optimality of small sample sizes. However the numerical results in Table 1 were certainly obtained with prices reasonably bounded relative to costs and values. Yet the sample size of 43 implies that the producer will receive a payment approximately twice in every three submitted batches, even if submitted quality is constant at 1% defective. In the terminology of acceptance sampling the producer is accepting a producer’s risk of 35%! Also apparent is the fact that the consumer cannot specify an arbitrary degree of protection against quality poorer than $p^*$. A less arbitrary objection is that the sample size in the Flehinger-Miller model is determined only by the selected profit division. If, for instance, the producer were content to accept a 30% share of $V(p^*) - h(p^*)$ there would exist a unique $n$ yielding an expected producer profit closest to this agreed on share.

The heart of the difficulty is that by using a policy with only a single payment when no defectives are observed in a sample of $n$, the sample size is a lever to force producer profit down,\textsuperscript{7} rather than an instrument to assure both parties that the payment received and quality claimed by the producer are in fact correct. We feel that a contingent pricing policy should include both flexibility in the determination of bounded prices, and controllable protection for producer and consumer, and we develop and explore the implications and applications of such a model.

\textsuperscript{7}It is in this sense that Flehinger and Miller state that their sampling is non-informative but rather motivational.
CHAPTER III
CONstrained CONTINGENT PRICING MODELS

Introduction

In Chapter II we noted that existing contingent pricing models yield unrealistic policies due to the lack of bounds on the prices and the lack of protection against sampling variation. We observed in Chapter I that maximum expected profit and minimum total product cost are not the sole concerns of producer and purchaser. Both consider risk, and we distinguish here between two types of risk. When the firm accepts a quality-contingent or fixed price contract it gambles on its cost estimates. This is risk based on uncertainty in the production environment and is the risk generally discussed in economic theory. Risk of this type is balanced by the minimum profit the producer considers appropriate for the product. In a contingent pricing contract the producer may actually achieve a specified quality but not receive the proper payment due to variation in the number of defectives in randomly drawn samples. This is statistical risk analogous to "producer's risk" in acceptance sampling. The possibility of overpayment for poor lots corresponds to "consumer's risk" in acceptance sampling. The existence of statistical risk might cause the producer to negotiate for more profit, but risk of this type is usually treated by designing plans that reduce consumer and producer risk below acceptable limits.
Guthrie and Johns [46] developed an elegant model to determine the optimal sample size and acceptance number in single sampling plans without explicit consideration of consumer or producer risk, but they assumed the existence and knowledge of an underlying distribution of the random variable $p$. While the assumption that users can readily state their AQL, producer's risk, LTPD, and consumer's risk is an oversimplification, it would also be unusual to find a consumer confident in his estimate of the underlying prior distribution of the quality of an incoming product—especially a new product. The prior distribution is likely to be nonstationary, and the sampling plan ought to adapt as the distribution changes. But we have noted that the producer's attempts to change the distribution are themselves a function of the sampling plan. We will formulate a model using the concept of consumer and producer risk.

**General Assumptions**

We postulate the existence of a production and procurement situation with the following characteristics.

1. A single type of item is produced. Each item has a constant probability, $p$, of being defective. Batches of $N$ are formed in such a way that the number of defectives, $X$, in each batch is a binomial random variable. Thus $p$ is also the expected fraction defective per batch.

2. The production process is controlled by the producer who chooses $p$ at a unit production cost, $h(p)$, which we assume
strictly convex and decreasing for \( p \) in some interval of interest.\(^1\)

3. After delivery of the items the consumer draws a sample of \( n \) (\( n \leq N \)) in such a way that the number of defectives in the sample, \( x \), is a binomial random variable. The consumer pays \( N\phi(x) \) for the batch of \( N \).

4. Sampling costs are paid by the consumer at \( c \) per unit inspected.

5. The consumer can tolerate some defective items in incoming lots and can state the expected value, \( NV(p) \), of batches of \( N \) items containing \( X \) defectives where \( X \) is a random variable whose distribution is given by the binomial probability law with parameters \( N \) and \( p \). We assume \( V(p) \) concave.

6. The consumer desires to maximize expected net gain by concluding a contract promising quality which maximizes the difference between expected product worth and expected total procurement cost.

7. The consumer desires protection against overpayment when quality is poor. He can state the quality, \( p_b \), he considers poor, and the maximum amount, \( Nv \), he is willing to pay for quality as poor as \( p_b \). Since any batch may yield a sample with

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\(^1\)The basic unit production cost, say \( \bar{h}(p) \), is known to the producer. During negotiation the producer adds a unit profit, say \( zh(p) \), and thus the existence of an \( h(p) = (1+z)\bar{h}(p) \) becomes known to the consumer. The strict convexity is used in obtaining sufficient conditions for the existence of solutions.
very few defectives, the consumer can state the relative
frequency, $\beta$, with which overpayment will be permissible.

8. The consumer will not agree to pay more than some amount, $M$.

9. Given a sample size, $n$, and payment schedule, $\varphi(x)$, the
producer chooses $p$ to maximize expected profit. Denote this
$p$ by $p'$.

10. The producer requires that the minimum price paid (or maximum
penalty levied) be bounded below by $m$.

11. The producer requires that prices be monotonically decreasing
in the number of defectives observed.

12. Since $h(p)$ includes unit profit, the producer will agree to a
contract only if at $p'$ the expected payment is at least equal
to $h(p')$.

13. The producer desires protection against underpayment. He can
state a minimum price, $N_w$, which the contingent pricing policy
should assure him of receiving when quality is no worse than
$p'$. Given that $p \leq p'$, payment less than $N_w$ is permissible
with frequency less than $\alpha$.\(^2\)

**Notation**

$N =$ batch or lot size

$X =$ number of defectives in the batch

$n =$ sample size

$x =$ number of defectives in the sample

\(^2\)We later relax these assumptions to discuss cases in which $h(p)$
contains a random component, the producer's control over $p$ is not
precise, and the producer must replace all defectives discovered during
inspection with nondefective items.
\( c = \text{unit sampling cost} \)

\( p = \text{probability each item produced is defective} \)

\( q = (1-p) \)

\( NV(p) = \text{expected value of a batch of } N \text{ items when batches are} \)

\( \text{formed according to Assumption 1} \)

\( N\Phi(x) = \text{price paid for the batch of } N \text{ when } x \text{ defectives are} \)

\( \text{observed} \)

\( h(p) = \text{unit production cost to attain quality level } p, \text{ including} \)

\( \text{unit profit} \)

\( p_b = \text{quality level considered poor} \)

\( N_v = \text{maximum payment desirable when } p \geq p_b \)

\( \beta = \text{permissible frequency that } N\Phi(x) \text{ may be greater than } N_v \)

\( \text{when } p \geq p_b \)

\( p' = \text{quality at which expected producer profit is maximized given} \)

\( n \text{ and } \Phi(x) \)

\( N_w = \text{minimum payment desirable when } p \leq p' \)

\( \alpha = \text{permissible frequency that } N\Phi(x) \text{ may be less than } N_w \)

\( \text{when } p \leq p' \)

\( g(p) = \text{expected value of } \Phi(x) \text{ given } p \text{ for fixed } n \)

\( m = \text{lower bound on possible payments, } \Phi(x) \)

\( M = \text{upper bound on possible payments, } \Phi(x) \)

\( T_c = \text{consumer expected profit} \)

\( T_p = \text{producer expected profit} \)

\( p_o = \text{quality which maximizes } V(p) - h(p) \)
\[ \pi_x = \binom{n}{x} p^x (1-p)^{n-x} \]
\[ \delta_x = \frac{d}{dp} \pi_x \]

**Basic Model**

The consumer and producer expected net gains are respectively

\[(3.1) \quad T_c = NV(p) - Ng(p) - nc \quad \text{and} \]
\[(3.2) \quad T_p = Ng(p) - Nh(p). \]

Assumptions 1-13 imply that no contract is acceptable to both parties unless (3.3)-(3.6) are satisfied.

\[(3.3) \quad \Pr[N \varphi(x) \geq Nv | p \leq p'] \geq 1 - \alpha \]
\[(3.4) \quad \Pr[N \varphi(x) \leq Nv | p \geq p_b] \geq 1 - \beta \]
\[(3.5) \quad M \geq \varphi(0) \geq \cdots \geq \varphi(n) \geq m \]
\[(3.6) \quad Ng(p') - Nh(p') \geq 0 \]

The consumer seeks to maximize (3.1) by choosing \( n, \varphi(0), \ldots, \varphi(n) \)

which satisfy (3.3)-(3.6) knowing that the producer will choose \( p \) to maximize (3.2).\(^5\)

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\(^3\) If it is necessary to stipulate the \((n,p)\) at which \( \pi_x \) and \( \delta_x \) are defined we write \( \pi_x(n,p) \) and \( \delta_x(n,p) \).

\(^4\) We also refer to \( \varphi(0), \ldots, \varphi(n) \) as \( \varphi_0, \ldots, \varphi_n \) and to \( n, \varphi_0, \ldots, \varphi_n \) as \( (n,\varphi) \).

\(^5\) In contrast to the Flehinger-Miller model, with this formulation of the problem, admissible policies are optimal for the consumer only if (3.6) holds with equality at \( p = p_0 \) for some \( n \), and there is no smaller \( n \) yielding a higher consumer profit.
Theorem 1

For fixed $n$ and $p$, say $(n, p_o)$ where $0 < p_o < 1$, it is necessary that a set of prices, $\varphi_x$, maximizing (3.1) while simultaneously maximizing (3.2) at $p_o$ and satisfying (3.3)-(3.6) be solutions to the following linear programming program.\(^6\)

\[
\text{(3.1') minimize } \sum_{i=0}^{n} \pi_x(n, p_o) \varphi_x
\]

\[
\text{(3.2') } \sum_{i=0}^{n} d_x(n, p_o) \varphi_x = h'(p_o)
\]

\[
\text{(3.3') } \varphi_{k_1} \geq \varphi_{k_2} \geq \varphi_{k_3} \geq \cdots \geq \varphi_{k_n} \geq \varphi_0
\]

\[
\text{(3.4') } \varphi_{k_1} \leq \varphi_{k_2} \leq \varphi_{k_3} \leq \cdots \leq \varphi_{k_n} \leq \varphi_{k_0}
\]

\[
\text{(3.5') } M = \varphi_0 = \cdots = \varphi_n = m
\]

\[
\text{(3.6') } \sum_{i=0}^{n} \pi_x(n, p_o) \varphi_x \geq h(p_o)
\]

Proof:

By (3.1) maximizing $T_c$ at $(n, p_o)$ is equivalent to minimizing $g(p_o)$ by choosing $\varphi_0, \ldots, \varphi_n$, and hence (3.1) is written as (3.1'). A condition necessary for $T_p$ to be maximized at $p_o \in (0, 1)$

\(^6\) $k_{\alpha}, k_{\beta}$ are described in the proof.
is $\frac{d}{dp} [g(p) - h(p)] = 0$ at $p = p_0$, or

$$\sum_{x=0}^{n} \lambda_x(n, p_0) \varphi_x = h'(p_0),$$

which is (3.2'). The constraints (3.5) and (3.6) are already linear in $\varphi_x$.

The constraint (3.3) requires specifically

$$\text{pr}[N\varphi(x) \geq N\varphi | p = p_0] \geq 1 - \alpha,$$

which is equivalent to

$$\sum_{k=0}^{n} \text{pr}[\varphi(k) \geq w] \text{pr}[x = k | p = p_0] \geq 1 - \alpha.$$

Let $k''$ be the largest integer for which $\varphi(k) \geq w$. From (3.8) we must have $\sum_{k=0}^{k''} \lambda_x(n, p_0) \geq 1 - \alpha$. Let $k_\alpha$ denote the smallest integer $k_\alpha$ such that $\sum_{k=0}^{k_\alpha} \lambda_x(n, p_0) \geq 1 - \alpha$. In order for (3.8) to be satisfied, it is necessary that $k'' \geq k_\alpha$, and by the monotonicity of $\varphi(x)$, we see that $\varphi(k_\alpha)$ cannot be less than $w$.

Thus (3.3) implies (3.3'). By similar reasoning (3.4) implies (3.4'). Note also that (3.3') with $k_\alpha$ defined above implies (3.7). We now show that (3.7) implies (3.3), and that (3.4') implies (3.4).

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7That $\varphi(x)$ satisfies (3.2') is not sufficient to insure that $\varphi(x)$ satisfies (3.2). However, as will be seen from the remainder of the proof, if a test of the points 0, 1, and all points for which $g'(p) = h'(p)$ shows that $T_p$ is maximized at $p_0$, then an optimal solution to (3.1')-(3.6') is also an optimal solution to (3.1)-(3.6). If $T_p$ is maximized at $p_2 \neq p_0$, we can add the constraint $T_p(p_0) > T_p(p_2)$ and determine a new set of prices.
For \( j \leq n \), \( \frac{d}{dp} \sum_{j}^{n} \pi_x(n, p) = \frac{-n!}{j!(n-j-1)!} p^j q^{n-j-1} \leq 0 \). Hence \( p < p_0 \)

implies \( \sum_{0}^{k} \pi_x(n, p) \geq \sum_{0}^{k} \pi_x(n, p_0) \). Thus if (3.7) is satisfied, then (3.3) is satisfied.

Let \( k_B \) be the largest integer for which

\[
\sum_{k_B}^{n} \pi_x(n, p_B) \geq 1 - \beta ,
\]

and require \( \phi_{k_B} \leq \nu \). We then have

\[
pr[N\phi(x) \leq N\nu | p = p_B] = \sum_{0}^{n} pr[\phi_k \leq \nu] pr[x = k | p_B]
\geq \sum_{k_B}^{n} pr[\phi_k \leq \nu] pr[x = k | p_B] \geq 1 - \beta .
\]

By reasoning as in the previous case we verify that (3.4') holding for \( k_B \) determined by (3.9) satisfies (3.4) for \( p > p_B \). This completes the proof that solutions to (3.1)-(3.6) for fixed \( (n, p) \) are solutions to (3.1')-(3.6').

**Variable Transformation**

Before demonstrating a computational simplification in (3.1')-(3.6')

we indicate the effect that (3.3), (3.4) have on the final prices and sample size. Let \( \alpha = .10, \beta = .15, p_0 = .05, p_B = .30, w = .50, \nu = .20, \) and \( n = 9 \). Then \( \sum_{0}^{1} \pi_x(9, .05) = .928 > .9, \) and \( \sum_{1}^{9} \pi_x(9, .3) = .96 > .85 \). Thus \( k_\alpha = 1, k_B = 1 \) and (3.3'), (3.4') require \( \phi_1 \geq .5 \) and \( \phi_1 \leq .2 \), a contradiction. Therefore \( n = 9 \) permits no feasible prices for \( p_0 = .05 \). At \( n = 10, k_\alpha = 1, k_3 = 2, \)
and there may exist a feasible set of prices. If so, they will be
constrained by $\Phi_1 \geq 5, \Phi_2 \leq 2$, indicating that both the minimum sample
size and the form of the pricing policy are affected by (3.3) and (3.4).

The transformation

$$y_j = \Phi_j - \Phi_{j+1}, \quad j = 0, 1, \ldots, n-1 \tag{3.10}$$
$$y_n = \Phi_n - m$$

yields a linear programming problem equivalent to (3.1')-(3.6') and
containing only 5 constraints. By (3.10), $\Phi_j = m + \sum_{i=j}^{n} y_i$, and choosing
$\Phi_x$ to minimize (3.1') is equivalent to choosing $y_i$ to minimize

$$m + \sum_{i=0}^{n} \sum_{x=0}^{i} y_i \pi_x.$$ 

Let $p_i^o = \sum_{x=0}^{i} \pi_i(n, p_o)$, and $d_i^o = \sum_{x=0}^{i} d_x(n, p_o)$. The transformed
problem follows.

$$(3.1'') \quad \text{Minimize} \quad \sum_{i=0}^{n} p_i^o y_i$$

$$(3.2'') \quad \sum_{i=0}^{n} d_i^o y_i = h'(p_o)$$

$$(3.3'') \quad \sum_{i=k}^{n} y_i \geq m - w$$

$$(3.4'') \quad \sum_{i=k}^{n} y_i \leq v - m$$

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\[ \sum_{i=0}^{n} y_i \leq M - m \]

\[ \sum_{i=0}^{n} \prod_{i} y_i \geq h(p_o) - m \]

Since this problem need have \( y_i \neq 0 \) for at most five \( y_i \), there need be at most six distinct price levels, \( \varphi_x \), one of which may be \( m \). If any of the constraints (3.3")-(3.5") are redundant the number of different price levels will be correspondingly reduced. If, for example, \( \alpha = 1 \), \( \beta = 1 \), \( M = L \), and \( m = -L \), where \( L \) is positive and very large, we obtain the price schedules characterized in Chapter II, since the constraint (3.5") will not be active if \( M - m \) is very large, and the constraints (3.3") and (3.4") can be satisfied by considering only \( y_o \) and \( y_n \). We can select \( y_o \), \( y_n \) to satisfy (3.2")-(3.4"), and (3.6"), setting all other \( y_i = 0 \). The price schedule is then

\[ \varphi_o = m + y_o + y_n \]

\[ \varphi_x = m + y_n \quad (x = 1, \ldots, n) \]

The number of active constraints, and thus the number of distinct price levels, \( \varphi_x \), depend on the parameters, \( \alpha \), \( \beta \), \( M \), \( m \), \( w \), and \( v \).

Policies Optimal at \( p_o \)

Every point \((n,p')\) is either infeasible for (3.2")-(3.6") or yields a value of (3.1") and a set of \( \varphi_x \) with the property that \((n, \varphi)\) induces the manufacturer to select \( p' \). Let \( R \) denote the set of \((n,p)\) for which \( \varphi_x \) exist satisfying (3.1")-(3.6"), and let \( G(n,p) \)
denote the value of (3.1') for \((n,p)\in R\). We consider conditions on
\(M, m, w, \text{ and } v\) under which there exists a policy \((n,\varphi)\) maximizing
\(T_p\) at \(p_o\), and yielding \(G(n,p_o) = h(p_o)\). Assume \(V(p), h(p), \alpha, \beta,\)
and \(p_o\) fixed. There are twelve possible cases, determined by the
relations \(w > v\), and the order relations among \([np], k_\alpha, \text{ and } k_\beta\).
The most restrictive cases are those for which \(w > v\). We will consider
one case with \(w > v\) and one with \(w \leq v\).

**Proposition 1**

Let \(D^b_a = \sum_a^b d_x(n,p_o)\)

\[ p^b_a = \sum_a^b \pi_x(n,p_o), \text{ and} \]

\[ s = [np]. \]

Case I:

Assume \(w \leq v\) and \(k_\alpha < s < k_\beta\). If

\[(3.11) \quad (M-m) \geq \frac{h'(p_o)}{D^s_o}, \text{ and} \]

\[(3.12) \quad m = [h(p_o) - w] + \frac{h'(p_o)}{(M-m)} \left[ s - (M-m) \frac{D^s_o}{p_o} \right] \]

then \(G(n,p_o) = h(p_o)\).

**Proof:**

Note first that \(w \leq v\) implies that \(\varphi_1\) satisfying (3.3) and
(3.4) can be found for all \(n \geq 1\). For \(\sum_{o}^{k_\alpha} \pi_x(n,p_o) \geq 1 - \alpha\) can always
be satisfied by $k_{\alpha} = n$ and $\sum_{k_B}^n \pi_{x}(n, p_b) \geq 1 - \beta$ can be satisfied with $k_B = 0$. In this event we have $v \geq \varphi_0 \geq \varphi_n \geq w$ which does not violate (3.5).

Therefore, for any $n$ we have $\max \sum_{\varphi}^n \delta_{i} \varphi_i = w \sum_{\varphi}^n \delta_{i} = 0$ for $\varphi_i$ satisfying (3.3)-(3.5). Also, $\min \sum_{\varphi}^n \delta_{i} \varphi_i = M \delta_{o}^s + m \delta_{s+1}^n = (M-m) D_{o}^s$. By (3.11) $(M-m) \geq h'(p_o)/D_{o}^s$, hence $(M-m) D_{o}^s \leq h'(p_o)$, and we obtain

\[ \min_{\varphi} \sum_{\delta_{i} \varphi_i} \leq h'(p_o) < \max_{\varphi} \sum_{\delta_{i} \varphi_i}. \]

Therefore there exist $\varphi_x$ satisfying (3.2)-(3.5). Assume that $\min_{\varphi} \sum_{\delta_{i} \varphi_i} < h'(p_o)$. Then for $0 < \lambda < 1$,

\[ h'(p_o) = \lambda (M-m) D_{o}^s + (1-\lambda) w D_{o}^n = \lambda (M-m) D_{o}^s, \quad \text{and} \]

\[ \lambda = \frac{h'(p_o)}{(M-m) D_{o}^s}. \]

We now show that $\varphi_x = \lambda M + (1-\lambda) w$, $0 \leq x \leq s$

$\varphi_x = \lambda m + (1-\lambda) w$, $s \leq x \leq n$

yield equality in (3.6). By (3.12)

\[ m = h(p_o) - w + \frac{h'(p_o)}{(M-m) D_{o}^s} [w - (M-m) p_{o}^s], \]

or by (3.14)

\[ m = h(p_o) - w + \lambda [w - (M-m) p_{o}^s], \]

which yields

\[ h(p_o) = (\lambda M + (1-\lambda) w) p_{o}^s + (\lambda m + (1-\lambda) w) p_{s+1}^n = G(n, p_o). \]

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8 The strict inequality holds by the original assumption that $h(p)$ is strictly convex and decreasing.
Case II:
Assume \( w > v, k_\alpha < s < k_\beta \), and assume that for some \( n_0 \leq N \),
\( k_\alpha(p_o, n) < k_\beta(p_b, n) \). Then if

\[
(3.15) \quad w - v \leq \frac{h'(p_o)}{k_B^{-1} D_o^S(n_o, p_o)}
\]

\[
(3.16) \quad M - m \geq \frac{h'(p_o)}{D_o^S(n_o, p_o)}
\]

\[
(3.17) \quad \frac{h(p_o) - m - (M-m) P_o^S}{(w-v) P_o B^{-1} + (v-m) - (M-m) P_o^S} = \frac{h'(p_o) - (M-m) D_o^S}{(w-v) D_o B^{-1} - (M-m) D_o^S}
\]

Then \( G(n_o, p_o) = h(p_o) \).

Proof: At \( n = n_o \), \( \max_{\varphi} \sum_i d_i \varphi_i = wD_o + vD_o^r = (w-v) D_o k_B^{-1} \).

By (3.15) \( (w-v) D_o \geq h'(p_o) \), and hence \( \max_{\varphi} \sum_i d_i \varphi_i \geq h'(p_o) \). At

\( n = n_o \), \( \min_{\varphi} \sum_i d_i \varphi_i = M D_o^S + m D_o^S_{s+1} = (M-m) D_o^S \).

By (3.16) \( (M-m) D_o^S \leq h'(p) \),

and hence \( \min_{\varphi} \sum_i d_i \varphi_i \leq h'(p_o) \). Thus there exist \( \varphi_i \) satisfying (3.2)-(3.5). In the most general case

\[
h'(p_o) = \lambda [(w-v) D_o k_B^{-1}] + (1-\lambda) [(M-m) D_o^S]
\]

yielding

\[
\lambda = \frac{h'(p_o) - (M-m) D_o^S}{(w-v) D_o k_B^{-1} - (M-m) D_o^S}
\]

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By (3.17)

\[
\frac{h(p_o) - m \cdot (M-m) P_o^S}{(w-v) P_o^B - (v-m) P_o^S} = \frac{h'(p_o) - (M-m) D_o^S}{(w-v) D_o^B - (M-m) D_o^S} = \lambda
\]

From which

\[
h(p_o) = \lambda \left( \frac{k_B-1}{wP_o^B + v(1-P_o^B)} \right) + (1-\lambda) \left[ M P_o^S + M (1-P_o^S) \right]
\]

or

\[
h(p_o) = \lambda \left( \frac{k_B-1}{wP_o^B + vP_o^n \frac{1}{k_B}} \right) + (1-\lambda) \left[ M P_o^S + m P_o^n \frac{1}{k_B+1} \right],
\]

and (3.6) holds with equality.

When either of these sets of conditions is satisfied, we need only compute \( p_o, n_o \), and generate the optimal prices, \( \varphi(x) \), by solving (3.1')-(3.6').

**Admissible Policies**

Recall that in Chapter II we required of an admissible policy that it maximize joint profit, \( \psi(n,p) = NV(p) - Nh(p) - nc \). Under our assumptions \( \psi(n,p) \) is concave in \( n \) and \( p \), and hence if the set \( R \) is convex in \( n \) for fixed \( p \) and convex in \( p \) for fixed \( n \), we may determine \((n_o, p_o)\) by any one of several algorithms or search methods.

We will later describe the Fibonacci search generalized to two variables and under the above assumptions, this method will determine \((n_o, p_o)\).

To see this, assume as proven that the Fibonacci search procedure will determine \( p_o(n) \) to maximize \( \psi(n,p) \) for fixed \( n \). We will show that
\( \psi(n, p_0(n)) \) is concave in \( n \). By the concavity of \( \psi(n, p) \) and the definition of \( p_0(n) \),

\[
\lambda \psi(n_1, p_0(n_1)) + (1-\lambda) \psi(n_2, p_0(n_2)) \\
\leq \psi(\lambda n_1 + (1-\lambda) n_2, \lambda p_0(n_1) + (1-\lambda) p_0(n_2)) \\
\leq \psi(\lambda n_1 + (1-\lambda) n_2, p_0(\lambda n_1 + (1-\lambda) n_2)) .
\]

Thus the convexity of \( R \) allows efficient determination of the unique \((n_0, p_0)\).

Let \( n = \text{minimum } n \text{ for which } p = 0 \text{ satisfies (3.3), (3.4)}. \)

\( \bar{n} = \text{minimum } n \text{ for which } p = p_0 \text{ satisfies (3.3), (3.4)}. \)

If \( R \) is convex in \( p \) for fixed \( n \) but not in \( n \) for fixed \( p \), we must evaluate \( \psi(n, p_0(n)) \) for every \( n \in [\underline{n}, \bar{n}] \) to determine \((n_0, p_0)\).

There is no feasible \( n < \underline{n} \), and for \( n > \bar{n} \) \( \psi(n, p_0) \geq \psi(n, p_0) \geq \psi(n, p_0(n)) \).

We state as a proposition the conditions under which \( R \) is convex in \( p \) for fixed \( n \). The proof is similar to the proof of proposition 1.

**Proposition 2**

Let \( \bar{p}(n) = \text{the greatest } p \text{ for which } \sum_{1}^{k_B-1} \pi_x(n, p) \geq 1 - \alpha \) where \( k_B \) is the greatest integer for which \( \sum_{1}^{n} \pi_x(n, p_b) \geq 1 - \beta \). Let \( p_1 \) be that \( p \) at which \( h'(p) = L \), where \( L \) is a specified constant.

Let \( p_2 \) be the minimum \( p \) at which \( v(p) = h(p) \). Let \( \underline{p} = \text{max } (p_1, p_2) \).

Let the maximum \( \sum d_i p_1 \) satisfying (4.3')-(4.5') be minimized at \( p' \), and the minimum \( \sum d_i p_1 \) be maximized at \( p'' \), where \( p', p'' \) are in the interval \([\underline{p}, \bar{p}(n)]\).
Assume $p \leq \overline{p}(n)$, $w > v$, and $[np] < k_{\alpha} < k_{\beta}$. Then if

$$\tag{3.18} (w-m) \leq \frac{h'(\overline{p}(n))}{D_{o}^{B-1}(p')} ,$$

$$\tag{3.19} (M-m) \geq \frac{1}{D_{o}^{S}(p'')} \left[ h'(\overline{p}) - \frac{h'(\overline{p}(n))}{D_{s+1}^{O}\left(\overline{p}'\right)} \right] \frac{k_{\alpha}}{D_{o}^{B-1}(p')},$$

and

$$\tag{3.20} h(p) \leq \lambda \left[ w p_{o}^{B-1}(p') + v p_{\alpha}^{n}(p') \right] + (1-\lambda) \left[ m p_{o}^{S}(p'') + w p_{s+1}^{B}(p'') + m p_{\alpha}^{n}(p'') \right]$$

are satisfied where $\lambda$ is given by (3.21) then $R$ is convex in $p$ for fixed $n$, for all $p \in \overline{p}(\overline{p}(n))$.

Proof:

Abbreviate $\overline{p}(n)$ by $\overline{p}$. By definition of $\overline{p}$, $\overline{p}$, (3.3) and (3.4) hold for $p \in \overline{p}(\overline{p})$.

Since $h''(p) > 0$, max $h'(p) = h'(\overline{p})$, and min $h'(p) = h'(\overline{p})$.

$$\text{Max} \sum_{\varphi} d_{i} \varphi_{1} = (w-v) D_{o}^{B-1}(p),$$

$$\text{min} \sum_{\varphi} d_{i} \varphi_{1} = (M-m) D_{o}^{S}(p) + (w-m) D_{s+1}^{B}(p).$$

Therefore by (3.18)

$$\max_{\varphi} \sum_{\varphi} d_{i} \varphi_{1} \geq (w-v) D_{o}^{B-1}(p') \geq h'(\overline{p}) .$$

By (3.19)

$$D_{o}^{S}(p'')(M-m) + \frac{h'(\overline{p})}{D_{s+1}^{O}\left(\overline{p}'\right)} \leq h'(\overline{p}),$$

and by (3.18)

$$\frac{k_{\alpha}}{D_{o}^{B-1}(p')}$$

79
\[ D^S_O(p'')(M-m) + (w-m) D^k_{s+1}(p') \leq h'(p), \text{ or } \min \sum d_i \varphi_i \leq h'(p). \] Hence for all \( p \in \overline{p}, \overline{p} \) there exist \( \varphi_x \) satisfying (3.2)-(3.5). For some \( p \in \overline{p}, \overline{p} \) let
\[
h'(p) = \lambda \left[ (w-v) D^k_{o-1}(p') \right] + (1-\lambda) \left[ (M-m) D^S_O(p'') + (w-m) D^n_{s+1}(p'') \right]
\]
yielding
\[
(3.21) \quad \lambda = \frac{h'(p) - (M-m) D^S_O(p'') + (w-m) D^k_{s+1}(p'')}{(w-v) D^k_{o-1}(p') - (M-m) D^S_O(p'') - (w-m) D^n_{s+1}(p'')}
\]
By (3.20), (3.6) is satisfied, and therefore \( R \) is convex in \( p \) for fixed \( n \) for \( p \in \overline{p}, \overline{p} \).

Conditions (3.18)-(3.20) are not restrictive and frequently equality holds in (3.20). When this is so we may determine \( p_o(n) \) by a search procedure. Convexity in \( n \) for fixed \( p \) is obtained only in small intervals \([n_1, n_2]\). Heuristically this is due to the changes in \( k_\alpha \) and \( k_\beta \) as \( n \) increases. For fixed \((n_1, p_1)\), \( k_\alpha \) can be less than \( k_\beta \), while at \( n_1 + 1 \ k_\alpha \) may increase while \( k_\beta \) remains constant, producing infeasibility at \((n_1+1, p)\). At \( n_1 + 2 \ k_\beta \) may increase, again creating a feasible point, \((n_1+2, p)\). Thus (3.3) and (3.4) can cause nonconvexity in \( n \) for fixed \( p \).

\( R \) will be convex in \( n \) for fixed \( p \) and \( n \in [n_1, n_2] \) when either \( w \leq v \), or \( F_X(k_\beta(n) - 1) \geq 1 - \alpha \). We have shown in proposition 1 that for \( w \leq v \) any \((n, p)\) satisfies (3.3') and (3.4') with \( k_\alpha = n \) and \( k_\beta = 0 \). For \( w > v \), \( F_X(k_\beta(n) - 1) \geq 1 - \alpha \) is easily satisfied for
large $\alpha$ or large $p_b$. For if in some interval $[n_1, n_2]$ $k_p(n)$ is
constant then as $n$ increases $F_x(k_\alpha)$ decreases, but for large $\alpha$
$F_x(k_\alpha) \geq 1 - \alpha$ can hold, and for large $p_b$, $F_x(k_p - 1) \geq 1 - \alpha$ can hold.
Of course, if as $n$ increases, $k_p(n)$ increases, then
$F_x(k_p(n) - 1) \geq 1 - \alpha$ can hold.

In the constrained maximization of $\psi(n, p)$ it will generally be
possible to use a search or programming algorithm to determine $p_o(n)$,
but a search method will not always determine $n_o$. In the latter case
$\psi(n, p_o(n))$ must be evaluated at every $n \in \mathbb{R}$.

Maximization of $T_c(n, p)$

Assume that by search or iteration we have determined $\psi(n_o, p_o)$.
This is an upper bound on $T_c(n, p)$ since $g(p) \geq h(p)$ implies
$G(n, p) \geq h(p)$, or

$$T_c(n, p) = NV(p) - NG(n, p) - nc \leq NV(p) - Nh(p) - nc = \psi(n, p).$$

The upper bound can be attained when equality holds in (3.20) at $(n_o, p_o)$.
In maximizing $T_c(n, p)$ any $(n, p)$ yielding $\psi(n_o, p_o)$ is a global
optimum. The conditions stated in Proposition 2 and the assumption that
the interval of interest is convex in $n$ insure only that $G(n, p)$ is
defined on some regular region. Without further restrictions on $G(n, p)$
we cannot insure that $T_c(n^*, p^*)$ determined by any search or nonlinear
programming method is globally optimal. We briefly summarize the struc-
ture of $R$ and then investigate the behavior of $G(n, p)$ in $r$ and $p$.

Assume that the conditions of Proposition 2 are satisfied and that
we need only remark on constraints (3.3) and (3.4). There is a minimum
n, p, in \( \mathbb{R} \), such that for \( n \geq n_p = 0 \) satisfies (3.3) and (3.4), and for each \( n \in \mathbb{R} \) there is a maximum \( p, \bar{p}(n) \) satisfying (3.3), (3.4). Although this \( \bar{p}(n) \) is not generally monotonically increasing in \( n \), there may be intervals \( [n_1, n_2] \) in which \( R \) is convex in \( n \) and \( p \). Proposition 3 states conditions under which \( G(n, p) \) decreases in \( n \) on such intervals.

**Lemma 1**

If \( k_\alpha > [np] \) and \( k_\alpha(n+1) = k_\alpha(n) \), \( \Delta n \sum_0^{k_\alpha} d_x \leq 0 \)

**Proof:**

\[
\Delta n \sum_0^{k_\alpha} d_x = \frac{-(n+1-k_\alpha)}{q} \pi_{k_\alpha}(n+1, p) + \frac{(n-k_\alpha)}{q} \pi_{k_\alpha}(n, p)
\]

\[
= \frac{\pi_{k_\alpha}(n, p)}{q} \left[ (np-k_\alpha) + (p-1) \right]
\]

\[
= (-1) \frac{\pi_{k_\alpha}(n, p)}{q} \left[ (k_\alpha-np) + (1-p) \right] \leq 0 .
\]

**Proposition 3**

Let \( P_1 = \sum_0^{k_\alpha} \pi_x(n, p) \quad D_1 = \sum_0^{k_\alpha} d_x(n, p) \)

\( P_2 = \sum_0^{k_\alpha} \pi_x(n+1, p) \quad D_2 = \sum_0^{k_\alpha} d_x(n+1, p) \)

Assume that for \((n, p), (n+1, p)\), we have

\[
(3.22) \quad G(n, p) = L_1 \sum_0^{k_\alpha(n)} \pi_x(n, p) + J_1 \sum_n^{k_\alpha(n)+1} \pi_x(n, p) .
\]
(3.23) \( G(n+1, p) = L_2 \sum_o \kappa(n+1) x(n+1, p) + J_2 \sum_o^{n+1} \kappa(n+1) x(n+1, p) \).

Let \( \kappa(n) = \kappa(n+1) \), \( \kappa > [np] \), and \( \epsilon = h'(p) \left[ \frac{P_1}{D_1} - \frac{P_2}{D_2} \right] \).

Then if \( J_2 \leq J_1 + \epsilon \), \( G(n+1, p) \leq G(n, p) \).

Proof:

By (3.2') \( L_1D_1 + J_1(-D_1) = h'(p) \) or \( (L_1-J_1) = \frac{h'(p)}{D_1} \), and similarly \( (L_2-J_2) = \frac{h'(p)}{D_2} \). Therefore by Lemma 1 \( (L_2-J_2) \leq (L_1-J_1) \).

Since \( 0 \leq P_2 \leq P_1 \), we have \( (L_2-J_2) P_2 \leq (L_1-J_1) P_1 \) or \( h'(p) \left[ \frac{P_1}{D_1} - \frac{P_2}{D_2} \right] \geq 0 \). By hypothesis, \( J_2 \leq J_1 + h'(p) \left[ \frac{P_1}{D_1} - \frac{P_2}{D_2} \right] \) and we obtain \( J_2 + (L_2-J_2) P_2 \leq J_1 + (L_1-J_1) P_1 \), which by (3.22), (3.23) is \( G(n+1, p) \leq G(n, p) \).

Behavior of \( G(n, p) \) in \( p \)

The behavior of \( G(n, p) \) in \( p \) depends on the changes in \( \frac{\sum_j \phi_x(n, p)}{\sum_j \phi_\lambda(n, p)} \) relative to changes in \( \frac{h(p)}{h'(p)} \) as \( p \) varies. If at some \( p_0 \), \( G(n, p_0) = h(p_0) \) and for \( p < p_0 \), \( h''(p) > 0 \), then as \( p \) decreases \( h'(p) \) becomes negative rapidly. If in the same interval \( \sum_j \phi_x(n, p) \) becomes more negative less rapidly so that \( \phi_x(x \leq j) \) increases in order that \( \sum \phi_1 \phi_1 \) becomes more negative, \( \sum \phi_1 \phi_1 \) increases and may yield \( \sum \phi_1 \phi_1 > h(p) \) for \( p < p_0 \). In Proposition 4 we indicate conditions for which \( G(n, p) - h(p) \) increases as \( p \) decreases.
Lemma 2

If \( k_\alpha > [np] \), and \( k_\alpha(p) \) is constant in some interval \( I \), then
\[
\frac{d}{dp} \sum_o k_\alpha d_x(n,p) \leq 0 \quad \text{for} \quad p \in I.
\]

Proof:

Let \( p_1, p_2 \in I \), \( p_1 < p_2 \), and \( k_\alpha(p_1) = k_\alpha(p_2) \). Then
\[
\sum_o k_\alpha d_x(n,p) = -\frac{n!}{(k_\alpha)! (n-k_\alpha-1)!} \cdot \frac{n-k_\alpha-1}{q} \cdot \frac{k_\alpha}{p} \cdot \frac{\pi_{k_\alpha}(n,p)}{q}.
\]

Since \( d_x(n,p) > 0 \) for \( x > [np] \), \( \pi_{k_\alpha}(n,p_2) > \pi_{k_\alpha}(n,p_1) \). Therefore
\[
\frac{\pi_{k_\alpha}(n,p_2)}{q_2} > \frac{\pi_{k_\alpha}(n,p_1)}{q_1} \quad \text{and} \quad \frac{-n-k_\alpha}{q_2} \pi_{k_\alpha}(n,p_2) < \frac{-n-k_\alpha}{q_1} \pi_{k_\alpha}(n,p_1) \quad \text{or}
\]
\[
\sum_o k_\alpha d_x(n,p_2) \leq \sum_o k_\alpha d_x(n,p_1).
\]

Proposition 4

Let \( I \) be an interval in which the conclusions of Lemma 2 are valid and \( \frac{d}{dp} \sum_o k_\alpha d_x(n,p) \leq 0 \). For \( p_1 < p_2 \) let \( G(n,p) \) be given by

\[
(3.24) \quad G(n,p_1) = L_1 \sum_o k_\alpha \pi_x(n,p_1) + J_1
\]

\[
(3.25) \quad G(n,p_2) = L_2 \sum_o k_\alpha \pi_x(n,p_2) + J_2
\]

Let \( p_1 = \sum_o k_\alpha \pi_x(n,p_1) \quad p_2 = \sum_o k_\alpha \pi_x(n,p_2) \quad p_1 = \sum_o k_\alpha d_x(n,p_1) \quad p_2 = \sum_o k_\alpha d_x(n,p_2) \)
Then if \( J_1 \geq J_2 \) and
\[
\begin{vmatrix}
  h(p_1) & p_1 \\
  h'(p_1) & D_1
\end{vmatrix} > \begin{vmatrix}
  h(p_2) & p_2 \\
  h'(p_2) & D_2
\end{vmatrix},
\]
\[G(n,p_1) - h(p_1) > G(n,p_2) - h(p_2).\]

Proof:

By hypothesis \( D_2 h(p_2) - p_2 h'(p_2) < D_1 h(p_1) - p_1 h'(p_1) \), and by Lemma 2 \( D_2 < D_1 \) or \(-D_2 > -D_1\). Therefore
\[
\frac{D_2 h(p_2) - p_2 h'(p_2)}{D_2} > \frac{D_1 h(p_1) - p_1 h'(p_1)}{D_1}
\]
or
\[
\frac{h(p_2) - h'(p_2)}{D_2} > \frac{h(p_1) - h'(p_1)}{D_1}.
\]

By (3.2') \( \frac{h'(p_2)}{D_2} = L_2 - J_2 \) and \( \frac{h'(p_1)}{D_1} = L_1 - J_1 \). Thus
\[(L_2 - J_2) p_2 - h(p_2) < (L_1 - J_1) p_1 - h(p_1).\] By hypothesis \( J_1 \geq J_2 \) and therefore \( G(n,p_2) - h(p_2) < G(n,p_1) - h(p_1) \) completing the proof.

Renegotiating Parameters

We have been concerned with stating the conditions under which \((n,p)\) can be found to maximize the consumer profit, \( T_c(n,p) \). Beyond stating conditions that yield \( G(n,p) = h(p) \) in some region, there is little we can say about \( G(n,p) \) that aids in negotiation of the parameters \( M, m, v, \alpha, \beta \). Of course as the constraints (3.3)-(3.6) are relaxed the set of points on which \( G(n,p) = h(p) \) increases, but conversely as the constraints become increasingly restrictive the resultant \( G(n,p) \) may become defined on a collection of unconnected
sets. In this event we could resort to a random search procedure to determine \((n^*,p^*)\),\(^9\) but it seems artificial in light of the original procurement situation to stop with such a solution.

The existence of a solution \((n^*,p^*) \neq (n_o,p_o)\) and such that \(T_c(n^*,p^*) < \psi(n_o,p_o)\) can occur in two ways. The discrimination required between \(p_o\) and \(p_b\) may be so fine that it is more profitable for the consumer to insist on higher quality to allow discrimination at a smaller sample size. Or the bounds \(w\) and \(m\) may be so high that in order to motivate production at \(p_o\), the expected payment is greater than \(h(p_o)\). Even if \((n^*,p^*)\) is correct for fixed parameters, the consumer might ask how far he would have to relax \(M, v,\) or \(\beta\), to obtain a policy closer to \((n_o,p_o)\). Whether such relaxation were worth the incremental profit would depend on the consumer's risk preferences. The consumer might also find it necessary and profitable to offer the producer a higher unit profit in return for agreement on decreased \(m, \alpha,\) and \(w\).

In the following example a minimum price of 60 per batch yielded a consumer profit of 350.60, and a minimum price of -51.75 per batch.

\(^9\)In Brooks' random search method we would choose, say, 30 points \((n,p)\) at random in the region \(R\), and evaluate \(T_c(n,p)\) at each point. Assume we are interested in the event that the greatest \(T_c\) thus obtained actually ranks among the top 10\% of all points \((n,p)\) in \(R\). The probability of this event is \(1 - (1-.1)^{30}\) or .958. This method makes no assumptions about the concavity of \(T_c(n,p)\) or the convexity of \(R\). Further refinement could be obtained by choosing points about the maxima obtained, or initiating a gradient procedure at the highest points obtained by random search. [47]
yielded a profit of 351.65. The consumer might offer the producer additional profit up to 1.05 per batch in return for the decreased lower bound.

Let $N = 500$, $\alpha = .1$, $w = h(p)$, $M = 1$, $c = .2$, $p_b = .2$, $p = .4$, $v = .3986$, $m = .12$, $V(p) = (1-p)$, and 

$$h(p) = .154 + .22p + .066e^{-0.02/p}$$

The solution with $m = .12$ is found at $n^* = 18$ and $p^* = .0528$.

$$\varphi_0 = .2918$$
$$\varphi_1 = .2171$$
$$\varphi_2 = .2171$$
$$\varphi_3 = .12$$
$$\vdots$$
$$\varphi_{18} = .12$$

Figure 1

The solution with $m = -.1035$ is found at $n^* = 13$ and $p^* = .0528$.

$$\varphi_0 = .2788$$
$$\varphi_1 = .2178$$
$$\varphi_2 = .2178$$
$$\varphi_3 = -.1035$$
$$\vdots$$
$$\varphi_{13} = -.1035$$

Figure 2
While the original intent in developing this basic model was the introduction of bounds on prices and assignable protection against sampling variation, we do not insist that parameter values be stated only once, but suggest that during negotiation the parameters of the model will be successively adjusted and the unit profit varied, until when the contract is concluded, \((n^*, p^*) = (n_o, p_o)\).

**Computational Procedure**

In order to successively adjust parameters, we must know the value of \(T_c(n^*, p^*)\) and \(\psi(n_o, p_o)\). In experimenting with representative data we found the conditions of Proposition 2 (convexity of \(R\) in \(p\) for fixed \(n\), with \(G(n, p) = h(p)\)) satisfied for a wide range of parameters. We found \(R\) convex in \(n\) for fixed \(p\) less often, and concluded that while the \(T_c(n, p)\) surface had too much structure to justify a random search, it did not have enough for a gradient method. We therefore used a generalization of the Fibonacci search technique.

The Fibonacci search in a single variable is fully developed by Wilde [48], and computational experience with the extension to more than one variable is treated by Krolak and Cooper.\(^{10}\) [49] The generalization

\(^{10}\)A Fibonacci search in a single variable using \(n\) search points can be described as follows. The Fibonacci sequence is \(f_0 = 1, f_1 = 1, f_n = f_{n-1} + f_{n-2}\) \((n \geq 2)\). For unimodal \(\psi(p)\) defined on \(a_1 \leq p \leq b_1\), let \(I_1 = b_1 - a_1\) and define \(\Delta_2\) to be \((f_{n-2}/f_n) I_1\). Let \(p_1 = \Delta_2 + a_1, p_2 = b_1 - \Delta_2\). Then if \(\psi(p_1) \geq \psi(p_2)\), \(a_2 \leftarrow a_1, b_2 \leftarrow p_2\), and \(p_3 = a_2 + \Delta_3\) where \(\Delta_3 = (f_{n-3}/f_{n-1}) I_2\). If \(\psi(p_2) > \psi(p_1)\) then \(a_2 \leftarrow p_1, b_2 \leftarrow b_1, \) and \(p_3 = b_2 - \Delta_3\). This continues until the maximum lies in an interval of length \(I_n = I_1(f_0/f_n)\).
to two variables requires that when searching in the variable \( n \), \( T_c(n,p) \) is evaluated at \( p^*(n) \), where \( p^*(n) \) maximizes \( T_c(n,p) \) for fixed \( n \). In order to avoid determining the region \( R \), we modify constraints (3.3'-3.5') by introducing artificial variables as follows.

\[
\begin{align*}
(3.3^0) & \quad \varphi_{k_1} + \theta_1 \geq w \\
(3.4^0) & \quad \varphi_{k_2} + \theta_2 \leq v \\
(3.5^0) & \quad \begin{cases} 
\varphi_0 + \theta_3 \leq M \\
\varphi_n + \theta_4 \geq m 
\end{cases}
\end{align*}
\]

The \( \theta_i \) are unrestricted in sign, and we change (3.1') to

\[
(3.1^0) \quad \text{Minimize } \sum_{i=0}^{n} \mu_i \varphi_i + Q \sum_{i=1}^{4} |\theta_i|
\]

where \( Q \) is a very large positive number. This problem will be feasible at all points of some conveniently defined region, and points feasible for the original problem (3.1')-(3.6') will be feasible for the modified problem with all \( \theta_i = 0 \).

The generalized Fibonacci search is computationally efficient.\(^{11}\)

Furthermore it provides a grid of the points evaluated which serves as a

\(^{11}\)The procedure was programmed in Fortran II for the IBM 7090 at Stanford University. The number of points evaluated is optional, and no attempt was made to optimize the program. Using a ten point search over the sample sizes and a ten point search over the interval \((0,p_b)\) the procedure may be viewed as 100 sequential evaluations to determine the maximum of a function on a \(143 \times 89\) lattice. Solutions almost always require less than 30 seconds.
check both for multimodality and for an irregular search region. If examination of the points evaluated suggests that the search did not terminate with an optimal solution, a more detailed evaluation may be necessary, and is possible by redefining the search region. The computational procedure can be succinctly written as

$$\max_{n} \max_{p} \left\{ NV(p) - nc - \min_{\varphi} \sum_{\varphi} \pi_1(n,p) \varphi_1 \right\}$$

We actually increase efficiency by proceeding as follows:

1. Find $p_0$ to maximize $V(p) - h(p)$.

2. After specifying a lower bound and an arbitrary number of search points, begin Fibonacci search in the variable $n$.

3. At each $n$ proceed directly to $p_0$ and evaluate $G(n,p_0)$ by solving the linear programming problem (3.1"')(3.6"'), as modified in (3.1')(3.3)(3.5).  
   a. If $G(n,p_0) = h(p_0)$, $p^*(n) = p_0$.
   b. If $G(n,p_0) > h(p_0)$, begin a Fibonacci search in the variable $p$ to determine $p^*(n)$.

For cases in which $G(n,p)$ is unimodal in $p$ and decreasing in $n$, or in which $G(n,p_0) = h(p_0)$ for some interval $[n_1, n_2]$, this procedure terminates with an optimal $(n,p)$. This will be the case when the conditions of Proposition 1 are satisfied on $[n_1, n_2]$. Otherwise the procedure provides no assurance that the global maximum will be attained, but due to the possible nonconvexity of $R$ and nonconcavity of $T_c(n,p)$, no procedure short of evaluating every sample size in $[n, \bar{n}]$ can promise a global maximum with certainty.
The following example illustrates points discussed during development of the basic model.

Example 1

Let \( V(p) = 1 - 3p, \ p \leq p_b \)

\[
= 0 \quad p > p_b
\]

\( h(p) = 0.4 - 0.2p - 0.01e^{-0.05/p} \),

\( c = 0.172, \ N = 200, \ M = 1, \ m = 0.15 \),

\( p_b = 0.175, \ v = 0.45V(p), \ \beta = 0.10 \),

\( w = 0.75h(p), \ \alpha = 0.05 \)

The optimal contingent pricing policy is found at \( n^* = 39 \), with \( p^* = 0.0325 \), and is shown in Figure 3.

\[\begin{align*}
\Phi_x &= 0.5243 \quad x = 0, 1 \\
\Phi_x &= 0.3580 \quad x = 2, 3 \\
\Phi_x &= 0.2137 \quad x = 4, \ldots, 7 \\
\Phi_x &= 0.15 \quad x = 8, \ldots, 39
\end{align*}\]

![Figure 3](image)

In Figure 4 and Table 1, the expected payment and profit resulting from the policy \((n,\Phi)\) is displayed.
Figure 4

Table 1

<table>
<thead>
<tr>
<th>$p$</th>
<th>Value at $p$</th>
<th>Cost at $p$</th>
<th>Expected Payment at $p$</th>
<th>Producer Profit at $p$</th>
<th>Consumer Profit at $p$</th>
</tr>
</thead>
<tbody>
<tr>
<td>.010</td>
<td>.921</td>
<td>2.4758</td>
<td>0.514</td>
<td>-1.961</td>
<td>0.4069</td>
</tr>
<tr>
<td>.015</td>
<td>.906</td>
<td>0.7894</td>
<td>0.504</td>
<td>-0.285</td>
<td>0.4019</td>
</tr>
<tr>
<td>.020</td>
<td>.891</td>
<td>0.5666</td>
<td>0.492</td>
<td>-0.074</td>
<td>0.3988</td>
</tr>
<tr>
<td>.025</td>
<td>.876</td>
<td>0.4984</td>
<td>0.479</td>
<td>-0.019</td>
<td>0.3969</td>
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In Figure 5 a grid of $T_c(n,p)$ in the region surrounding $(n^*,p^*)$ is shown, we denote infeasibility by *.
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</tr>
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</table>

Proportion Defective, \( p \)

Figure 5

Figure 6 indicates the increase in \( T_c(n^*, p^*) \) with increasing \( \alpha \).

![Graph showing increase in Tc(n*, p*) with increasing alpha](image)

Figure 6

**Piecewise Linear Price Schedules**

Existing incentive price schedules are usually piecewise linear.

In the simplest and most usual case incentive payments are determined by negotiating a sample size, maximum and minimum payments, and the number of defectives at which the maximum and minimum payment occur. The linear pricing policy shown in Figure 7 can be written as follows.
\[\phi_x = \phi_0 \quad 0 \leq x \leq a\]
\[\phi_x = \phi_0 \left[\frac{b - x}{b - a}\right] + \phi_n \left[\frac{x - a}{b - a}\right] \quad a < x < b\]
\[\phi_x = \phi_n \quad b \leq x \leq n\]

Number of Defectives Observed

Figure 7

We will modify the basic model to derive linear policies maximizing consumer net gain subject to the general assumptions 1-13. Assume that the consumer chooses the sample size, \(n\), and the pricing policy, \((a, b, \phi(x))\) to maximize (3.1) subject to (3.3), (3.4), (3.6), and (3.26), knowing that the producer will choose \(p'\) to maximize (3.2).

\[
\begin{align*}
\phi_0 &\leq M \\
\phi_x &= \phi_0 \quad x = 0, 1, \ldots, a \\
\phi_x &= \phi_0 \left[\frac{b - x}{b - a}\right] + \phi_n \left[\frac{x - a}{b - a}\right] \quad x = a + 1, \ldots, b - 1 \\
\phi_x &= \phi_n \quad x = b, \ldots, n \\
\phi_n &\geq m
\end{align*}
\]

(3.26)

For fixed \(n\), \(p_o\), \((0 < p_o < 1)\), a set of prices, \(\phi_x\), maximizing (3.1) while simultaneously maximizing (3.2) at \(p_o\) and satisfying (3.3), (3.4), (3.6), and (3.26) are solutions to the linear programming problem (3.1')-(3.6') with (3.5') replaced by (3.26). The justification for
this statement is identical to the proof of Theorem 1 and is not repeated.

The transformation (3.10) reduces the linear programming problem to one containing only 6 constraints.

\[(3.27)\quad \text{Minimize } \sum_{i} p_{0}^{i} y_{i}\]
\[(3.28)\quad \sum_{i} p_{0}^{i} y_{i} = h'(p_{0})\]
\[(3.29)\quad \sum_{\max(a, k_{a})}^{n} y_{i} \geq w - m\]
\[(3.30)\quad \sum_{\max(a, k_{b})}^{n} y_{i} \leq v - m\]
\[(3.31)\quad \begin{cases} 
\sum_{a}^{b-l} y_{i} \leq M - m \\
\left[\frac{b-a}{b-x}\right] \sum_{x}^{b-l} y_{i} = \sum_{a}^{b-l} y_{i}
\end{cases}\]
\[(3.32)\quad \sum_{a}^{b-l} p_{0}^{i} y_{i} \geq h(p_{0}) - m\]

We use the following notation in discussing linear pricing policies.

\[(3.33)\quad p_{0} = \sum_{x}^{o} \pi_{x} + \sum_{a+l}^{b-l} \left(\frac{b-x}{b-a}\right) \pi_{x}\]
\[(3.34)\quad p_{n} = \sum_{x}^{n} \pi_{x} + \sum_{a+l}^{b-l} \left(\frac{x-a}{b-a}\right) \pi_{x}\]
\[(3.35)\quad d_{0} = \sum_{x}^{o} d_{x} + \sum_{a+l}^{b-l} \left(\frac{b-x}{a+1}\right) d_{x}\]
\[ D_n = \sum_{b}^{n} d_x + \sum_{a+1}^{b-1} \left( \frac{x-a}{b-a} \right) d_x. \]

Let \( p_o \) maximize \( V(p) - h(p) \). For fixed \( a \) and \( b \), there do not in general exist \( \varphi_o \) and \( \varphi_n \) that maximize \( T_c \) at \( p_o \). Relations (3.1) and (3.2) can be written

\[ \varphi_o + \varphi_n \geq h(p_o) \quad \text{and} \]

\[ D_o \varphi_o + D_n \varphi_n = h'(p_o). \]

Assume that with equality in both equations, the solution is \( \varphi_o \).
\( \varphi_n \). If \( \varphi_n < m \), the solution is not feasible by (3.26), and hence \( \sum_1 \varphi_1 > h(p_o) \). Therefore the possibility again exists of an optimal solution at \( p^* \neq p_o \), and as in the basic case we search for \( p^*(n) \) and \( n^* \) with a generalized Fibonacci search procedure. For fixed \( (n,p) \) we find \( a^*(n,p) \), and \( b^*(a,n,p) \) by iteration after first restricting the number of points to be evaluated. We will develop the reduction procedure after presenting an example of a linear pricing policy derived by this method.

**Example 2**

Let \( V(p) = 1 - 1.5p \), \( p \leq .17 \)

\[ = 0 \quad \quad \quad \quad p > .17 \]

\( h(p) = .25 - .415p - .06 \ln p \)

\( c = .12, \ N = 500, \ M = .45, \ m = .10, \)

\( p_b = .17, \ v = .4V(p), \ \beta = .25, \)

\( w = h(p), \ \alpha = .25. \)

\(^{12}\) If \( b = a + 1 \) the second summation is deleted.
The optimal linear pricing policy is found at \( n^* = 46, \ p^* = .052 \), with \( a = 3, \ b = 7, \ \varphi_0 = .4209 \), and \( \varphi_{46} = .11 \). This is shown in Figure 8.

In Figure 9 and Table 2 we indicate the expected payments, costs and profits resulting from this policy.
Determining Feasible \((a, b)\)

For fixed \(n\) and \(p\) \(a\) may lie in the interval \([0, n-1]\), and \(b\) in the interval \([a+1, n]\), yielding \(n(n+1)/2\) points. We use constraints (3.3) and (3.4) to initially reduce the number of points. Assume that \(w > v\), and that \(n\) is sufficiently large so that \(p = 0\) satisfies (3.3) and (3.4), and let \(\bar{p}(n)\) be the greatest \(p\) satisfying (3.3), (3.4). Therefore for all \(p \in [0, \bar{p}(n)]\), \(k_\alpha(p) < k_B\). Now if \(a \geq k_B\), we would have \(v \geq k_B = \varphi_{k_B} \geq w\), a contradiction. Thus we need never consider \(a > k_B - 1\). Similarly if \(a < k_\alpha\) then if \(b < k_\alpha\) we would have \(w \leq \varphi_{k_\alpha} = \varphi_b = \varphi_{k_B} \leq v\) with a contradiction as before. Therefore for \(a < k_\alpha\) we need only consider \(b > k_\alpha\).

For fixed \(a\) we obtain bounds on feasible \(b\). For \(a < k_\alpha\) when \(\varphi_a = M\), we must have \(\varphi_{k_\alpha} \geq w\). Let \(s = \frac{M - w}{a - k_\alpha}\). Then \(b_{\min}\) is given
by the smallest integer greater than or equal to \( \frac{Y - W}{s} + k_\alpha \), as shown in Figure 10. If \( \varphi_{k_\alpha} = w \), \( b \) cannot exceed some \( b_{\text{max}} \) which would yield \( \varphi_{k_\beta} > v \) as shown in Figure 11. Therefore \( b_{\text{max}} \) is given by the greatest integer less than \( k_\beta + \frac{m - v}{s} \), where \( s = \frac{v - w}{k_\beta - k_\alpha} \).

If \( b_{\text{max}} < b_{\text{min}} \) there exist no prices for the current \( a' \) that satisfy (3.2'), (3.4'), (3.26), and (3.6'). Furthermore all \( a < a' \) are infeasible since \( b_{\text{max}} \) is constant in \( a \) but \( b_{\text{min}} \) increases as \( a \) decreases. For \( k_\alpha \leq a \leq k_\beta - 1 \) we see \( b_{\text{min}} = a + 1 \). Because \( \varphi_{k_\beta} \geq w \) and \( \varphi_{k_\beta} \leq v \), \( b_{\text{max}} \) is the greatest integer less than \( a + \frac{m - w}{s} \), where \( s = \frac{w - v}{a - k_\beta} \). See Figure 12.
Thus rather than \( n(n+1)/2 \) points we consider only \( a \in [0, k_B - 1], b \in [b_{\text{min}}, b_{\text{max}}] \), and if \( b_{\text{max}} < b_{\text{min}} \) evaluation terminates in \( b \). We further reduce the number of points evaluated by terminating iteration in \( b \) when \( T_c(b) \) first decreases in \( b \), and terminate iterating in \( a \) when \( T_c(a) \) first decreases in \( a \). This procedure is appropriate when \( T_c(a, b) \) is unimodal in \( a \) and \( b \). If \( T_c(a, b) \) is not unimodal in either variable this terminating rule is certain of attaining only a local maximum. If the value of \( T_c(a, b) \) at the local maximum is equal to \( \psi(n, p) \), then this maximum is global over \( (a, b) \), regardless of the unimodality of \( T_c(a, b) \).

The relatively few \( (a, b) \) remaining after the reduction process makes numerical determination of linear pricing policies practicable. Computation time is longer than in the basic model due to the search for \( a^* \) and \( b^* \) for each \( (n, p) \). Solutions for a "typical" problem involving 50 search points require about 40 seconds on the IBM 7090.
and use the resulting sample size as a lower bound in the linear model, reducing the number of sample sizes to be searched.

**Unimodality of** $T_c(a,b)$

We state the conditions in Proposition 5 under which $G(n,p,a,b) = h(p)$ for fixed parameters $M$, $m$, $w$, $v$. Alternatively these are the conditions on the parameters so that $G(n,p,a,b) = h(p)$ for fixed $a$, $b$. The proof is simply a restatement of (3.1')-(3.6') and will be omitted.

**Proposition 5**

Let $\mathcal{R}' = \begin{pmatrix} P_0 & P_n \\ D_0 & D_n \end{pmatrix}$, $\mathcal{R}_n = \begin{pmatrix} h(p) & P_n \\ h'(p) & D_n \end{pmatrix}$, $\mathcal{R}'_n = \begin{pmatrix} P_0 & h(p) \\ D_0 & h'(p) \end{pmatrix}$.

Then $G(n,p,a,b) = h(p)$ if and only if (3.39)-(3.42) hold.

(3.39) \[ \frac{\mathcal{R}'_n}{\mathcal{R}} \geq m \]

(3.40) \[ \frac{\mathcal{R}_n}{\mathcal{R}} \leq M \]

(3.41) \[ \mathcal{R}'_n [b - k_\alpha] + \mathcal{R}_n [k_\alpha - a] \geq \nu \mathcal{R}(b-a) \]

(3.42) \[ \mathcal{R}'_n [b - k_\alpha] + \mathcal{R}_n [k_\beta - a] \leq \nu \mathcal{R}(b-a) \]

Whenever $G(n,p) = h(p)$ for $(a,b)$ in some $A \times B$, and the problem is infeasible outside that region, the iterative procedure will always select the smallest $a$ and largest $b$. It is possible for $T_c(a,b)$ to be unimodal in $a$ and $b$, and in the next two
propositions we indicate conditions sufficient for \( T_c(a, b) \) to be unimodal in \( a \). Analogous conditions can be derived for \( b \).

**Lemma 3**

Let \( \Delta_a \) represent the first difference in \( a \) and let \( P_0', P_n', D_0, D_n \) be as defined in (3.35)-(3.36). Then \( \Delta_a P_0 > 0, \Delta_a P_n < 0 \), and if \( a + 1 > np \), \( \Delta_a D_0 > 0 \) and \( \Delta_a D_n < 0 \). This statement is also true if \( \Delta_a \) is replaced by \( \Delta_b \).

**Proof:**

\[
d_x = \begin{pmatrix} n \\ x \end{pmatrix} p^{x-1}(1-p)^{n-x-1}(x-np). \text{ Therefore if } a + 1 \geq np \text{ then } \\
\frac{d}{da}a + 1 \geq 0.
\]

\[
\Delta_a D_0 = \frac{1}{b-a} \sum b^{-1} d_x \frac{b-x}{(a-b)(a-b+1)}.
\]

Both terms are positive. Therefore \( \Delta_a D_0 > 0 \). The remaining statements of the lemma follow in the same manner by writing out the definitions of the quantities.

**Proposition 6**

Denote the dependence of \( \varphi_o, P_0, D_0, \varphi_n, P_n, D_n \) on \( a \) as \( \varphi_o^a, P_0^a, \) etc., and abbreviate \( \Delta_a \) by \( \Delta \). If for some \( a' > np + 1 \), \( \varphi_o^{a'} = w \), and if for \( np + 1 \leq a < a' \) the solution, \( \varphi_o^a \), to (3.37), (3.38) is increasing in \( a \), then \( G(a-1) > G(a) \) for \( np + 1 \leq a < a' \).

**Proof:**

By hypothesis \( \varphi_o^{a'} = w \). Either the solution to the equations (3.37), (3.38) yields \( w \), or the solution yields \( \varphi_o' < w \) and \( \varphi_o^{a'} = w \) by (3.3') and (3.26). In either case since the solution, \( \varphi_o' \), is
increasing in $a$, $\varphi^{a-1}_o = w$. Therefore, for every $a$ in the interval of interest, $\varphi^{a}_o = w$.

At $a$ we have $h'(p) = wD_o^n + \varphi^{a-1}_n n^n$, and an analogous expression at $a - 1$. Therefore

$$\varphi^{a-1}_n = \frac{wD_o^n + \varphi^{a-1}_n n^n}{D^{a-1}_n}, \quad \text{and}$$

$$G(a-1) = wP^{a-1}_o + \frac{P^{a-1}_n}{D^{a-1}_n} \left[ wD_o^n + \varphi^{a-1}_n n^n \right], \quad \text{from which}$$

$$G(a-1) - G(a) = P^{a-1}_n \left\{ w \left[ \frac{\Delta P_o}{P^{a-1}_n} + \frac{\Delta D_o}{P^{a-1}_n} \right] + \varphi^{a-1}_n \left[ \frac{D^{a}_n}{D^{a-1}_n} - \frac{P^{a}_n}{P^{a-1}_n} \right] \right\}.$$

We will show the quantity in brackets is positive, proving the proposition. By (3.4') $\varphi^{a-1}_n \leq w$, and it suffices therefore to show

$$\frac{\Delta D_o}{D^{a-1}_n} - \frac{P^{a}_n}{P^{a-1}_n} \leq \frac{\Delta P_o}{P^{a-1}_n} + \frac{\Delta D_o}{D^{a-1}_n}.$$  This is equivalent to

$$\frac{\Delta D_o - D^{a}_n}{D^{a-1}_n} \leq -\frac{\Delta P_o - P^{a}_n}{P^{a-1}_n}, \quad \text{which is} \quad \frac{D^{a}_o - D^{a-1}_o}{D^{a}_n} \geq \frac{D^{a}_o - D^{a-1}_o}{P^{a}_n} > -1, \quad \text{or} \quad \Delta D_o > \Delta a n.$$

But by Lemma 3 $\Delta D_o > 0$ and $\Delta a n < 0$.

**Proposition 7**

Assume $p \neq 0$ or 1, and for some $a'' \geq np + 1$, $\varphi^{a''}_n = m$. Assume also that for $a > a''$ the solution, $\varphi'_n$, to the equations (3.37), (3.38) decreases in $a$. Then $G(a) < G(a+1)$ for $a > a''$.  

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Proof:

Case I: \( m \leq 0 \). By (3.2) \( \varphi_o^a D_o^a + mD_n^a = \varphi_o^{a+1} D_o^{a+1} + mD_n^{a+1} \) and

\[
\varphi_o^a = \frac{\varphi_o^{a+1} D_o^{a+1}}{D_o^a} + \frac{mD_n^a}{D_o^a}. \quad \text{From which}
\]

\[
(3.43) \quad \frac{\Delta G(a)}{p_o^a} = \varphi_o^{a+1} \left[ \frac{p_o^{a+1}}{p_o^a} - \frac{D_o^{a+1}}{D_o^a} \right] - m \left[ \frac{\Delta D_n^a}{D_o^a} - \frac{\Delta P_n^a}{p_o^a} \right].
\]

By Lemma 3, \( \Delta D_n^a > 0 \), and thus \( \frac{p_o^{a+1}}{p_o^a} > 1 \). \( \Delta D_o^a > 0 \) and \( D_o^{a+1} < 0 \) imply \( \frac{D_o^{a+1}}{D_o^a} < 1 \). Therefore the quantity in the first bracket is positive.

By Lemma 3 \( \Delta D_n^a < 0 \), and \( \Delta P_n^a < 0 \). Since \( D_o^a < 0 \), and \( p_o^a > 0 \), the quantity \( \{-m \left[ \frac{\Delta D_n^a}{D_o^a} - \frac{\Delta P_n^a}{p_o^a} \right]\} \) is positive. Therefore \( \Delta G(a) \) is positive.

Case II: \( m > 0 \). From (3.43) since \( \varphi_o > m \) in any case of interest, it suffices to show \( \frac{p_o^{a+1}}{p_o^a} - \frac{D_o^{a+1}}{D_o^a} > \frac{\Delta D_n^a}{D_o^a} - \frac{\Delta P_n^a}{p_o^a} \). This statement is

\[
\frac{p_o^{a+1}}{p_o^a} > \frac{D_o^{a+1}}{D_o^a} \quad \text{which is} \quad \frac{p_o^a}{p_o^a} > \frac{D_o^a}{D_o^a}, \quad \text{proving Case II.}
\]

If the conditions of Propositions 6 and 7 are satisfied, and if \( a' \leq a'' \) with \( T_c(a) \) constant in \( a \) when \( a' < a < a'' \), then \( T_c(a) \) is unimodal in \( a \).
Once $a^*$ and $b^*$ are determined for fixed $n$ and $p$, the problem of determining $n^*$ and $p^*$ is identical to that in the basic case, except that the linear model is more restrictive. With the contingent pricing problem formulated in this manner, there is no advantage in restricting attention to linear policies. The same parameters must be negotiated in both cases, and the linear model must always require more computational effort and can never result in a higher consumer profit than the basic model.

For the data in Example 2 the basic model yielded the following pricing policy at $n^* = 36$ and $p^* = .0522$.

\[
\begin{align*}
\phi_x &= .4399, \quad x = 0, 1 \\
\phi_x &= .3933, \quad x = 2, 3 \\
\phi_x &= .2285, \quad x = 4, \ldots, 8 \\
\phi_x &= .1100, \quad x = 9, \ldots, 36
\end{align*}
\]

This policy is displayed in Figure 13, and the linear policy is displayed in Figure 14.
The expected payments are displayed in Figure 15 for both policies, and Table 3 contains the expected cost and profits for the nonlinear policy.
Table 3

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<tr>
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<th>Cost at p</th>
<th>Expected Payment at p</th>
<th>Producer Profit at p</th>
<th>Consumer Profit at p</th>
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Consumer unit expected profit using the linear policy is .5170, and using the less restricted price schedule derived by the basic model the expected unit profit is .5194. From Figure 15 it can be noted that the policy at sample size 36 does not provide the same discrimination in expected payment that the policy at n = 46 does, but both policies satisfy the conditions (3.3) and (3.4). In situations where the producer prefers a linear price schedule due to ease of interpretation, the consumer must decide whether to offer increased expected profit in return for producer acceptance of the non-linear price schedule.
Extensions and Applications

We indicate situations in which the basic assumptions 1-13 can be weakened and a more general class of production and procurement situations considered.

Replacement of Discovered Defectives

In most instances in which inspection or testing is non-destructive, a producer must replace all defectives discovered during the inspection process at no additional cost to the consumer. We will modify the basic model to treat this situation and exhibit a pricing policy derived from the data used in Example 1.

Assume that the cost to the producer of replacing discovered defectives is the production cost, \( h(p) \), plus a fixed charge, \( r \), reflecting additional transportation, testing, and special handling costs. The expected net gains are

\[
T_c(n,p) = (N-n) V(p) + n[V(0) - c] - Ng(p),
\]

\[
T_p(n,p) = Ng(p) - Nh(p) - np r.
\]

We will make the constraint (3.3) more realistic by providing the producer some assurance of regaining a portion of expected cost. Therefore the consumer chooses \((n,\varphi)\) to maximize (3.44) knowing that the producer will choose \(p'\) to maximize (3.45), and such that \((n,\varphi)\) satisfy

\[
pr[\varphi(x) \geq Nw(h(p')) + rx | p \leq p'] \geq 1 - \alpha
\]

\[
Ng(p') \geq Nh(p') + np'r
\]

in addition to (3.4) and (3.5).
Solutions, \( \Phi_x \), are obtained by solving the following problem for fixed \( n \) and \( p \).

\[
(3.48) \quad \text{Minimize } \sum_{\varphi}^n \pi_x \Phi_x
\]

subject to

\[
(3.49) \quad \sum_{\alpha}^n a_x \Phi_x = h'(p) + \frac{rn}{N}
\]

\[
(3.50) \quad \Phi_x \leq \left[ \frac{h(p')}{N} + \frac{r\alpha}{N} \right]
\]

\[
(3.51) \quad \sum_{\alpha}^n \pi_x \Phi_x \geq h(p') + \frac{rpm'}{N}
\]

and subject also to \((3.4')\) and \((3.5')\).

Using the data of Example 1 with \( r = 0.3 \), the following price schedule was obtained at \( n^* = 39, \; p^* = 0.0350 \).

\[
\begin{align*}
\Phi_x &= 0.5146 & x &= 0, 1 \\
\Phi_x &= 0.3869 & x &= 2, 3 \\
\Phi_x &= 0.2137 & x &= 4, \ldots, 7 \\
\Phi_x &= 0.15 & x &= 8, \ldots, 39
\end{align*}
\]

This policy is shown in Figure 16.
In Figure 17 and Table 4 the expected payment and profit resulting from this pricing policy are shown.

The consumer expected profit in this situation, .4113, is greater than the expected profit in the basic model, .3951, which certainly agrees with the procurement situation, which is more advantageous to the consumer.
Table 4

<table>
<thead>
<tr>
<th>p</th>
<th>Value at p</th>
<th>Cost at p</th>
<th>Expected Payment at p</th>
<th>Producer Profit at p</th>
<th>Consumer Profit at p</th>
</tr>
</thead>
<tbody>
<tr>
<td>.01</td>
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<td>2.4821</td>
<td>.5070</td>
<td>-.9751</td>
<td>.4201</td>
</tr>
<tr>
<td>.02</td>
<td>.9029</td>
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<td>.897</td>
<td>-. .0791</td>
<td>.4133</td>
</tr>
<tr>
<td>.03</td>
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<td>.4113</td>
</tr>
<tr>
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<td>.4554</td>
<td>.4554</td>
<td>.0000</td>
<td>.4113</td>
</tr>
<tr>
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<td>.4453</td>
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<td>-.0025</td>
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<tr>
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<td>.8305</td>
<td>.4335</td>
<td>.4172</td>
<td>-.0162</td>
<td>.4113</td>
</tr>
<tr>
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<td>.4266</td>
<td>.3916</td>
<td>-.0350</td>
<td>.4147</td>
</tr>
<tr>
<td>.07</td>
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<td>.4220</td>
<td>.3669</td>
<td>-.0551</td>
<td>.4153</td>
</tr>
<tr>
<td>.08</td>
<td>.7580</td>
<td>.4185</td>
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<td>-.0750</td>
<td>.4146</td>
</tr>
<tr>
<td>.09</td>
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<td>-.0939</td>
<td>.4121</td>
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<tr>
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<td>.7097</td>
<td>.4133</td>
<td>.3020</td>
<td>-.1113</td>
<td>.4078</td>
</tr>
</tbody>
</table>

Uncertainty in Process Control

Assume that the producer does not choose \( p \) precisely, but instead chooses a mean quality, \( \mu \), at a known cost, \( h(\mu) \). The distribution of the random variable \( p \) is known to be \( f(p|\mu, \theta_1, ... , \theta_n) \) where the \( \theta_i \) are known, uncontrollable parameters. We suppose that in this situation the consumer selects \( (n, \varphi) \) to maximize (3.52) subject to (3.54)-(3.57) knowing that the producer will select \( \mu \) to maximize (3.53).

(3.52) \( T_c(\mu, \varphi, n) = Nv(\mu) - Ng(\mu) - nc \)

(3.53) \( T_p(\mu, \varphi, n) = Ng(\mu) - Nh(\mu) \)

(3.54) \( \text{pr}[N_{\varphi}(x) \geq Nw|p \leq \mu'] \geq 1 - \alpha \)
\begin{align*}
(3.55) \quad \Pr[N\varphi(x) \leq Nv| p \geq p_b] & \geq 1 - \beta \\
(3.56) \quad M \geq \varphi_0 \geq \cdots \geq \varphi_n \geq m \\
(3.57) \quad Ng(u') \geq Nh(u')
\end{align*}

For fixed $n$, $\mu$, the $\varphi_x$ are solutions to the linear programming problem (3.1')-(3.6') with $h(p)$, $h'(p)$ replaced by $h(\mu)$, $h'(\mu)$, and with

\[\pi_x(n, \mu) = \int_0^1 \Pr[x|n, p] f(p|\mu) \, dp,\]

\[d_x(n, \mu) = \frac{d}{d\mu} \int_0^1 \Pr[x|n, p] f(p|\mu) \, dp,\]

$k_\alpha$ satisfying \( \sum_{j=0}^\infty \frac{n^j}{j!} \mu^j (1-\mu)^{n-j} \), and $k_B$ satisfying \( \sum_{j=k_B}^\infty \frac{n^j}{j!} p_b^j (1-p_b)^{n-j} \).

Assume that from jobs previously attempted by the producer of the items in Example 1 we estimate that the random variable $p$ has a beta distribution, and that while the producer can select the mean, $\mu$, the variance is uncontrollable and is estimated at $36 \times 10^{-6}$. From this model we obtain an optimal policy at $n^* = 39$, and $\mu^* = .0339$. The policy is shown in Figure 18.

\[\varphi_x = .4692 \quad x = 0, 1\]
\[\varphi_x = .4605 \quad x = 2, 3\]
\[\varphi_x = .2138 \quad x = 4, \ldots, 21\]
\[\varphi_x = .1500 \quad x = 22, \ldots, 39\]
In Figure 19 and Table 5 the expected payment and profit resulting from this policy and the basic policy is shown.

A price schedule based on uncertainty in process control does not provide discrimination as fine as that in the basic case.
Table 5

<table>
<thead>
<tr>
<th>$\mu$</th>
<th>Value at $\mu$</th>
<th>Cost at $\mu$</th>
<th>Expected Payment at $\mu$</th>
<th>Producer Profit at $\mu$</th>
<th>Consumer Profit at $\mu$</th>
</tr>
</thead>
<tbody>
<tr>
<td>.0100</td>
<td>.9212</td>
<td>2.4758</td>
<td>.4691</td>
<td>-2.0067</td>
<td>.4521</td>
</tr>
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<td>.0200</td>
<td>.8912</td>
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<td>-.1014</td>
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<tr>
<td>.0300</td>
<td>.8612</td>
<td>.4681</td>
<td>.4574</td>
<td>-.0107</td>
<td>.4038</td>
</tr>
<tr>
<td>.0339</td>
<td>.8495</td>
<td>.4544</td>
<td>.4544</td>
<td>.0000</td>
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</tr>
<tr>
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<td>-.0007</td>
<td>.3910</td>
</tr>
<tr>
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<td>-.0052</td>
<td>.3783</td>
</tr>
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<td>.3681</td>
</tr>
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<td>-.0295</td>
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</tr>
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<td>.6512</td>
<td>.4031</td>
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<td>.3253</td>
</tr>
</tbody>
</table>

Uncertainty in the Cost Function

We next assume that the cost of selecting the mean quality, $\mu$, is not known precisely, but is $h(\mu) + \eta$, where $\eta$ is a random variable with $E(\eta) = 0$, and known distribution. This situation differs from the previous case only in constraint (3.57) which we now write as

(3.58) \[ \Pr[g(\mu') \geq h(\mu') + F^{-1}_\eta(1-\gamma)] \geq 1 - \gamma \]

which is simply $g(\mu') \geq h(\mu') + F^{-1}_\eta(1-\gamma)$, where $F^{-1}_\eta(1-\gamma)$, is the minimum $u$ such that $F(\eta(u)) \geq 1 - \gamma$, given $\mu = \mu'$.

Using the data of Example 1, with the assumption that $\eta$ is normally distributed with $E(\eta) = 0$ and $\sigma_\eta = .0333 h(\mu)$, and that $p$ is beta distributed as in the previous example, we obtain the following policy for $\gamma = .10$. The optimal sample size is 39, and $\mu^* = .0339$. The policy is shown in Figure 20.
$\phi_x = .4869 \quad x = 0, 1$
$\phi_x = .4862 \quad x = 2, 3$
$\phi_x = .2137 \quad x = 4, \ldots, 21$
$\phi_x = .1500 \quad x = 22, \ldots, 39$

In Figure 21 the expected payment resulting from this policy is shown, and Table 6 contains the expected profits.
Table 6

<table>
<thead>
<tr>
<th>μ</th>
<th>Value at μ</th>
<th>Cost at μ</th>
<th>Expected Payment at μ</th>
<th>Producer Profit at μ</th>
<th>Consumer Profit at μ</th>
</tr>
</thead>
<tbody>
<tr>
<td>.01</td>
<td>.921</td>
<td>2.475</td>
<td>.487</td>
<td>-1.9880</td>
<td>.434</td>
</tr>
<tr>
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<td>.566</td>
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<td>- .079</td>
<td>.406</td>
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<tr>
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<td>.861</td>
<td>.468</td>
<td>.479</td>
<td>.011</td>
<td>.382</td>
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<td>.375</td>
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<td>.020</td>
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<td>- .014</td>
<td>.340</td>
</tr>
<tr>
<td>.08</td>
<td>.711</td>
<td>.410</td>
<td>.381</td>
<td>- .029</td>
<td>.330</td>
</tr>
<tr>
<td>.09</td>
<td>.681</td>
<td>.406</td>
<td>.360</td>
<td>- .046</td>
<td>.321</td>
</tr>
<tr>
<td>.10</td>
<td>.651</td>
<td>.403</td>
<td>.337</td>
<td>- .066</td>
<td>.314</td>
</tr>
</tbody>
</table>

The producer expected profit is positive due to the constraint (3.58), as was expected since the price schedule must provide protection against cost overruns.

Applications to Lifetesting

It is apparent that incentive price schedules can be derived from the models developed here in any situation in which items can be classified as defective or non-defective. We will note one such application to the situation where the mean life of an item is of interest.

Assume that the lifetimes, t, of items produced are random variables with known distribution f(t|μ), and that the producer controls μ at a known cost h(μ). Assume also that the value to the consumer of a unit with an expected life of μ can be expressed as V(μ). Several existing reliability incentive plans place a sample of n on test,
record the times until failure, and base the incentive payment on the average time to failure. An alternative approach is to fix the length of time, \( \tau \), for which units will be tested, consider survival to \( \tau \) success, failure prior to \( \tau \), failure, and base payment on the number of failures, \( x \). The problem is identical to the ones we have already treated with the consumer choosing \( n, \mu, \) and \( \phi_x \) to maximize \( NV(\mu) - nc - NG(n, \mu) \), and \( \int_0^\tau f(t|\mu)dt \) taking the place of \( p \). Alternatively the sample size, \( n \), might be predetermined and \( \tau \) selected, where the unit testing cost, \( c(\tau) \), depends on the test length selected.

A more interesting situation is possible in the operation called "running in." Assume, as before, that the producer controls mean life, \( \mu, \) at a cost of \( h(\mu) \), and that the expected value, \( V(\mu) \), is known. In the "running in" operation the producer must operate items successfully for a period \( \tau \) before declaring them satisfactory and delivering batches of \( N \). The consumer can base the incentive payments on the total number, \( N + x \), "run in" to achieve \( N \) satisfactory items.\(^{13}\) Let the producer's expected cost be the initial production cost of the \( N + x \) items, \( h(\mu)[N + E(x)] \), plus the running in cost of each item, \( c[N + E(x)] \). Let the consumer's problem be that of selecting \( \tau, \mu, \) and \( \phi(x) \) to maximize expected net gain. The only changes in the basic set of equations (3.1)-(3.6), occur in (3.2) which becomes

\(^{13}\) We ignore the possibility of \( x \) becoming larger than some pre-assigned number.
\[
T_p(\mu, \varphi, n) = Ng(\mu) - \frac{N[h(\mu) + c]}{\int_{\tau}^{\infty} f(t|\mu) \, dt}, 14
\]

and in (3.6) which becomes

\[
Ng(\mu') \geq \frac{N[h(\mu') + c]}{\int_{\tau}^{\infty} f(t|\mu') \, dt}
\]

where \( \mu' \) maximizes (3.59) for fixed \( \tau \). For arbitrary life distributions, the prices, \( \varphi_x \), can be obtained as the solution to a linear programming problem analogous to those in previous sections.

**Summary**

In this chapter we have developed a model of the contingent pricing situation which is more realistic than those previously reported. For fixed sample size and quality level we have derived a linear programming problem which yields as its solution an optimal pricing policy, \( \varphi_x \). We have indicated conditions sufficient to insure the existence of a point \( (n,p) \) yielding such a set of prices, and described a procedure for seeking an optimal point \( (n^*,p^*) \). We modified the basic contingent pricing model by requiring that price schedules be piecewise linear, and

\[14\] Since \( x \) is a random variable with a negative binomial distribution, \( E(x) = \frac{Np}{q} \), and since \( p = \int_{\tau}^{\infty} f(t|\mu) \, dt \),

\[
N + E(x) = \frac{N}{q} = \frac{N}{\int_{\tau}^{\infty} f(t|\mu) \, dt}
\]
next changed the objective function to reflect a more realistic procurement situation. Finally uncertainty was permitted in both the cost of attaining an average quality and in the control of the production process, and a brief illustration was included indicating how this contingent pricing model can be applied to any item classifiable as defective or non-defective.
CHAPTER IV
CONTINGENT PRICING POLICIES FOR CONTINUOUS VARIABLES

Introduction

In Chapter III we demonstrated that the model developed for the binomial case could be used to derive contingent pricing policies when the characteristic of interest varied continuously. In this chapter we develop the basic model for the continuous case, and while mathematical programming may be used to approximate the optimal contingent pricing policies, control theory provides interesting insights into the structure of the optimal pricing policies.

Let $\mu$ be some characteristic of interest and let $t$ be a continuous variable that estimates $\mu$. Assume the consumer desires the producer to produce at $\mu_0$. Recall that $h(\cdot)$ is defined to include a unit profit, and let excess profit denote the difference between $\varphi(t)$ and $h(\mu_0)$. By means of the maximum principle of L.S. Pontryagin [50] we show that subject to weak conditions on the distribution of $t$, the form of the pricing policy that both maximizes producer expected profit at $\mu_0$, and minimizes the mean square excess profit at $\mu_0$ is piecewise linear in $t$.

The production and procurement situation we now consider differs from that described in assumptions 1-13, Chapter III, only in the following particulars.
1. Each item produced has a characteristic of interest, $\mu$, which varies continuously in some interval $[\mu_1, \mu_2]$. This characteristic may be true reliability, actual propellant energy, etc.

2. Batches of $N$ items are delivered to a consumer. From each batch $n$ items are selected at random and tested. The result of testing the $i^{th}$ item is $\tau_i$, and a statistic $t(\tau_1, \ldots, \tau_n)$ is determined by the testing, where $t$ lies in $[t_1, t_2]$.

3. The payment schedule is $\varphi(t)$, a function of the statistic $t$.

4. The distribution of $\tau_i$ given $\mu$ is known and consequently $f(t|\mu)$ is known where $f(\cdot)$ denotes a probability density function. When $f(t)$ depends on parameters other than $\mu$ we assume them known and uncontrollable.

5. We assume $V(\mu)$ increasing and concave, and $h(\mu)$ increasing and strictly convex.

6. The consumer desires protection against overpayment when $\mu < \mu_p$, that value of $\mu$ considered poor.

In most acceptance plans and incentive arrangements, $\mu$ is estimated by $t$, and the problem considered is that of selecting the minimum sample size providing discrimination at an appropriate level of confidence between desirable and poor performance. Frequently test costs are high, tests are lengthy or destructive, and it is not feasible to use large sample sizes. This has led to the use of sequential sampling, accelerated life tests, and Bayesian acceptance sampling. The usual payment is either a fixed constant if the test is for the purpose of acceptance, or is generally linearly related to $t$ in incentive contracts.
Basic Model

We assume the consumer seeks a sample size and pricing policy to maximize consumer expected net gain, knowing that given \( n \) and \( \varphi(t) \) the producer will select \( \mu' \) to maximize his own expected profit. Considering assumptions 1-13, Chapter III, and assumptions 1-6, Chapter IV, we have the continuous analogue of (3.1)-(3.6) in

\[
T_c = N_\varphi(\mu) - N_g(\mu) - nc, \quad (4.1)
\]

\[
T_p = N_g(\mu) - N_h(\mu), \quad (4.2)
\]

\[
\text{pr}[N_\varphi(t) \geq w | \mu \geq \mu'] \geq 1 - \alpha, \quad (4.3)
\]

\[
\text{pr}[N_\varphi(t) \leq v | \mu \leq \mu_b] \geq 1 - \beta, \quad (4.4)
\]

\[
\varphi(t_1) \geq m, \quad (4.5)
\]

\[
\varphi(t_2) \leq M, \quad (4.6)
\]

\[
\frac{\partial \varphi(t)}{\partial t} = \dot{\varphi}(t) \geq 0, \quad (4.7)
\]

\[
g(\mu') \geq h(\mu'). \quad (4.8)
\]

Proposition 1

Let \( r(t | \mu, n) = \frac{\partial}{\partial \mu} r(t | \mu, n) \). If for fixed \( n \), \( \mu \), say \((\mu, \mu_0)\) there exists \( \varphi(t) \) maximizing \((4.1)\) and \((4.2)\) at \( \mu_0 \) and for which \((4.3)-(4.6)\) hold, then \( \varphi(t) \) minimizes \((4.1')\) and \((4.2')-(4.6')\) hold.

\[
g(\mu_0) = \int_{t_1}^{t_2} \varphi(t) r(t | \mu_0, n) \, dt \quad (4.1')
\]

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\[(4.2')\]
\[
\int_{t_1}^{t_2} \varphi(t) \ r(t|\mu_o,n) \ dt = h'(\mu_o)
\]

\[(4.3')\]
\[
\varphi(t_\alpha) \geq w, \quad \int_{t_\alpha}^{t_2} f(t|\mu_o,n) \ dt \geq 1 - \alpha
\]

\[(4.4')\]
\[
\varphi(t_\beta) \leq v, \quad \int_{t_1}^{t_\beta} f(t|\mu_b,n) \ dt \geq 1 - \beta
\]

\[
\varphi(t_1) \leq m
\]

\[(4.5')\]

\[
\varphi(t_2) \geq M
\]

\[
\dot{\varphi}(t) \geq 0
\]

\[(4.6')\]
\[
\int_{t_1}^{t_2} \varphi(t) \ r(t|\mu_o,n) \ dt \geq h(\mu_o)
\]

The proof is identical to that of Theorem 1, Chapter III, and is omitted.

Note that approximations to \( \varphi(t) \) are readily obtained by partitioning \([t_1, t_2]\) into a finite number of intervals, and using (3.10) to derive an equivalent linear programming problem containing only five constraints. The approximation can be made extremely precise with little loss in computational efficiency since the computing time required to derive solutions to linear programming problems increases much more rapidly with the number of constraints than the number of variables. Due to the number of constraints the discrete approximation obtained
through linear programming need have at most six distinct price levels. In the next section we will show that the maximum principle also yields the policy, \( \varphi(t) \), but linear programming procedures are more readily available in industrial and government computing centers than are the numerical analysis routines required to determine the exact \( \varphi(t) \) by the maximum principle. Consequently from the point of view of implementing contingent pricing policies, the use of mathematical programming is preferable.

Maximum Principle

Let \( x, p \) denote \( n \)-dimensional vectors, and \( x_i, p_i \) their components. Let \( \lambda, b, x_i^1 \) known constants, \( \varphi(t), \overline{\varphi}(t) \) known functions of \( t \), and \( R, q, \) known functions of their arguments, continuously differentiable with respect to each variable. We seek \( \varphi(t), \varphi(t) \leq \varphi(t) \leq \overline{\varphi}(t) \), which minimizes (4.7) and satisfies (4.10). We will call such a minimizing \( \varphi(t) \) optimal.

\[
(4.7) \quad \varphi = \sum_{i=1}^{n} b_i x_i(t_2)
\]

\[
(4.8) \quad \frac{dx_i(t)}{dt} = \dot{x}_i(t) = q_i(x, \varphi, t), \quad i = 1, \ldots, n
\]

\[
(4.9) \quad x_i(t_1) = x_i^1, \quad i = 1, \ldots, n
\]

\[
(4.10) \quad R(x_1(t_2), \ldots, x_n(t_2)) = 0
\]

\[
(4.11) \quad \dot{p}_i(t) = - \sum_{j=1}^{n} p_j \frac{\partial q_i}{\partial x_1}, \quad i = 1, \ldots, n
\]
\( p_i(t_2) = - \left[ b_i + \lambda_i \frac{\partial R}{\partial x_i(t_2)} \right] \), \( i = 1, \ldots, n \)

(4.13) \[ H(x,p,\varphi,t) = \sum_{i=1}^{n} p_i q_i \]

The maximum principle states that in order for \( \varphi(t) \) to be optimal it is necessary that there exist \( p_1(t), \ldots, p_n(t) \) not all zero, such that \( H(x,p,\varphi,t) \) is maximized at \( \varphi(t) \) for every \( t, t_1 \leq t \leq t_2 \).

When every \( \dot{x}_i(t) \) is linear in \( x(t) \) and contains only an additive function of \( \varphi(t) \) as in (4.14) the stated conditions are also sufficient.

(4.14) \[ \dot{x}_i(t) = \sum_{k=1}^{n} a_{ik} x_k(t) + \omega[\varphi(t)] . \]

**Application of the Maximum Principle**

The maximum principle is not directly applicable to (4.1')-(4.6') due to the constraint \( \dot{\varphi}(t) \geq 0 \), but by a transformation similar to (3.10) we seek \( y(t) \), the derivative of \( \varphi(t) \). Define \( y(t) \) and \( \varphi(t) \) by (4.15)-(4.17) where the integrals are Stieltjes integrals.

(4.15) \[ \int_{t}^{t_2} y(t) \, dt = \int_{t}^{t_2} d\varphi(t) \]

(4.16) \[ \varphi(t_2) = M - y(t_2) \]

(4.17) \[ \varphi(t) = \varphi(t_2) - \int_{t}^{t_2} y(t) \, dt \]

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Define \( F(t), D(t), x_i(t), i = 1, \ldots, 5, \) by

\[
F(t) = \int_{t_1}^{t} f(t|\mu, \nu) \, dt
\]

\[
D(t) = \int_{t_1}^{t} r(t|\mu, \nu) \, dt
\]

\[
x_1(t) = -\int_{t_1}^{t} y(t) \, F(t) \, dt
\]

\[
x_2(t) = -\int_{t_1}^{t} y(t) \, D(t) \, dt
\]

\[
x_3(t) = -\int_{t_1}^{t} y(t) \, J_1(t) \, dt, \text{ where } J_1(t) = 0, t_1 \leq t < t_\alpha,
\]

\[
= 1, t_\alpha \leq t \leq t_2,
\]

\[
x_4(t) = -\int_{t_1}^{t} y(t) \, J_2(t) \, dt, \text{ where } J_2(t) = 0, t_1 \leq t < t_\beta,
\]

\[
= 1, t_\beta \leq t \leq t_2,
\]

\[
x_5(t) = -\int_{t_1}^{t} y(t) \, J_3(t) \, dt, \text{ where } J_3(t) = 1, t_1 \leq t \leq t_2.
\]

We seek \( y(t) \geq 0 \) which maximizes \( x_1(t_2) \) and satisfies

\[
x_1(t_2) \geq h(\mu_0),
\]

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\[ x_2(t_2) = h'(\mu_0) - \phi(t_2), \]
\[ x_3(t_2) \geq w - \phi(t_2), \]
\[ x_4(t_2) \leq v - \phi(t_2), \]
\[ x_5(t_2) \geq m - \phi(t_2). \]

From (4.7), \( b_1 = 1, b_2, \ldots, b_5 = 0. \) The \( \dot{x}_i(t) \) defined by (4.20)-(4.24) are not functions of \( x(t) \). Therefore from (4.11) \( \dot{p}_1(t) = 0, i = 1, \ldots, 5, \) and from (4.12) \( p_1(t) = p_1(t_2). \) Thus \( p_1(t_2) = -[1 + \lambda_1], p_1(t_2) = -\lambda_i, i = 2, \ldots, 5. \) Denoting \( p_1(t) \) by constants \( \psi_i \), we obtain \( H \) from (4.13), and (4.20)-(4.24).

\[ H = \gamma[\psi_1 F + \psi_2 D + \psi_3 d_1 + \psi_4 d_2 + \psi_5 d_3]. \]

The maximum principle states that \( H \) as a function of \( y \) must be maximized at every \( t, t_1 \leq t \leq t_2. \) Abbreviate \( H \) by \( H = yK(t). \)

For \( K(t) < 0, y(t) = 0, \) and for \( K(t) \geq 0, y(t) = +\infty. \) If an optimal \( \phi(t) \) exists there must be \( \psi_1, \ldots, \psi_5 \) such that \( K(t) \) is never positive and is zero at only a finite number of points, \( t_j. \) At these points \( \int_{t_j^-}^{t_j^+} y(t) \, dt = \phi(t_j^+) - \phi(t_j^-), \) the jump in \( \phi(t) \) at \( t_j. \)

The exact number of such impulse points depends on \( f(t). \) For unimodal \( f(t), -D(t) \) is unimodal, and since \( F(t) \) increases in \( t, \)

\footnote{For the remainder of the chapter \( t \) will be understood to be in \([t_1, t_2]\) without additional specification.}
there can be at most five points of tangency with \( K(t) = 0 \). If there are exactly five, one must occur at \( t = t_1 \) due to \( j_3(t) \), one at \( t = t_\alpha \) due to \( j_1(t) \), and one at \( t = t_\beta \) due to \( j_2(t) \). The location of the remaining two depends on \( f(t) \). Thus the maximum principle indicates that \( \phi(t) \) is piecewise constant with no more than six distinct price levels, and yields additional information as to where the price levels may change.

Discontinuous price schedules are unattractive to a producer, which partially accounts for the use of piecewise linear policies in practice. With the problem formulated as in (4.15)-(4.29) we need only add the constraint \( y(t) \leq Y \) where \( Y \) is positive, to replace the jumps in \( \phi(t) \) by linear subarcs. This also reduces the set of points \( (n, \mu) \) on which solutions will exist. This may be thought of as "forcing" the selection of a piecewise linear policy.

From the following theorem we can infer that by seeking agreement to a piecewise linear pricing policy, the consumer is seeking that policy which both minimizes the producer's mean square excess profit at \( \mu_0' \) and maximizes producer expected profit at \( \mu_0 \). We define the following problem which is equivalent to the original contingent pricing problem (4.1)-(4.6) except that for clarity we delete the constraints (4.3), (4.4), and the monotonicity requirement, \( \dot{\phi}(t) \geq 0 \).

\[
(4.31) \quad x_0(t) = \int_{t_1}^{t} [\phi(t) - h(\mu_0)]^2 f(t \mid \mu_0, n) \, dt
\]

\[
(4.32) \quad x_1(t) = \int_{t_1}^{t} \phi(t) r(t \mid \mu_0, n) \, dt
\]

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(4.33)  \[ x_2(t) = \int_{t_1}^{t} \phi(t) f(t|\mu_0, n) \, dt \]

(4.34)  \[ m \leq \phi(t) \leq M \]

(4.35)  \[ x_1(t_2) = h'(\mu_0) \]

(4.36)  \[ x_2(t_2) \geq h(\mu_0) \]

We seek that \( \phi(t) \) satisfying (4.34) which minimizes \( x_2(t_2) \) and yields \( x_1(t_2), x_2(t_2) \) satisfying (4.35), (4.36). We call such a \( \phi(t) \) optimal for (4.31)-(4.36).

**Theorem 1**

If \( \phi(t) \) is optimal for (4.31)-(4.36) and if

\[ \frac{\partial}{\partial \mu} \ln f(t|\mu_0) = A_1(\mu_0) + A_2(\mu_0) t \]

then it is necessary and sufficient that \( \phi(t) \) be a piecewise linear policy of the form (4.37).

\[
\begin{align*}
\phi(t) &= m & t_1 \leq t < t_m \\
\phi(t) &= s(t) & t_m \leq t < t_M \\
\phi(t) &= M & t_M \leq t \leq t_2
\end{align*}
\]

**Proof:**

By (4.7)-(4.13) \( H(x,p,\phi,t) \) for this problem is

(4.38)  \[ H = - [\phi(t) - h(\mu_0)]^2 f(t) + \psi_1 \phi(t) r(t) + \psi_2 \phi(t) f(t) . \]

The \( \psi_1 \) are undetermined constants. The maximum principle states that if \( \phi(t) \) is optimal for (4.31)-(4.36) it is necessary and sufficient
that \( H \) as a function of \( \varphi \) be maximized at every \( t \). If \( \varphi(t) \) is not \( m \) or \( M \), a necessary and sufficient condition for \( H \) to be maximized by \( \varphi \) is

\[
(4.39) \quad \varphi(t) = \left( \frac{\psi_1}{2} \right) \frac{r(t)}{f(t)} + \frac{\psi_2}{2} + h(\mu_0)
\]

since it is easily verified that \( \frac{\partial^2 H}{\partial \varphi^2} < 0 \).

By definition \( r(t) = \frac{\partial}{\partial \mu} f(t|\mu) \), and \( \frac{r(t)}{f(t)} = \frac{\partial}{\partial \mu} \ln f(t|\mu) \) which is linear in \( t \) by hypothesis. Denote the right hand side of (4.39) by \( s(t) \).

If \( s(t) < m \) or \( s(t) > M \) for every \( t \), (4.35) cannot be satisfied for \( h'(\mu_0) > 0 \). Therefore if an optimal policy exists there is at least one \( t \) for which \( m < s(t) < M \). Let \( t_m \) satisfy \( s(t_m) = m \) and let \( t_M \) satisfy \( s(t_M) = M \). (4.37) follows.

**Remarks on Theorem 1**

The statement \( \frac{\partial}{\partial \mu} \ln f(t|\mu) = A_1(\mu) + A_2(\mu) \) implies that \( f(t|\mu) \) is a member of the exponential family of distributions,

\( f(t|\mu) = f(\mu) \exp \{ Q(\mu) U(t) \} h(t) \), but we require additionally that \( U(t) \) be linear in \( t \). The normal, gamma, binomial, poisson, and geometric distributions satisfy this linearity requirement.

By including the constraint \( g(\mu_0) \geq h(\mu_0) \) we introduced the constant term \( \frac{\psi_2}{2} \) into \( \varphi(t) \) but this does not affect the form of the price schedule. There is also no need to reintroduce the constraint \( \varphi(t) \geq 0 \) since by Theorem 1 if an optimal \( \varphi(t) \) exists and \( h'(\mu_0) > 0 \), \( s(t) \) must be positive.
The effect of introducing (4.3), (4.4) is to change the constraint set for \( \varphi(t) \). Jump discontinuities may occur at \( t = t_\alpha \) and \( t = t_\beta \). Since jump discontinuities are not seen in practice we infer that either the consumer and producer do not explicitly specify prices \( w \) and \( v \), or that \( n \) and \( (M - m) \) are sufficiently large so that (4.3') and (4.4') are not active constraints.

There is an interesting connection with estimation which we cannot pursue here.\(^2\) If \( f(t) \) is a member of the exponential family and \( U(t) \) is linear in \( t \), it can be shown that \( t \) is a sufficient statistic for the \( \tau_i \), and that \( t \) is the minimum variance estimator of \( \mu \) among all estimates with the same bias. Therefore if we define the quadratic loss function (4.40), and the risk function (4.41),

\[(4.40) \quad L = (\mu - t)^2, \]
\[(4.41) \quad R(\mu, t) = K(\mu) E(\mu - t)^2 \]

\( t \) is the estimate which minimizes risk.

**Existence of Piecewise Constant Pricing Policies**

In order for a price schedule satisfying (4.2')-(4.6') to exist at some \((n, \mu_0)\) the sample size must be sufficiently large to permit discrimination between \( \mu_0 \) and \( \mu_b \), and the price levels \( M, m, w, \) and \( v \) must be related to \( h(\mu_0) \) in a manner that permits \( \varphi(t) \) to satisfy (4.2')-(4.6'). Proposition 2 states these conditions. The proof is similar to that of Proposition 1, Chapter III and is sketched briefly.

\(^2\)See Lindgren, [51, p. 220].
Proposition 2

Let $n$ and $\mu_o$ be fixed and assume $t_B < t_\alpha < \mu_o$ where $t_\alpha$, $t_B$ are given by (4.3'), (4.4'). Let

$$D_a^b = \int_a^b r(t|\mu, n) \, dt$$

and

$$D_a^b = \int_a^b f(t|\mu, n) \, dt.$$  Then there exists a $\varphi(t)$ satisfying (4.2')-(4.6') at $(n, \mu_o)$ if

$$ (w - v) D_{t_B}^{t_\alpha} \leq h'(\mu_o) \leq (m - M) D_{t_\alpha}^{t_{t_\alpha}} + (w - M) D_{t_B}^{t_\alpha}$$

and

$$h(\mu_o) \leq \lambda [(m-M) D_{t_\alpha}^{t_{t_\alpha}} + (w-M) D_{t_B}^{t_\alpha}] + (1-\lambda) [(v-w) D_{t_B}^{t_{t_\alpha}} + w]$$

hold where $\lambda$ is given by (4.46).

Proof:

Let $g'(\mu_o)$, $g'(\mu_o)$ denote the maximum and minimum values of $g'(\mu_o)$ for any policy, $\varphi(t)$, satisfying (4.3')-(4.5'), and let $g(\mu_o)$, $g(\mu_o)$ denote the expected payment corresponding to these policies.

Since $t_\alpha < \mu_o$, $g'(\mu_o)$ is at most $m D_{t_\alpha}^{t_\alpha} + w D_{t_\alpha}^{t_\alpha} + M D_{t_\alpha}^{t_\alpha}$ or,

$$g'(\mu_o) = (m-M) D_{t_\alpha}^{t_\alpha} + (w-M) D_{t_\alpha}^{t_\alpha}.$$  

By (4.42) this is not less than $h'(\mu_o)$. Similarly

$$g'(\mu_o) = (w-v) D_{t_B}^{t_B}.$$  

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and by (4.42) this is not greater than \( h'(\mu_0) \). Therefore for some \( \lambda \), \( 0 \leq \lambda \leq 1 \), \( h'(\mu_0) = \lambda \bar{g}'(\mu_0) + (1-\lambda) \bar{g}'(\mu_0) \) and thus (4.2') is satisfied. From this we obtain

\[
(4.46) \quad \lambda = \frac{h'(\mu_0) - \bar{g}'(\mu_0)}{\bar{g}'(\mu_0) - \bar{g}'(\mu_0)}.
\]

By (4.43) \( h(\mu_0) \leq \lambda \left[ (m-M) P^\alpha_{\tau \lambda} + (w-M) P^\mu_{\tau \lambda} + M \right] + (1-\lambda) \left[ (v-w) P^\beta_{\tau \lambda} + w \right]. \)

After rearrangement

\[
(4.47) \quad h(\mu_0) \leq \lambda \bar{g}(\mu_0) + (1-\lambda) \bar{g}(\mu_0)
\]

which satisfies (4.6'), and completes the proof.

With the specification that \( t = \frac{\sum_{1}^{n} \tau_i}{n} \) and the assumption that \( t \) is normally distributed, we state conditions sufficient for the convexity of \( R \), the set of points \((n,\mu)\) feasible for (4.2')-(4.6'). The importance of having a convex constraint region is that we can be certain of determining \( (n_0,\mu_0) \) maximizing \( \psi(n,\mu) = NV(\mu) - Nh(\mu) - nc \).

**Proposition 3**

Let \( t = \frac{1}{n} \sum_{1}^{n} \tau_i \) and assume that \( t \) is normally distributed with mean \( \mu \) and known variance \( \frac{1}{n} \sigma^2 \). Let \( t_\alpha < \mu_0, t_\beta > \mu_b \) and let \( n(\mu_0) \) be the minimum \( n \) such that \( t_\beta < t_\alpha \) for \( \mu = \mu_0 \). Then \( t_\beta < t_\alpha \) for every \( n \geq n(\mu_0) \).
Proof:

It suffices to show that $t_\beta$ decreases and $t_\alpha$ increases as $n$ increases. Let $f_{N*}(\cdot)$ denote the standardized normal density function and $F_{N*}(\cdot)$ denote the cumulative function. Let $u_\beta$ satisfy $F_{N*}(u_\beta) = 1 - \beta$ and $u_\alpha$ satisfy $F_{N*}(u_\alpha) = \alpha$.

Since $t_\beta = \frac{u_\beta \sigma}{\sqrt{n}} + \mu_0$ and $u_\beta > 0$ by hypothesis, $t_\beta$ decreases in $n$. Similarly $t_\alpha = \frac{u_\alpha \sigma}{\sqrt{n}} + \mu_0$ and since $u_\alpha < 0$ by hypothesis $t_\alpha$ increases in $n$.

Proposition 4 states conditions under which $R$ is convex in $\mu$ for fixed $n_0$, and in $n$ for fixed $\mu_0$. The proof is identical to that of Proposition 2, Chapter III, and is omitted.

Proposition 4

Let $\underline{\mu}(n_0)$ be the minimum $\mu$ for which both

$$\int_{t_\alpha}^{t_2} f(t|\mu(n_0), n_0) \, dt \geq 1 - \alpha \quad \text{and} \quad t_\alpha > t_\beta.$$ 

Let $\bar{\mu}$ be chosen arbitrarily in $[t_1, t_2]$ so that $h'(\bar{\mu}) < \infty$ and $\mu(n_0) < \bar{\mu}$. Let $\bar{g}'(\mu)$ be minimized at $\mu'$ and $\bar{g}'(\mu)$ be maximized at $\mu''$, where $\mu', \mu'' \in [\underline{\mu}(n_0), \bar{\mu}]$. Then (4.2')-(4.6') can be satisfied at every $(n_0, \mu), \underline{\mu}(n_0) \leq \mu \leq \bar{\mu}$, if

$$\begin{align*}
(4.48) & \quad (m-M) D_{t_1}^{\alpha}(\mu') + (w-M) D_{t_\alpha}^{\bar{\mu}}(\mu') \geq h'(\bar{\mu}), \\
(4.49) & \quad (w-v) D_{t_2}^{\bar{\mu}}(\mu'') \leq h'(\underline{\mu}(n_0)),
\end{align*}$$

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\[(4.50) \quad h(\mu_o) \leq \lambda g(\mu') + (1-\lambda) g(\mu'') ,\]

hold where
\[(4.51) \quad \lambda = \frac{h'(\mu) - \bar{g}'(\mu'')}{\bar{g}'(\mu') - \bar{g}'(\mu'')} .\]

For \( n_2 > n_0 \) and fixed \( \mu_o \) let \( \bar{g}'(\mu_o) \) be minimized at \( n' \), \( n_0 \leq n' \leq n_2 \) and \( \bar{g}'(\mu) \) be maximized at \( n'' \), \( n_0 \leq n'' \leq n_2 \). If \((4.48)-(4.50)\) hold after replacing \( \mu' \) by \( n' \) and \( \mu'' \) by \( n'' \) then \( R \) is convex in \( n \), \( n_0 \leq n \leq n_2 \) for fixed \( \mu_o \).

While under these conditions we can compute \( (n_0, \mu_o) \) maximizing \( \psi(n,\mu) \), there is still little that can be said about the concavity of \( G(n,\mu) \). However the generalized Fibonacci search procedure described in Chapter III will yield better results for the conditions existing in Chapter IV. This is due to the fact that \( R \) is convex over a larger region than in the cases described in Chapter III, allowing us to search \( G(n,\mu) \) without introducing artificial variables, as in \((3.10)\).

**Payment Functions of Several Variables**

In many cases incentives apply to several characteristics simultaneously. An example of this is an incentive on production schedule, equipment reliability, and maintainability. Let \( \mu \) and \( t \) denote vectors, \( \mu^i \), \( t^i \) components, and \( \mu^i \) fixed values. We assume the desired values of \( k \) performance or quality characteristics, \( \mu^1_o, \ldots, \mu^k_o \) are known, that \( \mu^i \) is estimated by \( t^i = \frac{1}{n} \sum_{j=1}^{n} t^i_j \), and that \( f(t | \mu) \) is known.
A procedure frequently observed is the assignment of constant weights, \( b_i > 0, \sum_{i=1}^{k} b_i = 1 \), to the performance characteristics, and definition of the random variable, \( Z = \sum_{i=1}^{k} b_i t_i \). A single policy \( \varphi(Z) \) is then formulated. Assume the distribution of \( t \) to be multivariate normal with known parameters. Then \( Z \) is normally distributed with known mean and variance and the situation is similar to the single variable cases discussed in an earlier section of Chapter IV. For example we seek \( \varphi(Z) \) to minimize

\[
\int_{z_1}^{z_2} [\varphi(Z) - h(\mu_o)]^2 f(Z|\mu_o) \, dz,
\]

subject to

\[
\int_{z_1}^{z_2} \varphi(Z) \frac{\partial f(Z|\mu_o)}{\partial \mu_i} \, dz = \frac{\partial h(\mu_o)}{\partial \mu_i} \quad i = 1, \ldots, k,
\]

\[
\int_{z_1}^{z_2} \varphi(Z) f(Z|\mu_o) \, dz \geq h(\mu_o),
\]

\[
\Psi_{k+1} + \sum_{i=1}^{k} \frac{k}{2} \frac{\partial^2}{\partial \mu_i^2} \ln f(Z|\mu_o). \quad \text{By Theorem 1},
\]

Chapter IV, the optimal policy is given by

\[
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\]
\[ \varphi(z) = m, \quad s(z) < m, \]
\[ \varphi(z) = s(z), \quad m \leq s(z) \leq M, \]
\[ \varphi(z) = M, \quad M < s(z). \]

(4.56)

Since \( Z \) is normally distributed \( s(z) \) is linear in \( Z \). Constraints equivalent to (4.3), (4.4) can be introduced with no essential changes in the formulation, and as an alternative to (4.52) we could introduce a linear objective function to allow computation by means of the linear programming methods of Chapter III. The key simplification is the combination of \( t^1, \ldots, t^k \) into the single variable \( Z \), by consumer specification of the weights, \( b_i \).

It is possible to also include production schedule in \( \varphi(z) \) by defining a variable \( \mu^0 \) related to schedule in such a way that increases in \( \mu^0 \) are desirable. We then assume that if the producer aims at \( \mu^0 \) he attains \( t^0 \) where \( t^0, \ldots, t^k \) have a multivariate normal distribution with \( E(t) = \mu \) and known covariance matrix. Thus the situation is identical to that just described. However, pricing functions are usually formulated separately for performance variables and production schedule. That is \( \varphi(t^0, z) = \varphi_1(t^0) + \varphi_2(z) \). Since
\[
\int_{t^0}^{t^2} \int_{z}^{z_2} \varphi(t^0, z) f(t^0, z|\mu^0, \ldots, \mu^k) \, dt^0 \, dZ = \int_{t^0}^{t^2} \varphi_1(t^0) f_1(t^0) \, dt^0 + \int_{z}^{z_2} \varphi_2(z) f_2(z|\mu_0) \, dZ
\]

we now seek two pricing policies, \( \varphi_1(t^0) \) and \( \varphi_2(z) \) to minimize

(4.57)
\[
\int_{t^1}^{t^2} \varphi_1(t^0) f_1(t^0|\mu_0) \, dt + \int_{z_1}^{z_2} \varphi_2(z) f_2(z|\mu_0) \, dZ,
\]
subject to

\[ \frac{\partial}{\partial \mu_i} \left[ g_{\mu_o}(\mu_o) - h(\mu_o) \right] = 0, \quad i = 1, \ldots, k, \]

\[ \frac{\partial}{\partial \mu_o} \left[ g_{\mu_o}(\mu_o) - h(\mu_o) \right] = 0, \]

\[ g_{\mu_o}(\mu_o) + g_{\mu_o}(\mu_o) \geq h(\mu_o), \]

\[ m_z \leq \varphi_2(z) \leq M_z, \]

\[ m_{t_0} \leq \varphi_1(t_0) \leq M_{t_0}. \]

Note that in this formulation we are not required to choose the weights to be applied to overall performance, \( Z \), and schedule, \( t_0 \). Arbitrarily precise approximations to \( \varphi_1(\cdot) \) and \( \varphi_2(\cdot) \) can be obtained by the methods of Chapter III.

In many situations payment based on a constant weighting of the \( \mu_i \) does not reflect the consumer desires outside a small region about \( \mu_o \), and there may be considerable uncertainty concerning the correct weights within that region. We require a price schedule which is a joint function of the variables \( t_1, \ldots, t_k \). Specifically, consider the determination of a price function \( \varphi(t_1, t_2) \) where \( t_1, t_2 \) estimate different performance or quality characteristics, and the joint distribution of \( t_1 \) and \( t_2 \) is known. We seek \( \varphi(t_1, t_2) \) minimizing

\[ \int_{t_1}^{t_2} \int_{t_1}^{t_2} \varphi(t_1, t_2) f(t_1, t_2 | \mu_o, \mu_o) \, dt_1 \, dt_2, \]

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subject to

\[(4.64) \quad \frac{\partial g(\mu^l_o, \mu^2_o)}{\partial \mu^1_i} = \frac{\partial h(\mu^l_o, \mu^2_o)}{\partial \mu^1_i}, \quad i = 1, 2,\]

\[(4.65) \quad m \leq \varphi(t^1, t^2) \leq M,\]

\[(4.66) \quad g(\mu^l_o, \mu^2_o) \geq h(\mu^l_o, \mu^2_o),\]

\[(4.67) \quad \frac{\partial \varphi(t^1, t^2)}{\partial t^1_i} \geq 0, \quad i = 1, 2.\]

By partitioning the interval \([t^1_{l_1}, t^1_{l_2}]\) into \(v_1\) subintervals and the interval \([t^2_{l_1}, t^2_{l_2}]\) into \(v_2\) subintervals, we have the following approximating linear programming problem where \(t^1 = t^1_{l_1} + i(t^1_{l_2} - t^1_{l_1})/v_1,\)
\(t^2 = t^2_{l_1} + j(t^2_{l_2} - t^2_{l_1})/v_2.\) Minimize

\[(4.68) \quad \sum_{i=0}^{v_1} \sum_{j=0}^{v_2} \varphi_{ij} f^i_j(\mu^l_o, \mu^2_o),\]

subject to

\[(4.69) \quad \sum_{i=0}^{v_1} \sum_{j=0}^{v_2} \varphi_{ij} \frac{\partial f^i_j}{\partial \mu^1_i} = \frac{\partial h(\mu^l_o, \mu^2_o)}{\partial \mu^1_i}, \quad i = 1, 2,\]

\[(4.70) \quad g(\mu^l_o, \mu^2_o) \geq h(\mu^l_o, \mu^2_o),\]

\[(4.71) \quad m \leq \varphi_{0j} \leq \varphi_{1j} \leq \cdots \leq \varphi_{v_1j} \leq M, \quad j = 0, \ldots, v_2,\]
\[ m \leq \Phi_{10} \leq \Phi_{11} \leq \ldots \leq \Phi_{1v_2} \leq M, \quad i = 0, \ldots, v_1. \]

This problem can be solved by standard linear programming methods, but the computation required has increased considerably. Even a relatively coarse partition of \([t_1^1, t_2^1]\) and \([t_1^2, t_2^2]\) into 20 subintervals results in a problem containing 43 constraints and 441 variables, indicating extension to more than two variables is impractical. However the important two variable problem is that of determining an optimal joint price function of schedule performance, \(t^0\), and the overall measure of performance and quality, \(z\). This joint price function can be derived by the method just outlined.

A realistic and tractable alternative to seeking an optimal \(\varphi(t^1, \ldots, t^k)\) from the class of all bounded, monotonic, functions is to seek that \(\varphi(t^1, \ldots, t^k)\) yielding an expected payment satisfying specified constraints. One natural and desirable expected payment is

\[ g(\mu) = V(\mu) - [V(\mu_o) - h(\mu_o)]. \]

If the actual outcome \((t^1, \ldots, t^k)\) is expected to be close to \(\mu_o\), or if \(V(\mu)\) is not changing rapidly, \(g(\mu)\) can be approximated adequately by the first three terms of a Taylor series as

\[ g(\mu) \approx h(\mu_o) + \sum_{i=1}^{k} \frac{\partial V(\mu_o)}{\partial \mu_i^1} (\mu_i^1 - \mu_o^i) + \frac{1}{2} \sum_{i,j=1}^{k} (\mu_i^1 - \mu_o^i) \frac{\partial^2 V(\mu_o)}{\partial \mu_i^1 \partial \mu_j^1} (\mu_j^1 - \mu_o^j). \]
Under the assumption that the distribution of $t$ is multivariate normal, and that the covariance matrix $G = [g_{ij}]$ is known, $\varphi(t)$ need only be a quadratic function of the variables $t^i$ to yield $E[\varphi(t)] - g(\mu)$. For the $g(\mu)$ above,

$$\varphi(t) = a + \sum_{i=1}^{k} b_it^i + \sum_{i,j=1}^{k} c_{ij}t^it^j,$$

where

$$a = h(\mu_0) - \sum_{i=1}^{k} \frac{\partial V(\mu)}{\partial \mu_i} \mu_i^0 + \frac{1}{2} \sum_{i,j=1}^{k} (\mu_i^0 \mu_j^0 g_{ij}) \frac{\partial^2 V(\mu)}{\partial \mu_i \partial \mu_j},$$

$$b_i = \frac{\partial V(\mu)}{\partial \mu_i} - \sum_{j=1}^{k} \mu_j^0 \frac{\partial^2 V(\mu)}{\partial \mu_i \partial \mu_j},$$

and

$$c_{ij} = \frac{1}{2} \frac{\partial^2 V(\mu)}{\partial \mu_i \partial \mu_j}.$$

**Piecewise Linear Price Schedules**

We refer to a piecewise linear pricing policy of the form (4.79) as a "linear" policy, and to a piecewise constant pricing policy satisfying (4.2')-(4.6') as a "basic" policy.

$$\varphi(t) = m_0 \geq m \quad t \leq a$$

$$\varphi(t) = m_0 + (t-a)s \quad a \leq t \leq b$$

$$\varphi(t) = M_0 \leq M \quad t \geq b$$

$$s = \frac{M_0 - m_0}{b - a}$$
In Chapter III we described a procedure for deriving linear policies, and in Chapter IV indicated the criterion that leads naturally to such policies. The attractive features of linear policies are continuity and ease of interpretation, however linear policies can never be more profitable for the consumer than basic policies, and in this section we investigate the extra cost incurred by the consumer resulting from the use of an optimal linear policy.

Let $\varphi(t)$ denote a basic pricing policy, and $\varphi_k(t)$ a linear pricing policy. Let $g(\mu)$, $g_k(\mu)$ denote the expected payments, and $g'(\mu)$, $g'_k(\mu)$ denote the derivatives of the expected payments. The lesser profitability of the linear policies occurs in three ways.

If at fixed $(n, \mu)$ optimal linear and basic policies exist, it may be that $g_k(\mu) > g(\mu)$, and the loss to the consumer is $g_k(\mu) - g(\mu)$.

Neglecting (4.3'), (4.4'), it is possible that at some $(n, \mu)$ a linear policy cannot simultaneously satisfy (4.2') and (4.5') while a basic policy can. This is brought out by Proposition 5. One measure of the extra cost to the consumer in this event is the additional sampling required to increase $(n, \mu)$ to $(n', \mu)$, a point at which (4.2') and (4.5') can be simultaneously satisfied.

It is also apparent that even if a linear policy does satisfy (4.2') and (4.5') at $(n, \mu)$, constraints (4.3'), (4.4') may be active, and may prevent the existence of a linear policy. An increase in sample size also measures the loss in this case. In the remainder of this section we derive statements indicating the amount of extra sampling required for the existence of optimal linear policies as a function of the negotiated parameters.
Proposition 5

Recall that by $\bar{g}(\cdot)$, $\bar{g}(\cdot)$ we denote minimum and maximum expected payments. Let $t_\alpha < t_\mu$ and $t_\beta < t_\alpha$. If $a \leq t_\alpha$ and $b \geq a$ then $\bar{g}'(\mu) \leq \bar{g}'(\mu)$, $\bar{g}'(\mu) \geq \bar{g}'(\mu)$, and by definition of $g'(\mu)$ if more than one $q_\beta(t)$ is feasible then $\bar{g}'(\mu) < \bar{g}'(\mu)$.

Proof:

The maximum value of $g'(\mu)$ is

\[(4.30) \quad \bar{g}'(\mu) = m \int_{t_1}^{t_\alpha} r(t|\mu) \, dt + w \int_{t_\alpha}^{t_\mu} r(t|\mu) \, dt + M \int_{t_\mu}^{t_2} r(t|\mu) \, dt .\]

At that $(a,b)$ which maximizes $\bar{g}'(\mu)$

\[(4.31) \quad \bar{g}'(\mu) = \bar{g}'(\mu) + \int_a^{t_\alpha} s(t-a) \, r(t|\mu) \, dt + \int_{t_\alpha}^{t_\mu} [\varphi(t_\alpha) - w + s(t-t_\alpha)] \, r(t|\mu) \, dt + \int_{t_\mu}^{b} [M - \varphi(\mu) + s(t-\mu)] r(t|\mu) \, dt .\]

Since $s > 0$ and $r(t) \leq 0$ for $t < \mu$ the first integral is non-positive. By (4.3') $\varphi(t_\alpha) \geq w$ and the second integral is non-positive. By (4.5') $\varphi(\mu) \leq M$ and the third integral is non-positive, yielding

\[(4.32) \quad \bar{g}'(\mu) \leq \bar{g}'(\mu) .\]
Similarly

\[ g'(\mu) = v \int_{t_2}^{t_B} r(t|\mu) \, dt + w \int_{t_B}^{t_2} r(t|\mu) \, dt. \]

If \( a = b = t_B \), \( g'_A(\mu) = g'(\mu) \). If \( a \) is constrained to be less than \( b \), let \( s = \frac{w - m}{b - a} \) and consider the case in which \( b \leq \mu \) and \( \varphi(b) = w \).

For that \( a \leq t_B \) at which \( g'_A(\mu) \) is minimized we have

\[ g'_A(\mu) = g'(\mu) + \int_{a}^{t_B} [m - v + s(t-a)] \, r(t|\mu) \, dt \]

\[ + \int_{t_B}^{b} [\varphi(t_B) - w + s(t-t_B)] \, r(t|\mu) \, dt. \]

By (4.4) \( \varphi(t_B) \leq v \) and \( m + s(t_B-a) = \varphi(t_B) \). Thus the first integral is non-negative. Also \( \varphi(t_B) + s(b-t_B) = \varphi(b) = w \) and the second integral is non-negative, yielding

\[ g'(\mu) \geq g'_A(\mu) \]

and completing the proof.

We will now assume that \( g'_A(\mu) \leq h'(\mu) \), that \( m_0 = m \) and \( M_0 = M \) in (4.79), that \( t = \frac{1}{n} \sum_{i=1}^{n} \tau_i \), and that \( t \) is normally distributed with mean \( \mu \) and variance \( \frac{1}{n} \sigma^2 \). Under these assumptions we derive the minimum sample size, \( \bar{n} \), at which a linear policy satisfying (4.3')-(4.5') can yield equality in (4.2'). We then compare \( \bar{n} \) with the minimum sample size required by an optimal basic policy. From (4.79)
\[ g_\mu (\mu) = \int_{-\infty}^{a} m(t|\mu) \, dt + \int_{a}^{b} (m+s(t-a)) \, f(t|\mu) \, dt + \int_{b}^{\infty} M(t|\mu) \, dt , \]

or

\[ g_\mu (\mu) = M + s(a-\mu) \, F(a) - s(b-\mu) \, F(b) - s\sigma [f(b) - f(a)] . \]

Differentiating with respect to \( \mu \) we have

\[ g_\mu (\mu) = s[F(b) - F(a)] . \]

We seek \((a^*, b^*)\) to maximize (4.88). While it appears as if both \( s \) and either \( a \) or \( b \) are independent, a moment's reflection shows that if two policies have identical \( s \), and both satisfy \( \varphi(t_\alpha) \geq w \), that policy with the greater \( a \) will have the greater \( g_\mu (\mu) \). Therefore \( \varphi(t_\alpha) = w \) and thus

\[ a = t_\alpha + \frac{m - w}{s} , \]

\[ b = t_\alpha + \frac{M - w}{s} , \]

demonstrating that (4.88) is a function of only one variable. If \( s^* \) maximizes (4.88) we note that at the minimum sample size, \( n \), for which \( g_\mu (\mu) = h'(\mu) \), we must have \( g_\mu (\mu) = \overline{g}_\mu (\mu) \), and from (4.88) we have

\[ s^*[F(b) - F(a)] = h'(\mu_0) . \]
Let \( u_a = \frac{a - \mu}{\sigma_t} \), \( u_p = \frac{b - \mu}{\sigma_t} \), and \( u_\alpha = \frac{t_\alpha - \mu}{\sigma_t} \). After rearrangement (4.91) becomes

\[
(4.92) \quad \frac{\sqrt{n}}{u_\alpha - u_a} \left[ F_{N^*}(r_1 u_\alpha - r_2 u_a) - F_{N^*}(u_a) \right] = \frac{h'(\mu) \sigma_t}{w - m}
\]

where \( r_1 = \frac{M - m}{w - m} \), \( r_2 = r_1 - 1 \), and the subscript \( N^* \) denotes the standardized normal distribution. With only a table of the cumulative normal density function we can solve (4.92) for both the minimum \( n \), and that \( a \) which yield equality in (4.2') for a policy satisfying (4.3') and (4.5').

We derive approximations to \( n \) in terms of the parameters. By letting \( a = b = t_\alpha \) in (4.86) we obtain

\[
(4.93) \quad g'_\alpha(\mu) = \frac{\sqrt{n}}{\sigma_t} (M - m) f_{N^*}(u_\alpha),
\]

and since we seek the minimum \( n \) for which \( g'_\alpha(\mu) \geq h'(\mu) \), (4.93) yields

\[
(4.94) \quad \bar{n}_1 \geq \frac{\sigma_t h'(\mu)}{(M - m) f_{N^*}(u_\alpha)}.
\]

an upper bound to \( n \). The bound approaches \( n \) as \( w \) approaches \( M \).

We now obtain an approximation to \( n \) that is more accurate as \( w \) approaches \( m \). By setting \( b = \mu \) in (4.88),

\[
(4.95) \quad g'_\alpha(\mu) = \frac{M - w}{\mu - t_\alpha} \left[ \frac{1}{2} - F_{N^*}\left( u_\alpha \left[ M - m \right] \right) \right].
\]
From this we derive

\[(4.96)\quad n_2' \geq \left[ \frac{\sigma_t^2 h'(\mu)}{(W-M) \left[ \frac{1}{2} - \Phi_{N^*} \left( N^*_\alpha \left( \frac{M-m}{M-w} \right) \right] \right]} \right]^2.\]

The \( n \) determined by (4.92) is constrained by (4.4'), \( \psi(t_\beta) \leq v \).

Thus at both \( n \) and \( n_1 \) (4.2')-(4.5') hold, and we modify \( r_2' \) so that the linear policy obtained by setting \( b = \mu \) also satisfies (4.4').

The statement \( \psi(t_\beta) \leq v \) implies \( m + (t_\beta - a) \left[ \frac{v-m}{t_\alpha - a} \right] \leq v \), from which

\[(4.97)\quad n_2'' \geq \frac{\sigma_t^2}{(\mu_0 - \mu_b)^2} \left[ \frac{\sigma_t}{\mu_0 - \mu_b} \left( \frac{M-v}{M-w} \right) \right]^2.\]

Let \( n_2 = \max(n_2', n_2'') \). This is an upper bound on \( n \), and therefore for given parameters we choose \( \min(n_1, n_2) \) as the upper bound on \( n \).

One lower bound for \( n \) is given by \( n_3 \), the minimum sample size required for the existence of a basic policy satisfying (4.2')-(4.5').

From (4.80) \( \bar{g}'(\mu) = \frac{\sqrt{n}}{\sigma_t} \left[ (M-w) \cdot f_{N^*}(0) + (W-m) \cdot f_{N^*}(\alpha) \right] \), and since we seek \( \bar{g}'(\mu) = h'(\mu) \),

\[(4.98)\quad n_3 \geq \left[ \frac{\sigma_t h'(\mu)}{(M-w) f_{N^*}(0) + (W-m) f_{N^*}(\alpha)} \right]^2.\]

Finally the minimum sample size for which \( t_\beta < t_\alpha \) is given by

\[(4.99)\quad n_4 \geq \frac{s^2}{\sqrt{\frac{\mu_b - \mu_\alpha}{\mu_0 - \mu_b}}}.\]
The maximum of \( n_3 \), \( n_4 \) furnishes a lower bound on \( n \).

For convenience we collect the ratios \( \frac{n_1}{n_3} \), \( \frac{n_1'}{n_3} \), \( \frac{n_2''}{n_4} \) below. The ratio \( \frac{n_2''}{n_4} \) is an approximate measure of the extra sampling required by a linear policy to satisfy the price discrimination required by (4.3) and (4.4).

\[
\frac{n_1}{n_3} = \left[ \frac{\frac{2}{k+1} + \frac{r_{N}^{*}(u_{\alpha})}{r_{N}^{*}(u_{\beta})}}{r_{N}^{*}(u_{\alpha})} \right]^2, \quad k = \frac{M - w}{w - m}.
\]

\[
\frac{n_1'}{n_3} = \left[ \frac{\frac{2}{k+1} + \frac{r_{N}^{*}(u_{\alpha})}{r_{N}^{*}(u_{\beta})}}{r_{N}^{*}(u_{\alpha})} \right]^2, \quad k = \frac{M - w}{w - m}.
\]

\[
\frac{n_2''}{n_4} = \left[ \frac{u_{\beta} - u_{\alpha}}{u_{\beta} - u_{\alpha}} \right]^2, \quad r = \frac{M - v}{M - w}.
\]

Figure 1 displays the minimum of \( \frac{n_1}{n_3} \), \( \frac{n_1'}{n_3} \), \( \frac{n_2''}{n_4} \) as a function of \( k = \frac{M - w}{w - m} \) and \( \alpha \). This indicates approximately the degree of extra sampling required for linear policies to provide the same motivation as do basic policies. The figure overstates the amount of extra sampling required in all cases.

Thus under the assumptions that \( t \) is a sample mean, normally distributed with known variance, and that the linear policy is obtained by negotiating the maximum and minimum payments and the points \( a, b, \) at which these payments occur, we can calculate from (4.92) the minimum sample size required for the existence of an optimal linear pricing policy. Comparing this to the sample size required for the existence
of an optimal basic pricing policy provides one measure of the cost to the consumer of using optimal linear policies. Figure 1 indicates that we will not be very inaccurate by stating that linear policies will generally require from 1.1 to 1.6 times the sample size required by an optimal basic policy.

If linear policies are negotiated without consumer knowledge of the sampling distribution and production cost function, the expected producer profit will not in general be maximized at the desired performance level, \( \mu_0 \). Let \( a, b \) be selected symmetrically about \( \mu_0 \), and let
Let $M, m$, be selected symmetrically about $h(\mu_0)$. If the sampling distribution is symmetric about $\mu_0$, then $g(\mu_0) = h(\mu_0)$. If the probability of $t$ falling outside $[a, b]$ is negligible then $g'(\mu) \approx m + \frac{M-m}{b-a} \mu$, and by selecting $\frac{M-m}{b-a} = h'(\mu_0)$, motivation at $\mu_0$ is possible. But if the probability of $t$ falling outside the incentive region is not negligible, then $g(\mu) < m + h'(\mu_0) \mu$ for $\mu > \mu_0$ and $g(\mu) > m + h'(\mu_0) \mu$ for $\mu < \mu_0$, leading to a much weaker profit motivation in the case of convex production costs than the linearity of $\varphi(t)$ suggests. If the consumer intends to use a piecewise linear pricing policy, he must negotiate to keep $(b-a)$ large to compensate for lack of information describing the sampling distribution, and can reduce the sample size required to maximize expected producer profit at $\mu_0$ by keeping $M - m$ large relative to $\sigma_t h'(\mu_0)$. 

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CHAPTER V

CONCLUSIONS

We briefly review the factors leading to the contingent pricing model developed in this paper, summarize the results obtained, and comment on the implications of the results in negotiating or formulating contingent pricing contracts. In discussing the use of contractual incentives we noted that in addition to expected profit, a contractor or producer considers the risk of financial loss, the maintenance of corporate reputation, and the desire for both program continuation and future contracts in making program decisions. We concluded that expected profit alone is an uncertain motivator and that while the overall effect of contractual incentives on industrial efficiency has been beneficial, the government as consumer cannot predict the program decisions of the contractor on the basis of contractual profit alone. Thus incentive arrangements can be analyzed for consistency with government objectives, but not for motivation. Of course as sharing proportions increase, the relative weight of profit in program decisions will increase.

Following these qualitative considerations we concentrated on the specific and idealized problem of formulating consistent incentive payment plans, or contingent pricing policies, when the payment is based on the result of testing or inspecting a sample of the items contracted for. We reviewed two models which deal specifically with situations of this type, but which do not yield sufficiently realistic pricing policies.
We suggested that more realistic policies could be derived by assuming that while the consumer and producer each attempt to negotiate a contingent pricing contract which maximizes their individual expected net gains, both require that the contract include acceptable protection against possible overpayment or underpayment due to sampling variation. We postulate a model in which \( p \), the fraction of defective items in submitted batches, is controllable by the producer at known cost, \( h(p) \), and in which payments are based on the number of defectives observed in a sample drawn from each batch. We defined an optimal pricing function as one which motivates the producer to select that \( p_o \) desired by the consumer, yields a producer expected profit at \( p_o \) no less than that agreed on, and satisfies restrictions against the risk of incorrect payment. If an optimal price schedule exists at a fixed sample size, \( n \), and quality level, \( p \), it can be determined by linear programming. The linear programming formulation indicates that a price schedule optimal by our definition is a piecewise constant function with the number of distinct price levels related to the number of restrictions the producer and consumer include to control the risk of incorrect payment. We present several examples of such policies in Chapter III, and refer to these piecewise constant pricing policies as "basic" policies.

We next restricted pricing policies to be linear in the number of defectives observed within some interval \([a, b]\) and constant outside that interval. We refer to such piecewise linear policies as "linear," and illustrated with an example that optimal linear policies can be derived by solving a sequence of linear programming problems. We indicated briefly how the basic model could be extended to situations in
which production control and production cost are not precise, but can
be described in terms of known probability distributions. When considering
several variables, additive, independent, incentive payment schedules
do not adequately reflect the consumer's valuation of alternative out-
comes outside a small region. A joint pricing function of all of the
relevant variables is required rather than a linear combination of
independent payments. In Chapter IV we demonstrated that if the joint
density function of the sample outcomes is known, joint pricing functions
can be determined by linear programming, but that this seemed computa-
tionally impractical for more than two variables. A more realistic and
tractable approach is to define the desired expected payment function
in terms of the performance parameters of interest, and seek a pricing
function of known form such that the expected value of this pricing
function approximates the desired expected value function.

We modified the basic model to deal with test outcomes which are
continuous variables. In such cases approximations to optimal price
schedules can be computed by linear programming, but the maximum prin-
ciple of control theory yields the information that the form of the
optimal price function, $\varphi(t)$, is piecewise constant, and indicates the
points at which the price levels in $\varphi(t)$ can shift. More important,
the maximum principle enables us to explain the widespread use of linear
pricing functions in terms of a natural criterion. Assume that the
consumer desires the producer to produce at some quality or performance
level, $\mu^*_0$, and that $t$ estimates $\mu$. Let "excess" profit denote the
difference between actual profit received by the producer and the
negotiated target profit at $\mu^*_0$. Subject to very weak conditions on the
distribution of \( t \), that pricing policy which both maximizes producer
expected profit at \( \mu_o \) and minimizes mean square excess profit at \( \mu_o \)
is piecewise linear in \( t \). If the risk of incorrect payment due to
sampling variation is not explicitly considered, this result implies
that the piecewise linear policies in current use are in fact optimal
under a very natural criterion.

If the consumer uses linear policies that explicitly take account
of the risk of overpayment, he may incur extra costs. At some \((n, \mu_o)\)
a linear policy capable of motivating production at the desired \( \mu_o \),
satisfying a profit non-negativity constraint at \( \mu_o \), and remaining
within specified bounds, may not drive expected excess producer profit
as low as an optimal basic policy. Letting \( g_x(\mu) \) and \( g(\mu) \) denote
the expected payments of the linear and basic policies, and \( h(\mu) + z \)
denote the production cost at \( \mu \) plus the agreed on profit at \( \mu \), the
situation just described yields a resultant extra cost to the consumer
of \( g_x(\mu_o) - g(\mu_o) \). The occurrence of this situation is related to
the magnitude of \((M-m)\), the difference between the maximum and minimum
bounds within which the actual payment schedule is negotiated. As this
difference increases, the loss given by \( g_x(\mu_o) - g(\mu) \) decreases.

Extra cost may arise if at some fixed sample size, expected pro-
ducer profit cannot be maximized at \( \mu_o \) by a linear policy. Proposi-
tion 5, Chapter IV, states that a basic policy can always motivate pro-
duction at a specified point, \( \mu_o \), with a smaller sample size than can
a linear policy. The exact minimum sample size required for the linear
policy can be computed by (4.92) for specified parameters. Figure 1,
Chapter IV presents an approximate ratio of the sample sizes required

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for optimal linear and basic policies to motivate production at some \( \mu_0 \). We have defined \( w \) to be that payment which the producer is assured of receiving with probability \( 1 - \alpha \) if his quality or performance is at least as good as \( \mu_0 \). Let \( \sigma^2 \) be the variance of a single test or inspection outcome. The sample size required for motivation at \( \mu_0 \) is related to the derivative of the production cost function at \( \mu_0 \), \( h'(\mu_0) \).

We can say approximately that when \( \tau \) is an observation and \( t \) is a sample mean, the sample size required for the existence of an optimal policy linear as a function of \( t \) is directly proportional to \( \sigma_t h'(\mu_0) \), and inversely proportional to both \( (M-w) \) and \( (M-m) \). Therefore it behooves the consumer to negotiate for large \( (M-m) \) and the lowest price level \( w \) to which the producer will agree.

If the consumer insists that he will pay no more than \( v \) with probability \( 1 - \beta \) when true quality is no better than \( \mu_B \), a quality level designated as poor, and the producer insists on a high degree of assurance that he will receive at least \( w \) for appropriate performance, the linear policy will not provide the required price discrimination at as small a sample size as the basic policy. An approximate ratio of the sample sizes required for the linear and basic policies due to the need for price discrimination is

\[
\text{sample size required for linear} = \left( \frac{u_B - \frac{M-V}{M-w}}{u_B - u_\alpha} \right)^2 \frac{u_B - \frac{M-V}{M-w}}{u_B - u_\alpha} 
\]

where \( u_B \) and \( u_\alpha \) are the \( 1 - \beta \) and \( \alpha \) percentage points of the cumulative standardized normal density function. From (5.1) we see that the sample size required for the existence of an optimal linear policy
increases as the degree of price discrimination required in the incentive payment schedule increases. When price discrimination is not explicitly considered, the sample size depends only on \((M-m, \sigma_r, h'(\mu_0))\), and the confidence with which the consumer desires to distinguish between alternative performance levels.

Recall that \(a\) and \(b\) are the endpoints of the linear section of the policy \(\varphi(t)\). Assume the sampling distribution is symmetric about its mean, and that \(a, b\) are selected symmetrically about the sample outcome estimating the desired performance level \(\mu_0\). If \((b-a)\) is small relative to \(\sigma_r\), the linear portion of the policy does not apply over much of the region in which the sample outcomes are expected. The result is much weaker producer profit motivation toward high performance, and less financial deterrence from low performance than the linearity of \(\varphi(t)\) suggests. When sampling distributions are not known precisely the consumer must attempt to retain linearity over a wide region, by keeping \((b-a)\) large, in order to have any estimate of the profit motivation acting on the producer.

We conclude that piecewise linear pricing policies can yield an expected producer profit which is maximized at the performance level desired by the consumer, and in fact are optimal policies under one criterion. While linear policies are readily interpreted by the producer and have the desirable practical property of continuity, their use may require larger sample sizes and consequent extra cost. Moreover, at a fixed sample size producer expected profit resulting from a piecewise linear policy does not decrease as sharply for performance below \(\mu_0\) as does the expected profit resulting from a basic policy.
Therefore if a basic policy is unacceptable to a producer, we recommend that the linear policy negotiated include more than one linear section, with the slopes of successive linear sections decreasing. The object, of course, is to create a concave expected payment function in the neighborhood of the desired production point, thereby decreasing the financial attractiveness of performance below that desired by the consumer.

Throughout the paper we have assumed the consumer can state product value as a function of the relevant product characteristics and production or development schedule. Such knowledge does not customarily exist in functional form, and the derivation of optimal contingent pricing policies does not depend on this assumption. The essential requirements for the derivation of contingent pricing policies are knowledge of the sampling distribution and knowledge of the production cost function in the neighborhood of the performance or quality level desired. The consumer may of course derive contingent pricing policies at a variety of sample sizes and performance levels, and then select the preferred policy on economic and other grounds.

Contingent pricing policies can be implemented with little problem since producers have long been familiar with acceptance sampling, and are becoming increasingly familiar with incentive contracts. Where existing procurement contracts call for acceptance sampling, the AQL, producer's risk, LTPD, and consumer's risk can be inferred from the sampling plan in use. Starting with the requirement that the "w" of this paper be the price currently paid when batches are accepted, the contingent pricing policy can readily be computed by the methods of
Chapter III. In new procurement actions where no acceptance sampling arrangements exist, the problems involved in negotiating the parameters of a contingent pricing policy do not appear more intractable than those required to determine the proper payment and acceptance sampling plan. The only additional essential item of information required is the production cost function in the neighborhood of the desired AQL, and this does not seem to be a sufficiently damaging disclosure to preclude negotiation of such contracts.

Furthermore there may be no a priori reason to prefer one price level \( w \) or \( v \) to another. Therefore while \( M \) and \( m \) are to be negotiated, the sample size \( n \) can be determined by the specification of \( \mu_0, \mu_b \), and the probabilities of erroneous statements about the parameters. The prices \( w \) and \( v \) are then determined by \( \varphi(t) \). This has the effect of creating a more symmetric policy about \( \mu_0 \).

As we remarked in Chapter II, contingent pricing policies provide attractions to both consumer and producer. When the consumer cannot reject items due to urgent operational need or planned production requirements, a single acceptance price does not provide a sufficiently powerful tool to motivate superior producer effort. If the producer purchases the same item from several producers, all may submit acceptable items, but contingent pricing policies allow the conscientious producer to be rewarded with an immediate increase in dollar payments rather than a slightly greater percentage of accepted lots in the long run. Lots previously rejected by the consumer and scrapped may in some instances be accepted at reduced cost, thereby averting a total loss for the producer.
The approach developed in this paper can be used to test existing or proposed incentives for consistency with consumer objectives, and to formulate optimal contingent pricing policies in situations where true performance must be inferred from sample data.
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Contingent Pricing Policies

Technical Report, December 1964

Durbin, Eugene P.

Release without limitations on dissemination

A production and procurement situation is described in which the producer and consumer of some commodity negotiate a contract under which the consumer purchases a batch of \( N \) items and pays a unit price, \( \phi(x) \), determined by the result of sampling \( n (\leq N) \) items from the batch. It is assumed that the producer can control the quality, \( p \), of the commodity to a known degree at a known cost, \( h(p) \), and further that given \( \phi(x) \) and \( n \) the producer selects \( p \) to maximize expected net return. Therefore the consumer must choose \( \phi(x) \) and \( n \) to maximize his own net return under the producer's optimal action. Under the additional assumptions that \( \phi(x) \) is bounded and that the contract must provide for protection against incorrect payment due to sampling variation it is shown that at a fixed \((n,p)\) optimal price schedules must be solutions to a specified linear programming problem, and that for fixed \( n \), if the consumer desires the producer to select \( p = p_0 \), and seeks \( \phi(x) \) to both maximize producer expected net return at \( p_0 \) and minimize the mean square difference between actual producer profit at \( p_0 \) and negotiated profit, under weak conditions the optimal \( \phi(x) \) is piecewise linear. A computational procedure for seeking optimal price functions is described, and several hypothetical examples are presented.
## Key Words

- Contracts, Incentive
- Acceptance Sampling
- Control Theory
- Game Theory
- Mathematical Programming
- Statistics: Inference

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